

Arithmetic-Geometric Mean Inequality

 $\frac{1}{2} \mathbf{a} + \frac{1}{2} \mathbf{b} \geq \mathbf{a}^{\frac{1}{2}} \mathbf{b}^{\frac{1}{2}}$

For example, let a=2 & b=8. Then this inequality is

 $5 = \frac{1}{2} \times 2 + \frac{1}{2} \times 8 \ge \sqrt{2 \times 8} = 4$ Arithmetic mean Geometric MeanIf a=4 & b=9, $6.5 = \frac{1}{2} \times 4 + \frac{1}{2} \times 9 \ge \sqrt{4 \times 9} = 6$ Arithmetic mean Geometric Mean

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Simplest case: Given two positive numbers a & b, their arithmetic mean $\frac{1}{2}a + \frac{1}{2}b$ is greater than or equal to their geometric mean \sqrt{ab}

The Arithmetic-Geometric Mean Inequality

i.e.,
$$\frac{1}{2} a + \frac{1}{2} b \ge a^{\frac{1}{2}} b^{\frac{1}{2}}$$

with equality if & only if a = b

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The Arithmetic-Geometric Mean Inequality
The General Case: Let
$$x_1, x_2, ..., x_n > 0$$

and $\delta_1, \delta_2, ..., \delta_n \ge 0$ and $\sum_{i=1}^n \delta_i = 1$
Then $\sum_{i=1}^n \delta_i x_i \ge \prod_{i=1}^n x_i^{\delta_i}$

with equality if & only if $X_1 = X_2 = \dots = X_n$

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The Arithmetic-Geometric Mean Inequality

$$\sum_{i=1}^{n} \delta_i x_i \ge \prod_{i=1}^{n} x_i^{\delta_i}$$

If we let n=2, and δ_i = $\frac{1}{2}$, then we obtain the earlier inequality,

$$\frac{1}{2}a + \frac{1}{2}b \ge a^{\frac{1}{2}}b^{\frac{1}{2}}$$

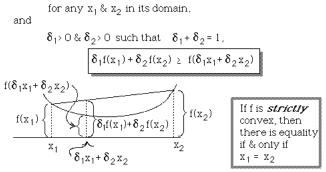
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For the proof of the Arithmetic–Geometric Mean Inequality, we need the fact that $f(x) = -\ln x$ is a strictly convex

function of x in its domain, namely x such that x > 0:

$$\begin{split} f(x) &= -\ln x \\ \implies f'(x) &= -x^{-1} \\ \implies f''(x) &= x^{-2} \\ \implies f''(x) > 0 \text{ for } x > 0 \\ \implies f \text{ is strictly convex} \end{split}$$

A convex function f has the property that



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More generally, for a convex function f,

 $\text{if} \ x_1, x_2, \ldots x_n \ \text{ are in its domain,} \\$

and we are given a "weight" δ_i for each x_i such that

$$\sum_{i=1}^n \, \delta_i = 1$$
 , $\delta_i \ge 0$

then

$$\underbrace{\delta_{1}f(x_{1}) + \delta_{2}f(x_{2}) + ... + \delta_{n}f(x_{n})}_{on \, vex \, combination \\ of \, function \, values} \xrightarrow{f(\delta_{1}x_{1} + \delta_{2}x + ... + \delta_{n}x_{n})}_{function \, vex \, combination \, of \\ the \, x_{1}s}$$

If f is strictly convex, then there is equality above if & only if $x_1 = x_2 = \dots = x_n$

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$$\sum_{i=1}^{n} \delta_{i} x_{i} \ge \prod_{i=1}^{n} x_{i}^{\delta_{i}}$$

Find the dimensions of the open rectangular box Example: with a fixed surface area So having the greatest volume.

Let the dimensions be denoted by $x_1, x_2, \& x_3$ Then the volume is $V(x) = x_1 x_2 x_3$, and the surface area is

So = x_1x_2 + $2x_2x_3$ + $2x_1x_3$ of area of area of front & two endsarea of bottom back

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in

From x > 0, the function
$$f(x) = -\ln x$$
 is strictly convex.
Consequently, if $x_1, x_2, ..., x_n > 0$ and $\sum_{i=1}^n \delta_i = 1$, $\delta_i \ge 0$
then $\delta_1 f(x_1) + \delta_2 f(x_2) + ... + \delta_n f(x_n) \ge f(\delta_1 x_1 + \delta_2 x + ... + \delta_n x_n)$
i.e., $-\delta_1 \ln x_1 - \delta_2 \ln x_2 - ... - \delta_n \ln x_n \ge -\ln(\delta_1 x_1 + \delta_2 x + ... + \delta_n x_n)$
 $\delta_1 \ln x_1 + \delta_2 \ln x_2 + ... + \delta_n \ln x_n \le -\ln(\delta_1 x_1 + \delta_2 x + ... + \delta_n x_n)$
 $\ln \prod_{i=1}^n x_i \delta_i \le -\ln(\sum_{i=1}^n \delta_i x_i)$
Since the log function is strictly increasing, $\prod_{i=1}^n x_i \delta_i \le \sum_{i=1}^n \delta_i x_i$

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$$S_{0} = x_{1}x_{2} + 2x_{2}x_{3} + 2x_{1}x_{3} = 3\left[\frac{x_{1}x_{2} + 2x_{2}x_{3} + 2x_{1}x_{3}}{3}\right]$$

$$S_{0} = 3\left[\frac{1}{3}(x_{1}x_{2}) + \frac{1}{3}(2x_{2}x_{3}) + \frac{1}{3}(2x_{1}x_{3})\right]$$

$$\stackrel{2}{\longrightarrow} 3\left(x_{1}x_{2}\right)^{\frac{1}{3}}\left(2x_{2}x_{3}\right)^{\frac{1}{3}}\left(2x_{1}x_{3}\right)^{\frac{1}{3}}$$

$$= 3 \cdot 4^{\frac{1}{3}}\left(x_{1}^{2}x_{2}x_{3}^{2}\right)^{\frac{1}{3}} = 3 \cdot 4^{\frac{1}{3}}\left[\nabla(x)\right]^{\frac{2}{3}}$$
inequality

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By the A-G Mean Inequality, then,

$$3 \cdot 4^{\frac{1}{3}} [V(x)]^{\frac{2}{3}} \leq S_0$$
 for all x,

with equality if & only if the three terms

 (x_1x_2) , $(2x_2x_3)$, and $(2x_1x_3)$

are equal .

 $S_0 = x_1 x_2 + 2 x_2 x_3 + 2 x_1 x_3$

$$\implies$$
 $x_1x_2 = 2x_2x_3 = 2x_1x_3 = \frac{1}{3}S_0$

We can solve for ${\tt x}_1$, ${\tt x}_2$, & ${\tt x}_3\,$ by using logarithms

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Let
$$z_1 = \ln x_1$$
 and $K = \ln (\frac{1}{3} S_0)$
Then

$$\begin{cases}
\ln x_1 + \ln x_2 = \ln (\frac{1}{3} S_0) \\
\ln 2 + \ln x_2 + \ln x_3 = \ln (\frac{1}{3} S_0) \Rightarrow \\
\ln 2 + \ln x_1 + \ln x_3 = \ln (\frac{1}{3} S_0) \Rightarrow \\
\begin{pmatrix}
1n 2 + z_2 + z_3 = K \\
\ln 2 + z_1 + z_2 = K \\
2z_2 + z_3 = K - \ln 2 \Rightarrow \\
z_3 = \frac{K}{2} - \ln 2
\end{cases} \begin{cases}
z_1 = \frac{K}{2} \\
z_2 = \frac{K}{2} \\
z_3 = \frac{K}{2} - \ln 2
\end{cases}$$

 $x_1 x_2 = \frac{1}{3} S_0 \implies ($

 $x_1x_2 = 2x_2x_3 = 2x_1x_3 = \frac{1}{3}S_0$

$$\begin{array}{c} x_{1}x_{2} = \frac{1}{3} S_{0} \implies \\ 2x_{2}x_{3} = \frac{1}{3} S_{0} \implies \\ 2x_{1}x_{3} = \frac{1}{3} S_{0} \implies \\ 2x_{1}x_{3} = \frac{1}{3} S_{0} \implies \\ \end{array} \begin{array}{c} \ln x_{1} + \ln x_{2} = \ln \left(\frac{1}{3} S_{0}\right) \\ \ln 2 + \ln x_{2} + \ln x_{3} = \ln \left(\frac{1}{3} S_{0}\right) \\ \ln 2 + \ln x_{1} + \ln x_{3} = \ln \left(\frac{1}{3} S_{0}\right) \end{array}$$

This is a linear system of equations in $(\ln x_i)$,

We can solve for x_1 , x_2 , & x_3 by using logarithms

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$$\ln x_{1} = z_{1} = \frac{K}{2}$$

$$\ln x_{2} = z_{2} = \frac{K}{2}$$

$$\ln x_{3} = z_{3} = \frac{K}{2} - \ln 2$$

$$K = \ln (\frac{1}{3}S_{0})$$

which yields the solution

$$x_1 = x_2 = \sqrt{\frac{1}{3}} S_0$$

 $x_3 = \frac{1}{2} \sqrt{\frac{1}{3}} S_0$

so that the volume of the box is

$$V(x) = x_1 x_2 x_3 = \frac{1}{2} \left(\frac{S_0}{3} \right)^{3/2}$$

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Example

Maximize the volume of a cylindrical can given a total cost $\, {\rm C}_{0} \,$ if

the cost of the top & bottom of the can is $\,C_1\,$ cents/square inch the cost of the side of the can is $\,C_2\,$ cents/square inch

Volume is
$$V(r,h) = \pi r^2 h$$

Cost is $2\pi r^2 C_1 + 2\pi r h C_2 =$

is $2\pi r^2 C_1 + 2\pi rh C_2 = C_0$ cost of top cost of side & cost of side

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"Split" the total cost into

$$C_{0} = 2\pi r^{2}C_{1} + 2\pi rhC_{2} = 4 \left(\frac{\pi r^{2}C_{1}}{2} + \frac{\pi rhC_{2}}{2}\right)$$

$$\frac{\pi r^{2}C_{1}}{2} + \frac{\pi rhC_{2}}{2} \geq (\pi r^{2}C_{1})^{\frac{1}{2}} (\pi rhC_{2})^{\frac{1}{2}} = \pi r^{\frac{3}{2}} h^{\frac{1}{2}} (C_{1}C_{2})^{\frac{1}{2}}$$

$$\frac{h}{hequality}$$
Not a constant multiple of the volume:

$$V(r,h) = \pi r^{2}h$$

Unfortunately, we cannot proceed as before. We need to "split" \mathbb{C}_0 into a sum of terms differently.

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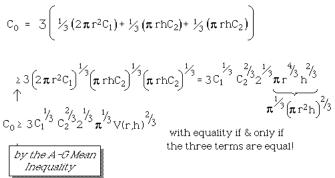
$$C_0 = 2\pi r^2 C_1 + 2\pi rhC_2$$

$$C_0 = 2\pi r^2 C_1 + \pi rhC_2 + \pi rhC_2$$
Spin the of three terms

=
$$\Im \left[\frac{1}{3} (2\pi r^2 C_1) + \frac{1}{3} (\pi rh C_2) + \frac{1}{3} (\pi rh C_2) \right]$$

We now again apply the Arithmetic-Geometric Mean Inequality to the sum within the braces, with weights equal to $\frac{1}{3}$ for each term.

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Geometric Programming

"Geometric Programming" is an optimization technique which determines what fraction of the cost is to be attributed to each term, i.e., determines the "weights" in the Arithmetic-Geometric Mean Inequality.

That is, by the A-G Mean Inequality,

$$C_0 \ge 3C_1^{1/3} C_2^{2/3} 2^{1/3} \pi^{1/3} V(r,h)^{2/3}$$
 for all r & h

with equality if & only if the terms are equal:

$$(2\pi r^2 C_1) = (\pi rhC_2) = (\pi rhC_2) = \frac{1}{3} C_0$$

$$\Rightarrow r = \sqrt{\frac{l_3 C_0}{2\pi C_1}} \Rightarrow h = \frac{l_3 C_0}{\pi r C_2} = \frac{1}{C_2} \sqrt{\frac{2C_0 C_1}{3\pi}}$$

The "trick" is knowing how to "split" the total cost into several terms.

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Example
Minimize
$$f(x) = C_1 x^3 + \frac{C_2}{x}$$

 $f(x) = C_1 x^3 + \frac{1}{3} \frac{C_2}{x} + \frac{1}{3} \frac{C_2}{x} + \frac{1}{3} \frac{C_2}{x}$
 $= 4 \left(\frac{1}{4} \left(C_1 x^3 \right) + \frac{1}{4} \left(\frac{1}{3} \frac{C_2}{x} \right) + \frac{1}{4} \left(\frac{1}{3} \frac{C_2}{x} \right) + \frac{1}{4} \left(\frac{1}{3} \frac{C_2}{x} \right) \right)$
 $\geq 4 \left(\left(C_1 x^3 \right)^{1/4} \left(\frac{1}{3} \frac{C_2}{x} \right)^{1/4} \left(\frac{1}{3} \frac{C_2}{x} \right)^{1/4} \left(\frac{1}{3} \frac{C_2}{x} \right)^{1/4} \left(\frac{1}{3} \frac{C_2}{x} \right)^{1/4} \right) = \underbrace{4 \left(\frac{1}{3} \right)^{3/4} C_1^{-1/4} C_2^{-3/4}}_{doesn't depend on x'}$
by A-G Mean
Inequality

Minimize $f(x) = C_1 x^3 + \frac{C_2}{x} \ge 4 \left(\frac{1}{3}\right)^{3/4} C_1^{-1/4} C_2^{-3/4}$ with equality if & only if $C_2 = \frac{3}{2} - \frac{1}{2} C_2 = \frac{1}{2} C_2 C_2 = \frac{1}{2} \left(\frac{1}{3}\right)^{3/4} C_1^{-1/4} C_2^{-3/4}$

$$C_{1}x^{2} = \frac{1}{3}\frac{C_{2}}{x} = \frac{1}{3}\frac{C_{2}}{x} = \frac{1}{3}\frac{C_{2}}{x} = \frac{1}{3}\frac{C_{2}}{x} = (\frac{1}{3})^{74}C_{1}^{74}C_{$$

Again, the "trick" is knowing how to "split" the objective, i.e., allocating a fourth of the total cost to the first term, and three-fourths to the second term.

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