

## One-Dimensional Search Methods

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In addition to solving nonlinear optimization problems with a single variable, we require an algorithm to do "line searches" as part of a multi-dimensional nonlinear optimization problem:

$$\text{Minimize } f(x^k + t d^k)$$

where  $x^k$  = the  $k^{\text{th}}$  iterate,  $x^k \in \mathbb{R}^n$   
 $d^k$  = (feasible) direction of descent  
 $t$  = step size  
 $x^{k+1} = x^k + t^* d^k$  for optimal stepsize  $t^*$

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- an analytic expression for  $f(x)$  might be unknown...  $f(x)$  might be "evaluated" by performing a laboratory or simulation experiment, for example
- it is assumed that the function  $f$  is *unimodal*, i.e., a local optimum will be globally optimal.
- the result of minimization will be a "sufficiently small" *interval of uncertainty* containing the optimum.
- the derivative of  $f$  need not be computed in many of these methods.

- Three-Point Equi-Interval Search
- Golden-Section Search
- Fibonacci Search
- Polynomial Interpolation

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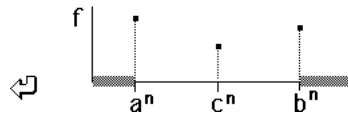
### Three-Point Equi-Interval Search

*Simple, but inefficient....  
not recommended!*

Assume that at the  $n^{\text{th}}$  iteration we have the interval of uncertainty  $[a^n, b^n]$  and its midpoint  $c^n = \frac{a^n + b^n}{2}$ ,

$$c^n = \frac{a^n + b^n}{2}$$

along with the function values  $f(a^n)$ ,  $f(b^n)$ , &  $f(c^n)$

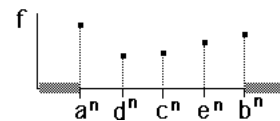


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Find the midpoints of the two subintervals  $[a^n, c^n]$  and  $[c^n, b^n]$ :

$$d^n = \frac{3a^n + b^n}{4}, \quad e^n = \frac{a^n + 3b^n}{4}$$

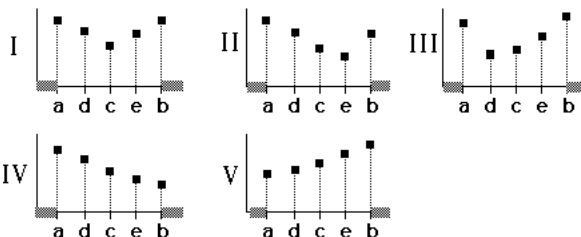
and evaluate  $f(d^n)$  and  $f(e^n)$ :



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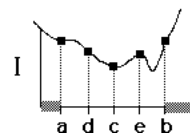
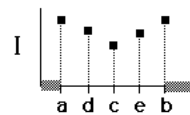
Consider the relative magnitudes of the function at these five points:

There are several cases to consider:



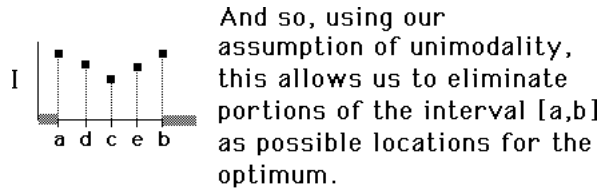
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For example, suppose that  $f(c)$  is lower than  $f(d)$  and  $f(e)$ .

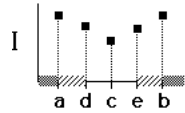


*Assuming that the function is unimodal, the minimum cannot be in the interval  $[e, b]$ !*

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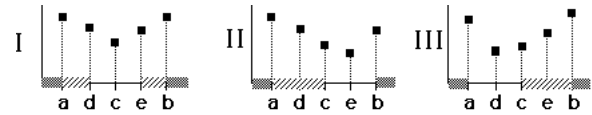


$x^*$  cannot be in  $[a,d]$  or  $[e,b]$ :



We therefore choose  $a^{n+1} = d^n, c^{n+1} = c^n, b^{n+1} = e^n$  to begin iteration  $n+1$ .

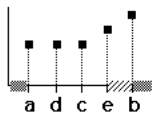
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In cases I, II, and III, 50% of the interval is eliminated.

In cases IV and V, 75% of the interval is eliminated!

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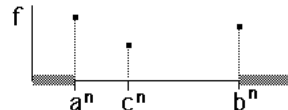


In the event that the smallest of  $f(a), f(b), f(c), f(d),$  &  $f(e)$  is not unique, less than 25% of the interval can be eliminated.

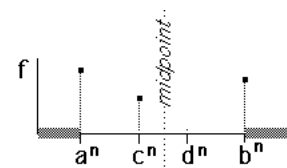
This event will generally be very rare, especially given round-off errors, etc.

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**Golden-Section Search**



As in 3-point equi-interval search, at the  $n$  iteration we have an interval of uncertainty  $[a^n, b^n]$  and an interior point  $c^n \in (a^n, b^n)$ , but  $c^n$  is *not the midpoint*! We insert a *single* additional point  $d^n$  so that  $c^n$  and  $d^n$  are *symmetric* about the midpoint of the interval  $[a^n, b^n]$  and compare the values of  $f(a^n), f(b^n), f(c^n),$  &  $f(d^n)$ .



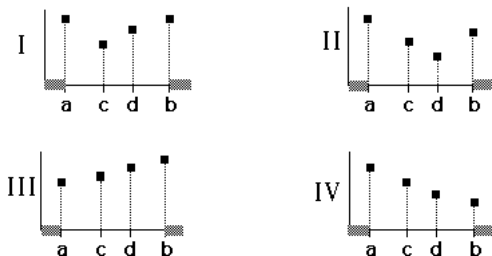
The new point  $d^n$  is selected so as to be symmetric to  $c^n$  in the interval.

$$b^n - d^n = c^n - a^n \Rightarrow d^n = a^n + b^n - c^n$$

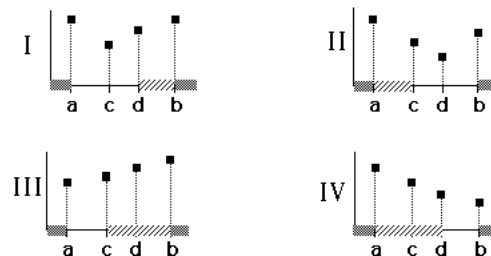
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Assuming unimodality, we can eliminate a segment from the interval of uncertainty



The shaded segments (////) can be eliminated from the interval of uncertainty:



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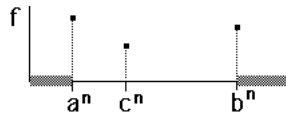
Once the location of  $c^0$  has been determined within the original interval of uncertainty  $[a^0, b^0]$ , the location of subsequent points is determined (by symmetry).

How should  $c^0$  be located within  $[a^0, b^0]$ ?

In "Golden Section" search, this is done so that the ratio  $\frac{c^n - a^n}{b^n - a^n}$  is constant ( $\alpha$ )  $\forall n$  (assuming points are labeled so that  $c^n < d^n$ )

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**Fibonacci Search**

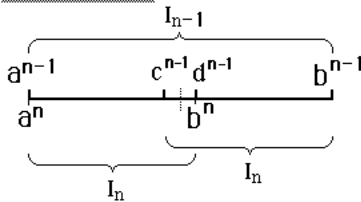


As is the case with Golden Section Search, this method begins each iteration with an interval of uncertainty  $[a, b]$  and one interior point  $c$ , and then inserts another interior point  $d$  which is symmetric to  $c$ .

In Fibonacci search, however, the ratio  $\frac{c^n - a^n}{b^n - a^n}$  is not constant, but converges to  $1/2$ !

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At the last iteration



The distance between  $c^n$  and  $d^n$  will be  $\epsilon$

$I_{n-1} = 2 I_n - \epsilon$

The final interval of uncertainty will be one of these two intervals

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$I_{k-1} = I_k + I_{k+1}$

"Fibonacci Numbers"

$\Rightarrow I_{n-1} = 2 I_n - \epsilon = F_2 I_n - F_0 \epsilon$   
 $I_{n-2} = 3 I_n - \epsilon = F_3 I_n - F_0 \epsilon$   
 $I_{n-3} = 5 I_n - 2 \epsilon = F_4 I_n - F_1 \epsilon$   
 $I_{n-4} = 8 I_n - 3 \epsilon = F_5 I_n - F_2 \epsilon$   
 $I_{n-5} = 13 I_n - 5 \epsilon = F_6 I_n - F_3 \epsilon$   
 $\vdots$   
 $I_{n-k} = F_{k+1} I_n - F_{k-1} \epsilon$   
 $\vdots$   
 $1 = I_1 = F_n I_n - F_{n-2} \epsilon$

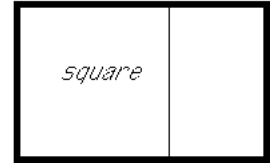
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This requirement uniquely determines

$\alpha = \frac{c^n - a^n}{b^n - a^n} = \frac{3 - \sqrt{5}}{2} \forall n$   
 $= 0.381966$   
 $\beta = 1 - \alpha = 0.618034$

known to early Greek mathematicians as the "Golden Section"

If a rectangle with ratio width:length =  $\beta$  is cut to yield a square, the other rectangle also has width:length =  $\beta$



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Given:  $[a^1, b^1]$  = initial interval of uncertainty

$I_k = b^k - a^k$

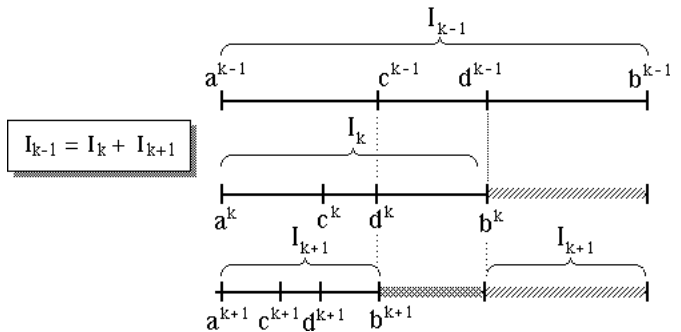
$I_n = b^n - a^n$  = desired length of interval of uncertainty

$\epsilon$  = "distinguishability constant"  $> 0$  (i.e.,  $x$  &  $y$  are indistinguishable if  $|x - y| < \epsilon$ )

For ease of discussion, assume  $I_1 = b^1 - a^1 = 1$

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In general, we have



$I_{k-1} = I_k + I_{k+1}$

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**Fibonacci Numbers**

Leonardo of Pisa, son of Bonacci ("Fibonacci") 1202 AD.

Rule for Generating the Sequence:

$F_0 = F_1 \equiv 1$   
 $F_n = F_{n-1} + F_{n-2}, n \geq 2$

$\Rightarrow F_2 = F_1 + F_0 = 1 + 1 = 2$   
 $F_3 = F_2 + F_1 = 2 + 1 = 3$   
 $F_4 = F_3 + F_2 = 3 + 2 = 5$   
 $F_5 = F_4 + F_3 = 5 + 3 = 8$   
 $\vdots$

n	F <sub>n</sub>
0	1
1	1
2	2
3	3
4	5
5	8
6	13
7	21
8	34
9	55
10	89

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$$I_1 = F_n I_n - F_{n-2} \epsilon$$

Solving for the "reduction ratio"  $\frac{I_1}{I_n}$

$$\frac{I_1}{I_n} = \frac{F_n}{1 + F_{n-2} \epsilon}$$

Given a desired reduction ratio, we can find  $n$ , the required number of iterations.

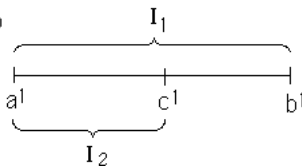
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Once we have determined  $n$  (the # of iterations), we can compute

$$I_2 = F_{n-1} \left[ \frac{1 + F_{n-2} \epsilon}{F_n} \right] - F_{n-3} \epsilon$$

$$= \frac{F_{n-1}}{F_n} + \left[ \frac{F_{n-1} F_{n-2}}{F_n} \right] \epsilon$$

This will tell us where to put our initial interior point  $c^1$  within  $[a^1, b^1]$ .



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**Polynomial Interpolation**

In quadratic & cubic interpolation methods, we use information about the function at two or more points to determine a polynomial in agreement with the known information about the function  $f$ .

A minimum point is then computed for the interpolating polynomial to obtain a new point interior to the interval of uncertainty  $[a, b]$ .



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**Polynomial Interpolation**

- Lagrange's Interpolating Polynomials  
*polynomials  $p(x)$  with  $p(a)=f(a)$ , etc.*
- Quadratic Interpolation
- Hermite Interpolating Polynomials  
*polynomials  $p(x)$  with  $p(a)=f(a)$ ,  $p'(a)=f'(a)$*
- Cubic Interpolation

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**Example**

Suppose we desire  $I_n \leq 0.01 I_1$

i.e.,  $\frac{I_1}{I_n} \geq 100$

and suppose  $\epsilon \approx 0$

Then  $\frac{I_1}{I_n} = \frac{F_n}{1 + F_{n-2} \epsilon} \approx F_n$

Choose  $n$  so that  $F_n \geq 100$

$\Rightarrow n = 11$

$n$	$F_n$
0	1
1	1
2	2
3	3
4	5
5	8
6	13
7	21
8	34
9	55
10	89
11	144
...	...

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**Example**

For example, suppose  $n = 11$  and  $\epsilon \approx 0$

$$I_2 \approx \frac{F_{10}}{F_{11}} \approx 0.6180555$$

Throughout the remainder of the iterations, the other interior points are located to retain the symmetry of  $c^k$  and  $d^k$ .



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Polynomial	Given information
Quadratic	$a, f(a), b, f(b), c(a, b), f(c)$
Cubic	$a, f(a), f'(a), b, f(b), f'(b)$

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**Lagrange's Interpolating Polynomials**

Assume that we are given the  $n+1$  values

$$\{x_0, x_1, x_2, \dots, x_n\}$$

and function values  $f(x_i)$ ,

What is the polynomial  $p(x)$  of degree  $n$  which agrees exactly with  $f(x)$  at the values  $x_i, i=0, 1, 2, \dots, n$  ?



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Define

$$L_j(x) = \prod_{k \neq j} \frac{x - x_k}{x_j - x_k} = \frac{x - x_0}{x_j - x_0} \times \frac{x - x_1}{x_j - x_1} \times \dots \times \frac{x - x_n}{x_j - x_n}$$

Properties:

$L_j(x)$  is a polynomial of degree  $n$   
 $L_j(x_j) = 1$   
 $L_j(x_k) = 0$  for  $k \neq j$

That is, for each  $x_j$  we define a polynomial of degree  $n$  which is 1 at  $x_j$  but zero for  $x_k, k \neq j$ .

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**Hermite Interpolating Polynomials**

Assume that we are given the  $n+1$  triplets of values

$$\begin{matrix} x_0, f(x_0), f'(x_0) \\ x_1, f(x_1), f'(x_1) \\ \vdots \\ x_n, f(x_n), f'(x_n) \end{matrix}$$

We want to find a polynomial  $p(x)$  such that

$$\begin{matrix} p(x_i) = f(x_i) \\ p'(x_i) = f'(x_i) \end{matrix} \quad \forall i=0,1,\dots,n$$



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Differentiating  $p(x) = \sum_{j=0}^n f_j h_j(x) + \sum_{j=0}^n f'_j \bar{h}_j(x)$

yields  $p'(x) = \sum_{j=0}^n f_j h'_j(x) + \sum_{j=0}^n f'_j \bar{h}'_j(x)$

We would therefore like  $h_j$  and  $\bar{h}_j$  to satisfy

$$\begin{matrix} h'_j(x_k) = 0 \quad \forall k \neq j \\ \bar{h}'_j(x_k) = \begin{cases} 1 & \text{if } k = j \\ 0 & \text{if } k \neq j \end{cases} \end{matrix}$$

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**Quadratic Interpolation**

Given:  $a, b, c$   
and  $f(a), f(b), f(c)$

The interpolating quadratic polynomial is

$$p(x) = f(a) \frac{x-b}{a-b} \times \frac{x-c}{a-c} + f(b) \frac{x-a}{b-a} \times \frac{x-c}{b-c} + f(c) \frac{x-a}{c-a} \times \frac{x-b}{c-b}$$

We want to find the minimum of this polynomial,

and so we will find  $x$  such that  $\frac{dp(x)}{dx} = 0$



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Lagrange's interpolating polynomial is

$$p(x) = \sum_{j=0}^n f(x_j) L_j(x)$$

This polynomial agrees exactly with the function  $f$  at  $x_0, x_1, x_2, \dots, x_n$

i.e.,  $p(x_k) = \sum_{j=0}^n f(x_j) L_j(x_k) = f(x_k) \quad \forall k=0,1,2,\dots,n$



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Notation  $f_i \equiv f(x_i)$   
 $f'_i \equiv f'(x_i) = \frac{df}{dx}(x_i)$

The polynomial  $p(x)$  will be of the form

$$p(x) = \sum_{j=0}^n f_j h_j(x) + \sum_{j=0}^n f'_j \bar{h}_j(x)$$

In order that  $p(x_i) = f(x_i)$   
 $h_j$  and  $\bar{h}_j$  will satisfy

$$\begin{matrix} h_j(x_j) = 1 \\ h_j(x_k) = 0 \quad \forall k \neq j \\ \bar{h}_j(x_k) = 0 \quad \forall k \neq j \end{matrix}$$

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The following functions have the desired properties:

$$\begin{matrix} h_j(x) = [1 - 2(x - x_j) L'_j(x_j)] L_j^2(x) \\ \bar{h}_j(x) = (x - x_j) L_j^2(x) \end{matrix}$$



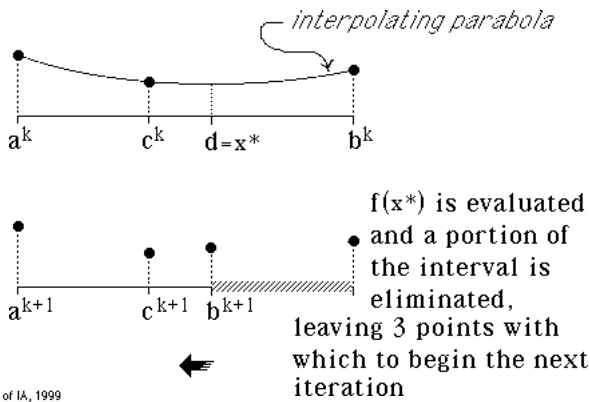
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$$\begin{aligned} \frac{dp(x)}{dx} &= \frac{f(a)}{(a-b)(a-c)}(2x-b-c) + \frac{f(b)}{(b-a)(b-c)}(2x-a-c) \\ &\quad + \frac{f(c)}{(c-a)(c-b)}(2x-a-b) = 0 \end{aligned}$$

$$\Rightarrow x^* = \frac{1}{2} \frac{f(a)(b^2 - c^2) + f(b)(c^2 - a^2) + f(c)(a^2 - b^2)}{f(a)(b-c) + f(b)(c-a) + f(c)(a-b)}$$

Having located this 4<sup>th</sup> point (call it  $d$ ), evaluate  $f(d)$  and proceed as in Golden Section or Fibonacci search, eliminating a portion of the interval  $[a,b]$

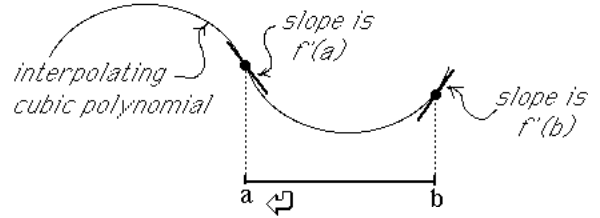
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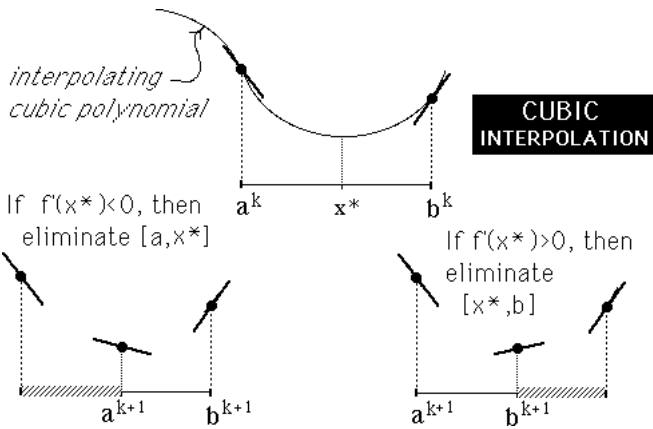
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**Cubic Interpolation**

Given  $a, f(a), f'(a),$  and  $b, f(b), f'(b),$  there is a unique cubic polynomial which passes through the points  $(a, f(a)), (b, f(b))$  and is tangent to the graph at these points.



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**Cubic Interpolation, Using Hermite Polynomials**

Given  $a, f(a), f'(a),$  and  $b, f(b), f'(b)$

The interpolating CUBIC polynomial is

$$p(x) = f(a)h_0(x) + f(b)h_1(x) + f'(a)\bar{h}_0(x) + f'(b)\bar{h}_1(x)$$

where

$$h_0(x) = \left[1 - 2 \frac{x-a}{a-b}\right] \left[\frac{x-b}{a-b}\right]^2 \quad \bar{h}_0(x) = (x-a) \left[\frac{x-b}{a-b}\right]^2$$

$$h_1(x) = \left[1 - 2 \frac{x-b}{b-a}\right] \left[\frac{x-a}{b-a}\right]^2 \quad \bar{h}_1(x) = (x-b) \left[\frac{x-a}{b-a}\right]^2$$

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**Finding the stationary point of  $p(x)$  in  $[a,b]$**

Step 1:  $z = f(a) + f(b) + 3 \left[ \frac{f(a) - f(b)}{b - a} \right]$

Step 2:  $w = \sqrt{\max \{0, z^2 - f'(a)f'(b)\}}$

Step 3:  $x^* = b - \frac{(b-a)(f'(b) + w - z)}{f'(b) + 2w - f'(a)}$

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