Notes on Formulation of Optimization Problems

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Outline

- Conventions
- Formulation Rules
- Typical Constraint Forms
- Core Models

Conventions I (generally) Use

- INDICES
  - letters like i, j, k, (middle of alphabet)
  - used as subscripts
  - index items in sets with corresponding upper case letter (e.g, I, J, K)
- INPUTS
  - lower case letters near beginning of alphabet
- DECISION VARIABLES
  - upper case letters near end of alphabet

Indices, Inputs and Dec. Var

- INDICES
  - used for enumerating items (e.g., demand nodes, candidate sites, scenarios, time periods)
- INPUTS
  - you **know** these before you start the problem or can readily compute them from other inputs
  - demand values, distances, costs, probabilities, coverage distance, indicators of whether nodes are covered by others, number of sites to locate

Indices, Inputs, and Dec. Var.

- DECISION VARIABLES
  - these are what you want to know or what you must determine **within** the model along the way to determining what you <u>really</u> want to know
  - locations of facilities, whether a node is covered, assignment of demand nodes to facilities, maximum distance

Objective function

- OBJECTIVE FUNCTION
  - this is what you want to minimize or maximize
  - may be a single decision variable (e.g., W for maximum distance between a node and the facility serving it as in P-center)
  - more often will be a function of decision variables (e.g., the total number of facilities as in set covering, the total demand weighted distance as in P-median)

Formulation rules

- Daskin's 10 (or 11) rules of formulation

Rule 1

- Clearly define all subscripts (at least in your own mind) and sets. For example:
  - I: set of demand nodes indexed by *i*
  - J: set of candidate sites indexed by *j*
  - K: set of scenarios indexed by *k*
  - T set of time periods indexed by *t*

Rule 2

- Clearly separate the definitions of
  - indices and sets
  - inputs (or parameters)
  - decision variables

Rule 3

- In defining inputs or decision variables in words, if an index appears in the input or decision variable it should appear in the verbal definition as well

d_{ij} = distance between demand node i and candidate site j

This one (above) is fine

d_{ij} = distance

This one (above) is BAD

Rule 4

- Do not leave dangling subscripts in the objective function

minimize $\sum_{i \in I} \sum_{j \in J} h_i d_{ij} Y_{ij}$ Fine

minimize $\sum_{i \in I} c_{ij} X_{ij}$ **BAD; j index is dangling**

Rule 5

- At least some decision variable must appear in the objective function and in each constraint

minimize $\sum_{i \in I} \sum_{j \in J} h_i d_{ij} Y_{ij}$ Fine

subject to $d_{ij} \geq 0$ **BAD if d_{ij} is an input distance. No decision variable here**

Rule 6

- Be sure all variables are linked in some way to each other (otherwise the problem is separable and you probably have an error)

Rule 6 example

maximize $\sum_{i \in I} h_i Z_i$

subject to $\sum_{j \in J} X_j = P$

$Z_i \in \{0,1\} \quad \forall i \in I$

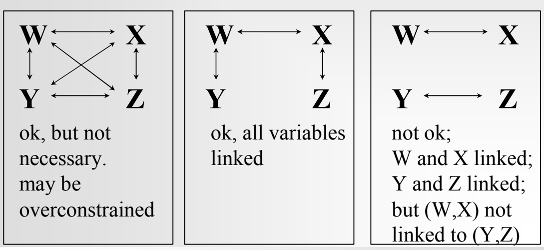
$X_j \in \{0,1\} \quad \forall j \in J$

X and Z variables are unlinked. You need an additional constraint. e.g.,

$$Z_i - \sum_j a_{ij} X_j \le 0 \qquad \forall i \in I$$

More on Rule 6

- Each variable does not have to be directly linked to each other variable



Rule 7

- If a variable or constant used in a constraint includes some index, then either
  - you should be summing over the index OR
  - you should specify the values of the index to which the constraint applies
  - **DO NOT DO BOTH in the same constraint**

Rule 7 examples

$\sum_{j \in J} Y_{ij} = 1$	$\forall i \in I$	**Ok**
$\sum_{j \in J} h_{ik} d_{ijk} Y_{ijk} \le D$	$\forall i \in I$	**BAD; Need to specify what is going on with index *k***
$\sum_{j \in J} Y_{ij} = 1$	$\forall j \in J$	**BAD; Summing over *j* and specifying constraint applies to all *j*; Also, what is going on with index *i*?**

Rule 8

- Try to keep it linear (**IF POSSIBLE**)
  - avoid multiplying decision variables in the objective function or in constraints
  - avoid raising a decision variable to some power
  - avoid logs, trig functions, ...
  - be creative in transformations

Rule 9

- Avoid big M type constraints (**IF POSSIBLE**)
  - constraints with a big value of some constant multiplied by a binary variable
  - often used to turn on or off a constraint depending on the value of the variable
  - may be unavoidable (e.g., α-reliable minimax regret formulation)

Rule 10

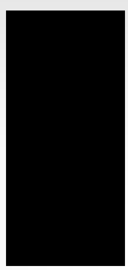
- Disaggregate constraints when possible

$Y_{ij} \le X_j \quad \forall i \in I, \forall j \in J$ **Good, disaggregate constraint**

$\sum_{i \in I} Y_{ij} \le |I| X_j \quad \forall j \in J$ **Not so good, aggregate constraint. Will lead to weaker LP relaxations**

Rule 11

- Know which of rules 1-10 can be bent and when and how to do so



Typical Constraint Forms

- TOTAL CONSTRAINT

$\sum_{j \in J} X_j = p$

- the total of all the X_j variables must be *p*
- e.g., Pick *p* of the X_j variables and set them to 1, set all others to 0 (for X_j a binary variable)

Typical Constraint Forms

- SELECTION or ASSIGNMENT CONSTRAINT

$\sum_{j \in J} Y_{ij} = 1 \quad \forall i \in I$

- For each row *i* (e.g., each demand node), the total of the Y_{ij} variables (for that *i*) must be 1
- each node *i* must be assigned to exactly one facility node

Typical Constraint Forms

- SELECTION or ASSIGNMENT CONSTRAINT

$\sum_{k \in K_j} X_{jk} \le 1 \quad \forall j \in J$

- pick at most one capacity for each site *j* (where K_j is a set of available capacities at candidate site *j*)

Typical Constraint Forms

- SUPPLY Constraints

$$\sum_{j \in J} X_{ij} \le S_i \quad \forall i \in I$$

where

X_{ij} = flow from i to j

- The total flow out of node *i* must be less than or equal to the supply at node *i* (S_i)
- **Note the definition of X_{ij} and that *i* is being used as a supply node and *j* is being used as a demand node**

Typical Constraint Forms

- DEMAND Constraints

$$\sum_{i \in I} X_{ij} \ge D_j \quad \forall j \in J$$

where

X_{ij} = flow from i to j

- The total flow into node *j* must be greater than or equal to the demand at node *j* (D_j)
- **Note the definition of X_{ij} and that *i* is being used as a supply node and *j* is being used as a demand node**

Typical Constraint Forms

- DEMAND-LIKE Constraints

$$\sum_{k \in K} q_k Z_k \ge \alpha$$

- Total probability of selected scenarios must be at least α where K is a set of scenarios
- Used in α-reliable minimax regret model

Typical Constraint Forms

- LINKAGE or FORCING CONSTRAINTS

$$Y_{ij} \le X_j \quad \forall i \in I, \forall j \in J$$

- X_j must be at least as large as Y_{ij} OR
- Y_{ij} must be no bigger than X_j for each pair of *i* and *j*
- You cannot assign demands at *i* to a facility at *j* (Y_{ij}=1) unless you locate at *j* (X_j=1)

Typical Constraint Forms

- LINKAGE or FORCING CONSTRAINTS

$$Z_i - \sum_{j \in J} a_{ij} X_j \le 0 \quad \forall i \in I$$

- Node *i* cannot be counted as being covered (Z_i=1) unless there is at least one facility that is located that is capable of covering node *i* ($\sum_{j \in J} a_{ij} X_j \ge 1$)

Typical Constraint Forms

- LARGEST OF Constraints

$$W \ge \sum_{j \in J} d_{ij} Y_{ij} \quad \forall i \in I$$

- W must be larger than the largest value of $\sum_{j \in J} d_{ij} Y_{ij}$
- Typically $\sum_{j \in J} d_{ij} Y_{ij}$ would represent the distance between node *i* and the facility to which it is assigned
- Used in P-center problems in which we minimize *W* subject to this and other constraints

Typical Constraint Forms

- CONSTRAINTS THAT SWITCH ON and OFF

$$R \ge V_k - \hat{V}_k - M(1 - Z_k) \qquad \forall k \in K$$

where

M = a very large number

so

if $Z_k = 1$ then $R \ge V_k - \hat{V}_k$

but if $Z_k = 0$ then $R \ge V_k - \hat{V}_k - M$

and constraint is "inactive"

Typical Constraint Forms

- CONSTRAINTS THAT SWITCH ON and OFF
  - If Z_k=1 then constraint is active, otherwise it is "inactive"
  - Used in α-reliable minimax regret model
  - Note that in this case, the remainder of the constraint (without the term in *M*) is a LARGEST OF constraint
  - Try to avoid big-M constraints (see rule 9)

Typical Constraint Forms

- DISTANCE Definition

minimize $\sum_{i \in I} \sum_{j \in J} h_{ij} (x_{ij}^{+} + x_{ij}^{-})$

subject to $x_j - x_i = x_{ij}^{+} - x_{ij}^{-} \qquad \forall i \in I, \forall j \in J$

$x_{ij}^{+}, x_{ij}^{-} \ge 0 \qquad \forall i \in I, \forall j \in J$

other constraints

where

$x_j, x_i =$ x coordinates of points j and i

$x_{ij}^{+}, x_{ij}^{-} =$ positive and negative components of the distance between i and j

Typical Constraint Forms

- DISTANCE Definition
  - Note definitions of decision variables (different from normal)
  - Used in some layout formulations
  - Note interaction between objective function and constraints is critical
  - example:

$$x_j = 3; x_i = 7$$
$$x_{ij}^{+} = 0; x_{ij}^{-} = 4$$

Set Covering Model

minimize $\sum_{j \in J} X_j$

NUMBER SELECTED

subject to $\sum_{j \in J} a_{ij} X_j \ge 1 \qquad \forall i \in I$

DEMAND - LIKE constraint

$X_j \in \{0,1\} \qquad \forall j \in J$

INTEGRALITY

Maximal Covering Model

maximize $\sum_{i \in I} h_i Y_i$

Number Covered

subject to $\sum_{j \in J} a_{ij} X_j \ge Y_i \qquad \forall i \in I$

Coverage Constraint (linkage)

$\sum_{j \in J} X_j = p$

Number to Locate

$X_j \in \{0,1\} \qquad \forall j \in J$

$Y_i \in \{0,1\} \qquad \forall i \in I$

Integrality

P-median Model

minimize $\sum_{i \in I} \sum_{j \in J} h_i d_{ij} Y_{ij}$ Demand Wtd Total Dist

subject to

$\sum_{j \in J} Y_{ij} = 1 \quad \forall i \in I$ ASSIGNMENT constraint

$\sum_{j \in J} X_j = p$ TOTAL constraint

$Y_{ij} - X_j \leq 0 \quad \forall i \in I, \forall j \in J$ LINKAGE constraint

$X_j \in \{0,1\} \quad \forall j \in J$

$Y_{ij} \in \{0,1\} \quad \forall i \in I, \forall j \in J$

INTEGRALITY

Fixed Charge Loc. Model

minimize $\sum_{j \in J} f_j X_j + \beta \sum_{i \in I} \sum_{j \in J} h_i d_{ij} Y_{ij}$ Fixed + Transport Cost

subject to

$\sum_{j \in J} Y_{ij} = 1 \quad \forall i \in I$ ASSIGNMENT constraint

$\sum_{j \in J} X_j = p$ TOTAL constraint

$Y_{ij} - X_j \leq 0 \quad \forall i \in I, \forall j \in J$ LINKAGE constraint

$X_j \in \{0,1\} \quad \forall j \in J$

$Y_{ij} \in \{0,1\} \quad \forall i \in I, \forall j \in J$ INTEGRALITY

P-center Model

minimize W Maximum Distance

subject to

$\sum_{j \in J} Y_{ij} = 1 \quad \forall i \in I$ ASSIGNMENT constraint

$\sum_{j \in J} X_j = p$ TOTAL constraint

$Y_{ij} - X_j \leq 0 \quad \forall i \in I, \forall j \in J$ LINKAGE constraint

$W \geq \sum_{j \in J} d_{ij} Y_{ij} \quad \forall i \in I$ MAXIMUM constraint

$X_j \in \{0,1\} \quad \forall j \in J$

$Y_{ij} \in \{0,1\} \quad \forall i \in I, \forall j \in J$ INTEGRALITY