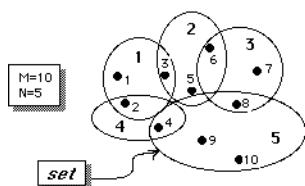


Given M points, and N sets each containing one or more points:

Let C_j = cost of set #j

$$a_{ij} = \begin{cases} 1 & \text{if point } i \text{ is an element of set } j \\ 0 & \text{otherwise} \end{cases}$$



$$A = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

Applications

- = potential facility site
- = customer location
- C_j = cost of building facility j

Facility Location

Facilities must be selected to serve every customer. Possible links are indicated.
(customers are "points", facilities are "sets")

- Problem definition & formulation
- Applications
- Lagrangian relaxation
- Lagrangian dual problem
- Solving dual by subgradient method
- Solving dual by dual ascent method
- Heuristic based on Lagrangian relaxation
- Eliminating sets during Lagrangian relaxation
- Example: subgradient optimization
- Example: dual ascent
- Computational results

Define variables $X_j = \begin{cases} 1 & \text{if set } #j \text{ is selected} \\ 0 & \text{otherwise} \end{cases}$

Set Covering Problem

$$\begin{aligned} \text{Minimize} \quad & \sum_{j=1}^N C_j X_j \\ \text{subject to} \quad & \sum_{j=1}^N a_{ij} X_j \geq 1 \quad \text{for each } i=1, 2, \dots, M \\ & X_j \in \{0,1\} \quad \text{for each } j=1, 2, \dots, N \end{aligned}$$

Applications

Information Retrieval

Retrieve a given set of m requests for information from a set of n files so that the length of the search is minimized.

C_j = length of file j

$a_{ij} = 1$ if the i^{th} information requested is in file j ,
0 otherwise

(information requests are "points", and the files are the "sets")