

#### Search Trees

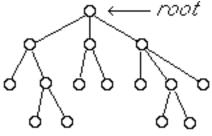
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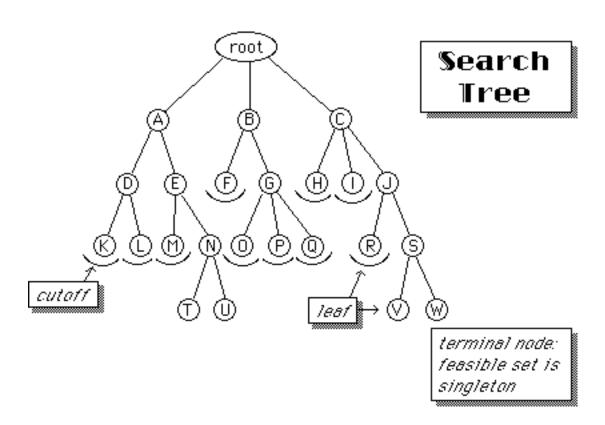




- Each node of the search tree for a problem represents a subset of feasible solutions of the problem
- The root of the tree represents the set of all feasible solutions of the problem
- The descendents of each node of the tree represent a partition of the set represented by that node

A collection of subsets  $B_i$  of set A (i=1,2,...t) is a **partition** if

$$B_1 \cup B_2 \cup B_3 \cdots \cup B_t = A$$
  
and  $B_i \cap B_j = \emptyset$  if  $i \neq j$ 



#### Example: Ranking Nodes in a Preference Graph

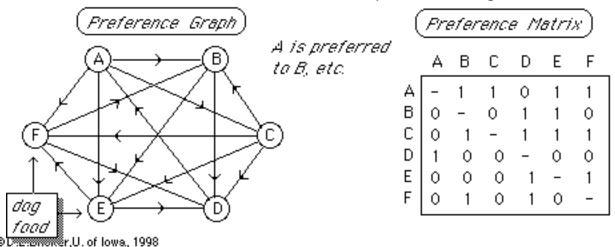
In many experiments (especially in the social sciences, when numerical measurement of attributes are difficult or impossible), one is required to rank a set of objects by comparing only two at a time.

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# Example

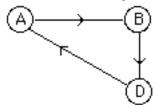
Six different dog foods are to be ranked according to their appeal to dogs.

Each day, 2 of the 6 are served to a dog, who indicates his preference by finishing it first.



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In the dog food example, the dog exhibited some inconsistency: for example,



he preferred A over B, B over D,

and Dover A!

How can we establish a "good" ranking?

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#### Methods for Ranking

- ranking by score: the score of an object is the number of pairs in which it is preferred (i.e., the row-sum of the preference matrix).
  - ties may occur
  - assumes every possible pair was compared

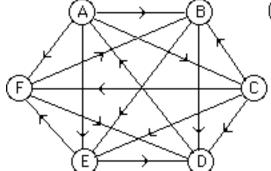
	Α	В	С	D	Ε	F
Α	-	1	1	0	1	1
A B C D E F	0	-	0	1	1	0
С	0	1	-	1	1	1
D	1	0	0	-	0	0
Ε	0	0	0	1	-	1
F	0	1	0	1	0	-

```
4
2
4
1
2
2
```

#### Methods for Ranking

 ranking by Hamiltonian path: find a path through every node of the preference graph such that each node is preferred over its successor.

For example,  $A \rightarrow C \rightarrow B \rightarrow E \rightarrow F \rightarrow D$ or  $A \rightarrow C \rightarrow F \rightarrow F \rightarrow B \rightarrow D$ 



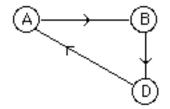
(several such paths may exist!)

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## Methods for Ranking

• ranking with minimum discrepancies

A discrepancy is an instance in which X is ranked above Y, but Y is preferred to X



For example, the ranking A > B > D has one discrepancy (i.e., A>D)

- does not assume that every pair was compared!
- is a difficult problem to solve

## Using a Search Tree for Minimum Discrepancy Ranking

Two different methods for partitioning:

 choose a pair of objects X & Y which have not been ranked.

Form two subsets of rankings:

- --those in which X > Y, i.e., X is ranked above Y
- --those in which Y > X, i.e., Y is ranked above X

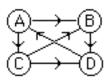
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#### Second method of partitioning:

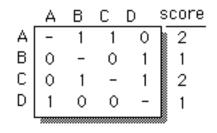
 an object is assigned to a position in the ranking e.g., in the first partition, n nodes are created, in each of which one of the n objects is assigned to the first position in the ranking, and

in the second partition, **n-1** nodes are created, one for each of the remaining **n-1** objects which might be assigned to the **second** position in the ranking, etc.

# Example



First Partitioning Method

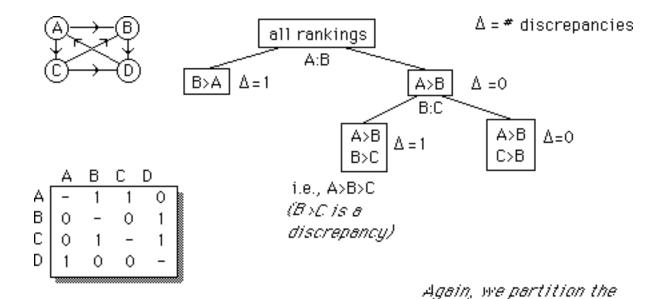


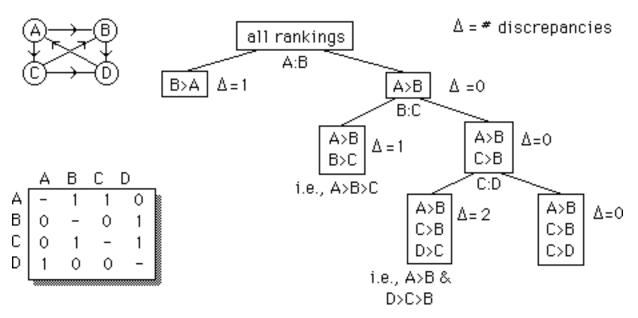
 $\Delta$  = # discrepancies

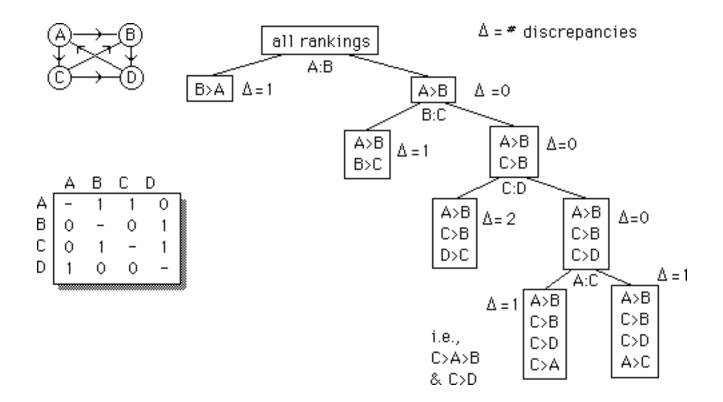


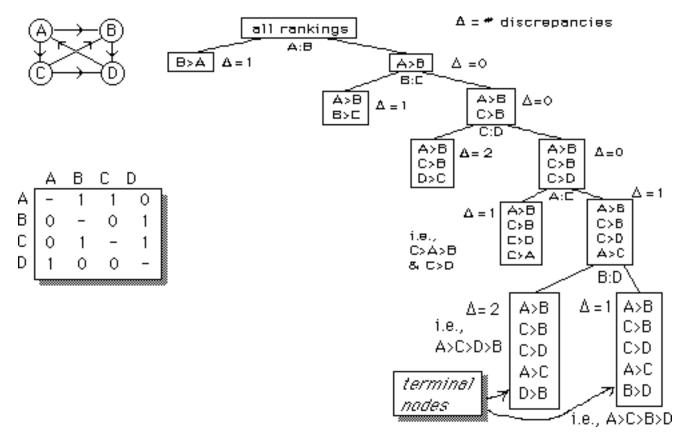
We will partition the most promising node, that with no discrepancies

most promising node

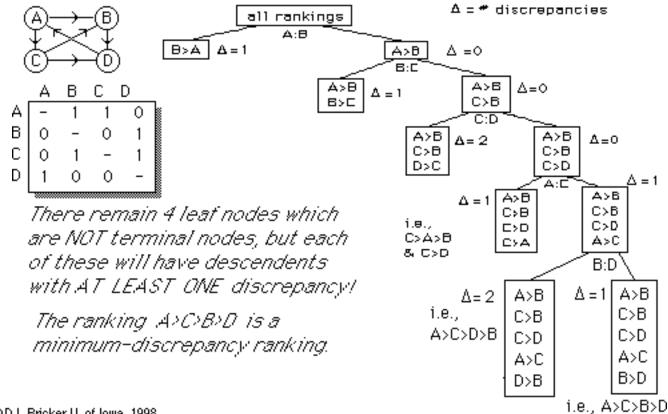




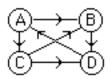




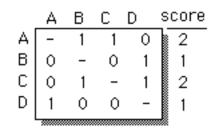
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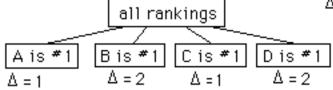
### Example



Second Partitioning Method

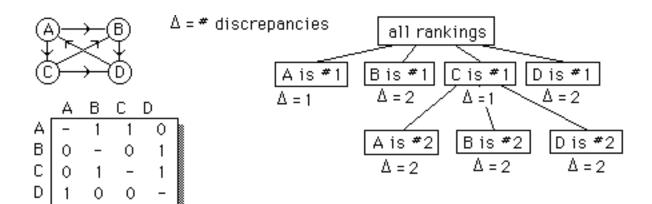


 $\Delta$  = # discrepancies



We will partition the most promising node, that with one discrepancy

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Second Partitioning Method

We will partition the most promising node, that with one discrepancy

