

Chance-Constrained LP

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A "chance constraint" is a modification of a constraint in which the right-hand-side is *random*.

Rather than guaranteeing that the constraint is satisfied for every possible right-hand-side value (which may be impossible, if the random variable is unbounded), a restriction is imposed that the constraint be satisfied by the optimal solution with *at least* a certain specified probability.

Consider the constraint

$$\sum_{j=1}^n a_{ij} x_j \leq b_i \quad \text{where } b_i \text{ is a random variable.}$$

For example, suppose x_j is the production time for process j , and a_{ij} is the consumption rate of raw material i by process j . The right-hand-side b_i could be the (random) quantity of resource i which will be available.

The above constraint requires that the scheduled production time by the processes not consume more raw material than will be available.

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CHANCE CONSTRAINT

$$P \left\{ \sum_{j=1}^n a_{ij} x_j \leq b_i \right\} \geq \alpha$$

i.e., we require that the original constraint

$$\sum_{j=1}^n a_{ij} x_j \leq b_i$$

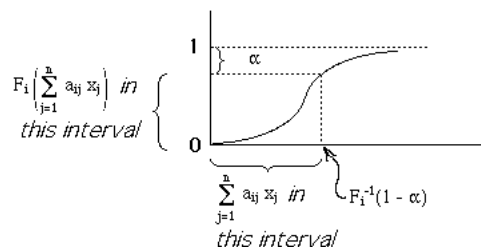
be satisfied with at least probability α .

As stated, this is not a valid LP constraint!

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But
$$F_i \left(\sum_{j=1}^n a_{ij} x_j \right) \leq 1 - \alpha \iff F_i^{-1}(1 - \alpha) \geq \sum_{j=1}^n a_{ij} x_j$$

The inequality on the right is linear!



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LINEARIZING A CHANCE CONSTRAINT

Given the distribution function (cdf)

$$F_i(y) = P\{b_i \leq y\}$$

our chance constraint is equivalent to

$$P \left\{ \sum_{j=1}^n a_{ij} x_j \leq b_i \right\} = 1 - P \left\{ b_i \leq \sum_{j=1}^n a_{ij} x_j \right\} = 1 - F_i \left(\sum_{j=1}^n a_{ij} x_j \right)$$

i.e.,
$$1 - F_i \left(\sum_{j=1}^n a_{ij} x_j \right) \geq \alpha \quad \text{or} \quad F_i \left(\sum_{j=1}^n a_{ij} x_j \right) \leq 1 - \alpha$$

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EXAMPLE

Water Resources Planning Under Uncertainty

A water system manager must allocate water from a stream to three users:

- municipality
- industrial concern
- agricultural sector

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Use	Request	Net Benefit per unit
1. Municipality	2	100
2. Industrial	3	50
3. Agricultural	5	30

Let X_i = amount of water allocated to use # i
 The optimal allocation might be found by solving the LP:
But the decision must be made before the quantity Q of the available water is known!

$$\begin{aligned} \text{Max } & 100X_1 + 50X_2 + 30X_3 \\ \text{subject to } & X_1 + X_2 + X_3 \leq Q \\ & 0 \leq X_1 \leq 2 \\ & 0 \leq X_2 \leq 3 \\ & 0 \leq X_3 \leq 5 \end{aligned}$$

$$\begin{aligned} \text{Max } & 100X_1 + 50X_2 + 30X_3 \\ \text{subject to } & X_1 + X_2 + X_3 \leq Q \\ & 0 \leq X_1 \leq 2 \\ & 0 \leq X_2 \leq 3 \\ & 0 \leq X_3 \leq 5 \end{aligned}$$

Random variable with known probability distribution, namely, $N(7, 1.5)$ i.e., normal, with mean $\mu=7$ and std deviation $\sigma=1.5$.

How should the water be allocated before the quantity available is known?

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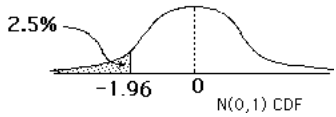
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$$X_1 + X_2 + X_3 \leq Q$$

$$\begin{aligned} P\{Q \geq X_1 + X_2 + X_3\} & \geq \alpha \\ \Leftrightarrow 1 - F(X_1 + X_2 + X_3) & \geq \alpha \\ \Leftrightarrow F(X_1 + X_2 + X_3) & \leq 1 - \alpha \\ \Leftrightarrow X_1 + X_2 + X_3 & \leq F^{-1}(1 - \alpha) \end{aligned}$$

Suppose $\alpha = 97.5\%$
 $\mu = 7$
 $\sigma = 1.5$

$$X_1 + X_2 + X_3 \leq \mu - 1.96\sigma = 4.06$$



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MAX      150 X1 + 50 X2 + 30 X3
SUBJECT TO
          2)  X1 + X2 + X3 <= 4.06
END
SUB      X1          2.00000
SUB      X2          3.00000
SUB      X3          5.00000
    
```

OBJECTIVE FUNCTION VALUE
 1) 403.000000

VARIABLE	VALUE	REDUCED COST
X1	2.0000	-100.0000
X2	2.0600	.0000
X3	.0000	20.0000

ROW SLACK OR SURPLUS DUAL PRICES
 2) .0000 50.0000

LINDO

JOINT CHANCE CONSTRAINTS

Suppose that the RHSs of several constraints are random:

$$\sum_{j=1}^n a_{ij} x_j \leq b_i \quad \text{for } i=1, 2, \dots, k$$

We might impose a chance constraint for *each* of the k random right-hand-sides

$$\sum_{j=1}^n a_{ij} x_j \leq F_i^{-1}(1 - \alpha) \quad \text{for } i=1, 2, \dots, k$$

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Assume that the k random variables are independent, and that we require

$$P\left\{\sum_{j=1}^n a_{1j} x_j \leq b_1 \text{ and } \sum_{j=1}^n a_{2j} x_j \leq b_2 \text{ and } \dots \text{ and } \sum_{j=1}^n a_{kj} x_j \leq b_k\right\} \geq \alpha$$



$$P\left[\sum_{j=1}^n a_{1j} x_j \leq b_1\right] \times P\left[\sum_{j=1}^n a_{2j} x_j \leq b_2\right] \times \dots \times P\left[\sum_{j=1}^n a_{kj} x_j \leq b_k\right] \geq \alpha$$



$$\left[1 - F_1\left(\sum_{j=1}^n a_{1j} x_j\right)\right] \times \left[1 - F_2\left(\sum_{j=1}^n a_{2j} x_j\right)\right] \times \dots \times \left[1 - F_k\left(\sum_{j=1}^n a_{kj} x_j\right)\right] \geq \alpha$$

These chance constraints will *not* guarantee that the optimal solution is feasible with probability α .

Rather, if the right-hand-sides are independent random variables, then the optimal x would satisfy *all* of the constraints with probability α^k .

For example, if $\alpha = 95\%$ and there are $k=10$ chance constraints, then x is feasible with probability $\alpha^k = 59.9\%$

For example, if b_i has an exponential distribution with mean $1/\lambda_i$, i.e.,

$$F_i(y) = 1 - e^{-\lambda_i y}$$

then the joint chance-constraint has the form

$$\left[\exp\left(-\lambda_1 \sum_{j=1}^n a_{1j} x_j\right)\right] \times \left[\exp\left(-\lambda_2 \sum_{j=1}^n a_{2j} x_j\right)\right] \times \dots \times \left[\exp\left(-\lambda_k \sum_{j=1}^n a_{kj} x_j\right)\right] \geq \alpha$$

which is a highly *nonlinear* constraint.

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$$\left[\exp\left(-\lambda_1 \sum_{j=1}^n a_{1j}x_j\right) \right] \times \left[\exp\left(-\lambda_2 \sum_{j=1}^n a_{2j}x_j\right) \right] \times \dots \times \left[\exp\left(-\lambda_k \sum_{j=1}^n a_{kj}x_j\right) \right] \geq \alpha$$

By using a log transformation, we can simplify to

$$\ln \left[\exp\left(-\lambda_1 \sum_{j=1}^n a_{1j}x_j\right) \right] + \dots + \ln \left[\exp\left(-\lambda_k \sum_{j=1}^n a_{kj}x_j\right) \right] \geq \ln \alpha$$

or

$$\left(-\lambda_1 \sum_{j=1}^n a_{1j}x_j\right) + \left(-\lambda_2 \sum_{j=1}^n a_{2j}x_j\right) + \dots + \left(-\lambda_k \sum_{j=1}^n a_{kj}x_j\right) \geq \ln \alpha$$

$$\Rightarrow \sum_{j=1}^n \sum_{i=1}^k (-a_{ij}\lambda_i)x_j \geq \ln \alpha$$

which is, in fact linear!

In cases other than the exponential distribution, however, the constraint *cannot* be linearized by a log transformation.

In the case of the normal distribution, the constraint will remain nonlinear, and cannot even be written in closed form!

Frequently, however, the nonlinear constraint will have a *convex* feasible region, e.g. when b_i 's have normal, gamma, or uniform distributions, so that multiple local optima don't exist.