

Center Problems in a Network



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Center of a Network

Define the function $\sigma(x) = \max_{j \in N} d(x, j)$

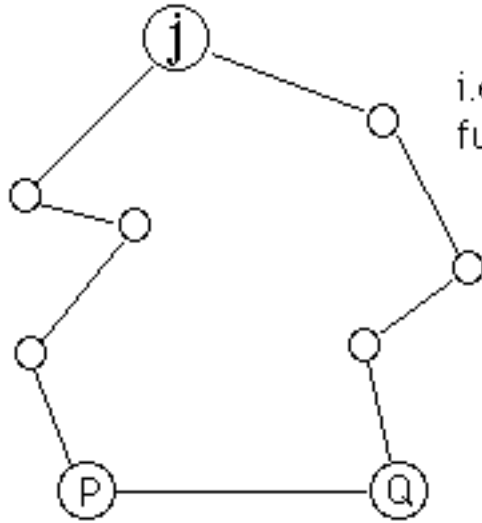
where

$d(x, j)$ = shortest path from x to node j

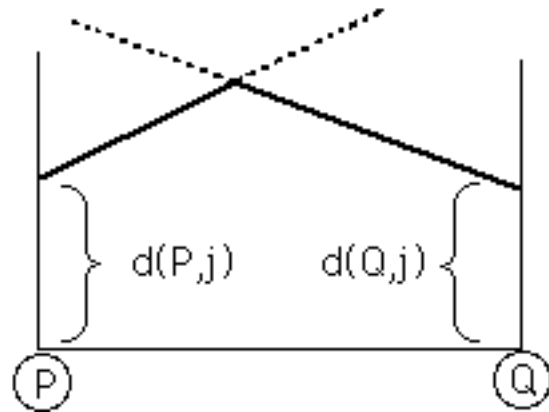
i.e., the distance from x to the farthest node of the network.

Suppose $x \in \text{edge } [P, Q]$

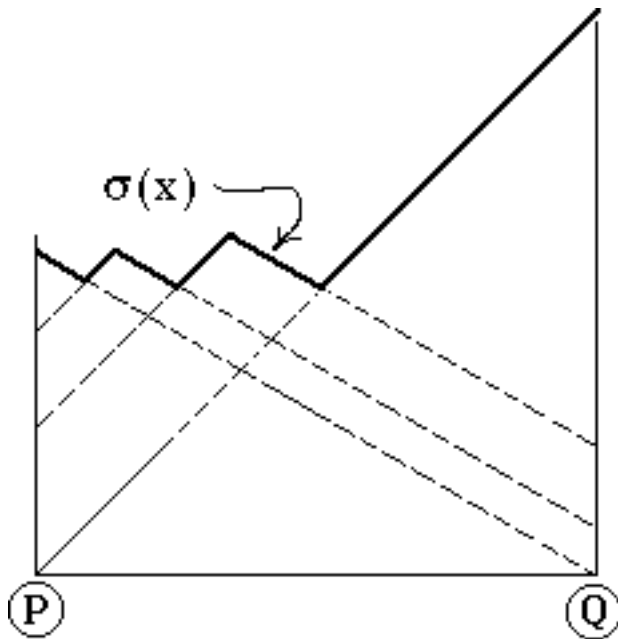
$d(x, j)$ = shortest path from x to node j
 = minimum of $d(x, P) + d(P, j)$
 and $d(x, Q) + d(Q, j)$



i.e., the lower envelope of two linear functions



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= maximum $d(x, j)$
 $j \in N$

For $x \in \text{edge } [P, Q]$,

$\sigma(x)$ is the upper envelope of the functions $d(x, j)$ for $j \in N$

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The *Vertex Center* is the point $x \in N$ which solves

$$\underset{x \in N}{\text{minimize}} \quad \sigma(x)$$

i.e., the point which solves the *minimax* problem

$$\underset{x \in N}{\text{minimize}} \left\{ \underset{j \in N}{\text{maximum}} \, d(x, j) \right\}$$

The *Edge Center* of an edge $[J, K]$ is the point z on edge $[J, K]$ which solves

$$\underset{x \in [J, K]}{\text{minimize}} \quad \sigma(x)$$

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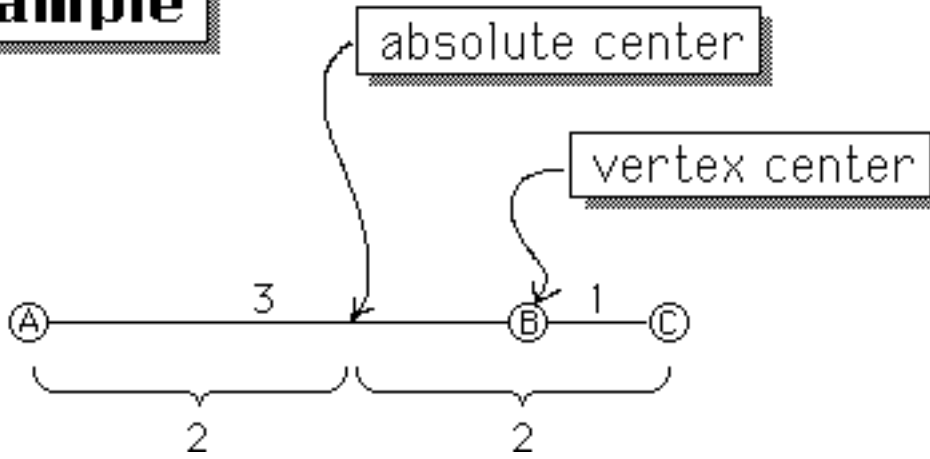
The *Absolute Center* of a network is the point z (a node or a point on an edge) which solves

$$\underset{x \in G}{\text{minimize}} \quad \sigma(x)$$

where $G = N \cup A$ is the set of nodes and points on edges in the edge set A

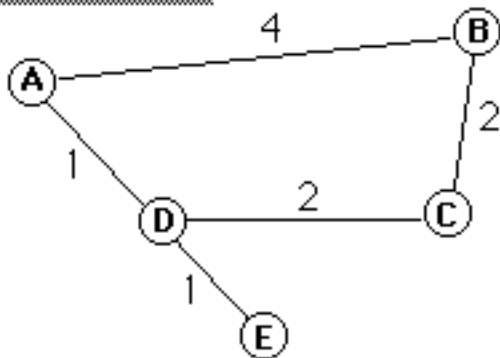
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Example



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Example

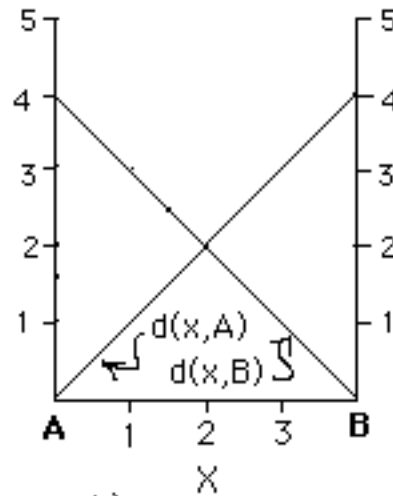
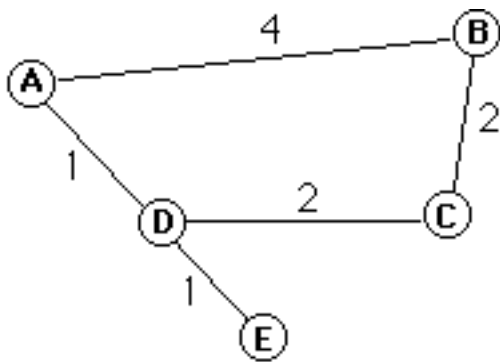


Where should a fire station be located so as to minimize the distance to the farthest village?

$d(x,J)$ = shortest path from point x (on the network) to village J , $J \in N = \{A,B,C,D,E\}$

$$\text{Minimize } \left\{ \text{Max}_{J \in N} d(x,J) \right\}_x$$

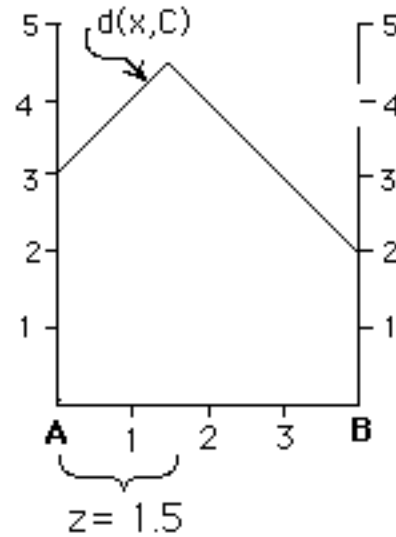
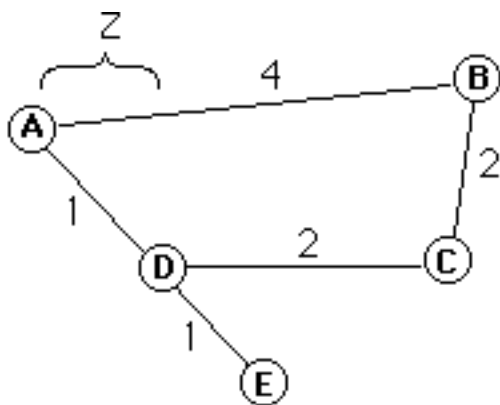
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Consider $d(x,J)$ for points x on the edge (A,B)

$d(x,A)$ is monotonically increasing (slope: +1) as x moves from A to B , while $d(x,B)$ is monotonically decreasing (slope: -1)

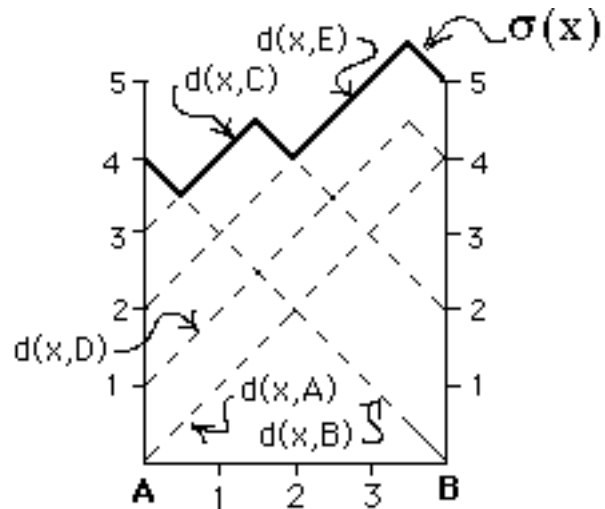
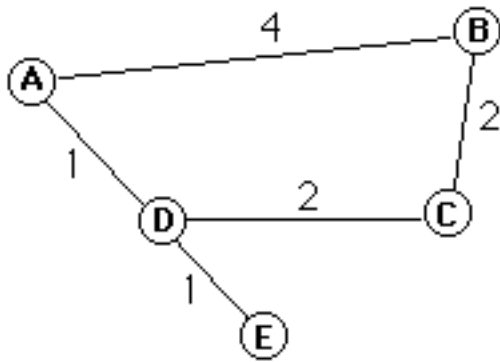
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$d(x,C) = 3$ at $x=A$, and increases (slope= +1) as x moves toward B . At the point x where $d(x,A)+1+2 = d(x,B)+2$, the function begins to decrease (slope= -1).

$$z+1+2=(4-z)+2 \implies z=1.5$$

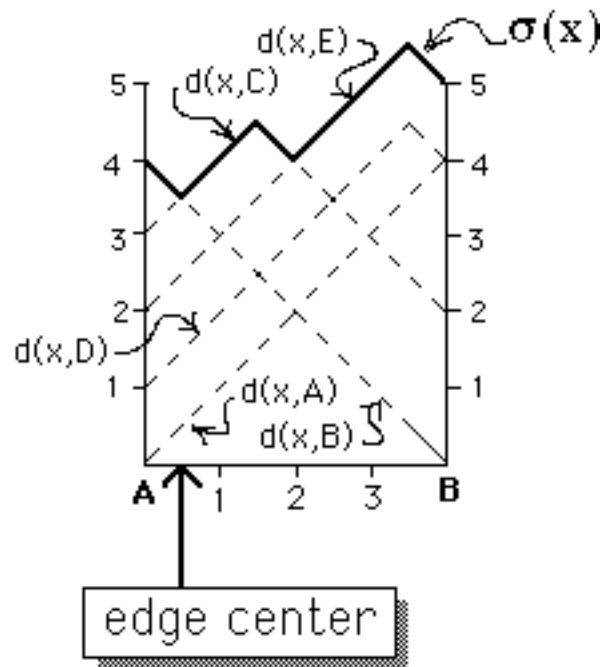
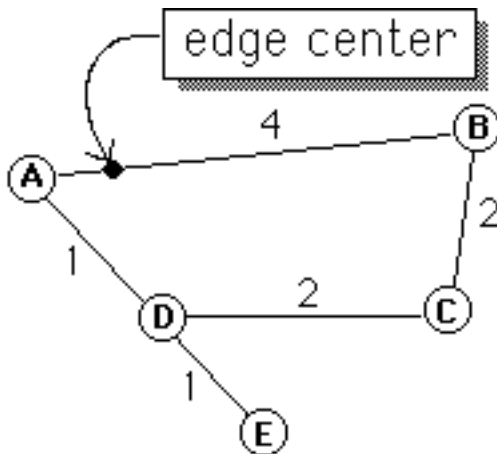
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$$\sigma(x) = \text{maximum}_{j \in N} d(x, j)$$

$\sigma(x)$ is the upper envelope of the family of functions $d(x, J), J \in N$

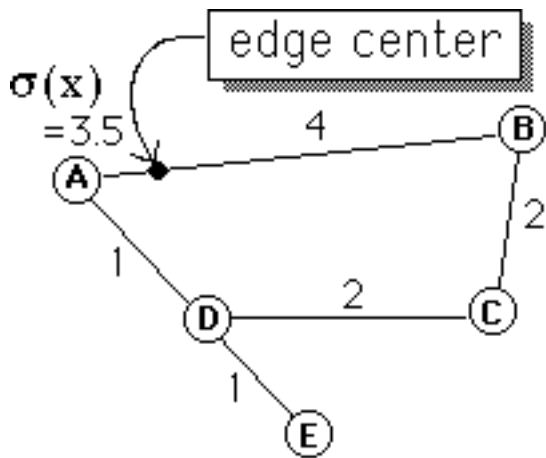
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$$\sigma(x) = \text{maximum}_{j \in N} d(x, j)$$

The point which minimizes the function σ on $[A, B]$ lies 0.5 miles from A

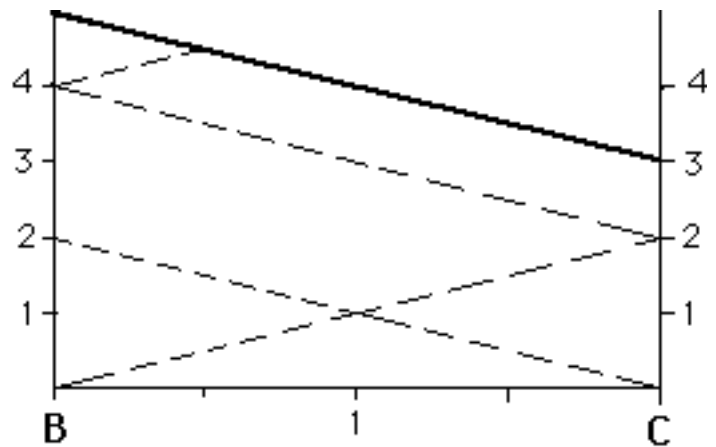
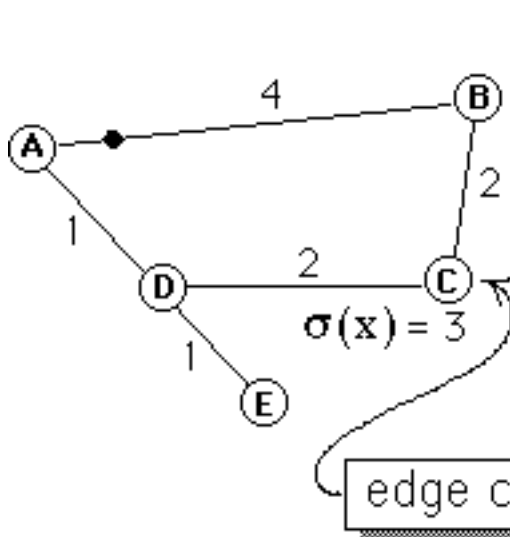
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$$\sigma(x) = \text{maximum}_{j \in N} d(x, j)$$

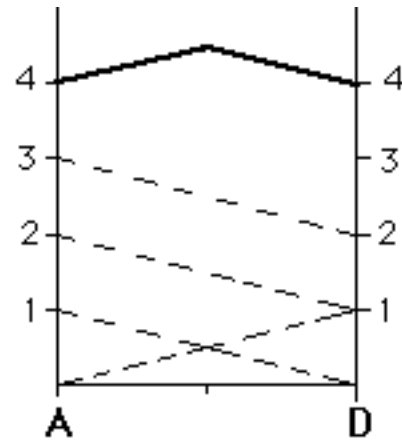
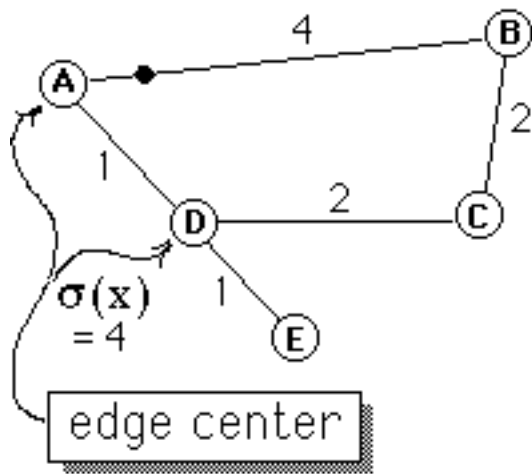
The absolute center may be found by computing each edge center, and selecting the best.

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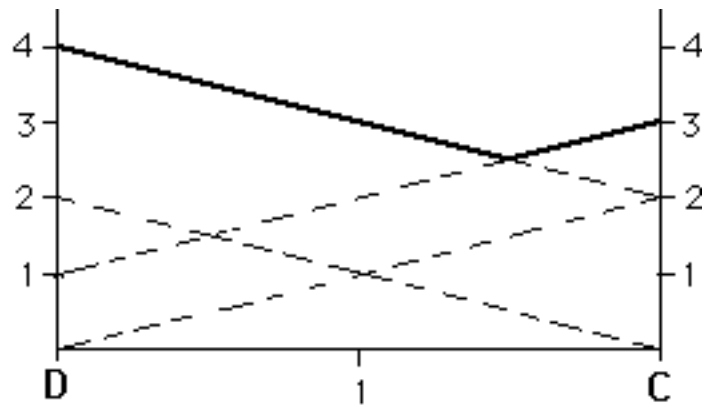
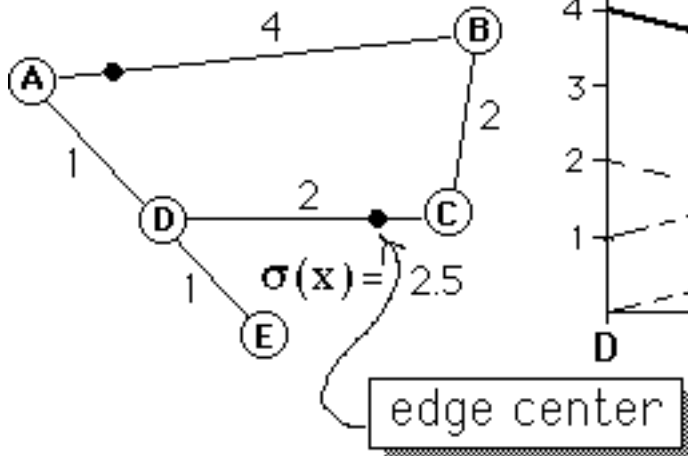
$$\sigma(x) = \text{maximum}_{j \in N} d(x, j)$$

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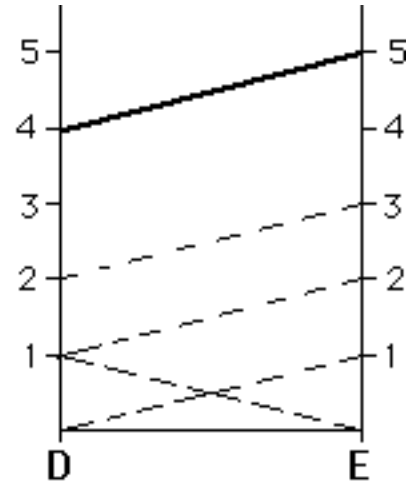
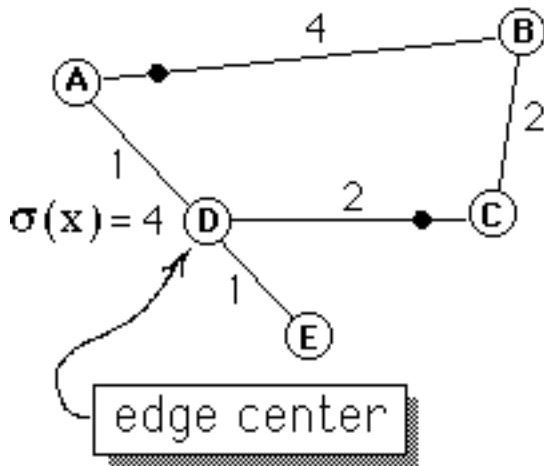
$$\sigma(x) = \text{maximum}_{j \in N} d(x, j)$$

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$$\sigma(x) = \text{maximum}_{j \in N} d(x, j)$$

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$$\sigma(x) = \max_{j \in N} d(x, j)$$

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| edge | edge center | $\sigma(x)$ |
|-------|--------------|-------------|
| [A,B] | 0.5 from A | 3.5 |
| [B,C] | Vertex C | 3 |
| [C,D] | 0.5 from C | 2.5 |
| [A,D] | Vertices A&D | 4 |
| [D,E] | Vertex D | 4 |

absolute center of network

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Searching some edges for their centers may be avoided by using the lower bound provided by

Theorem

Let X_{pq} be the edge center of $[P,Q]$.

Then

$$\sigma(X_{pq}) \geq \frac{\sigma(P) + \sigma(Q) - d(P,Q)}{2}$$

If this lower bound exceeds σ at the vertex center of the network, then the absolute center cannot lie on this edge!



Proof

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Proof:

$$d(X,j) \leq \sigma(X) \quad \forall j$$

$$d(P,j) \leq d(P,X) + d(X,j)$$

$$d(P,j) \leq d(P,X) + \sigma(X)$$

$$\text{But } \sigma(P) = \max_j \{ d(P,j) \},$$

$$\Rightarrow \sigma(P) \leq d(P,X) + \sigma(X)$$

$$\text{Likewise, } \sigma(Q) \leq d(Q,X) + \sigma(X)$$

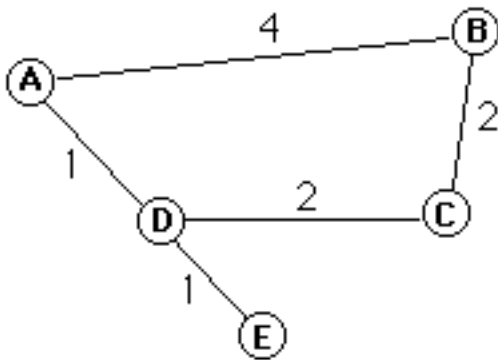
Sum these two inequalities:

$$\sigma(P) + \sigma(Q) \leq 2 \sigma(X) + \underbrace{d(P,X) + d(X,Q)}_{d(P,Q)}$$

$$\Rightarrow \sigma(X) \geq \frac{\sigma(P) + \sigma(Q) - d(P,Q)}{2}$$



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| vertex | σ |
|--------|----------|
| A | 4 |
| B | 5 |
| C | 3 |
| D | 4 |
| E | 5 |

vertex center



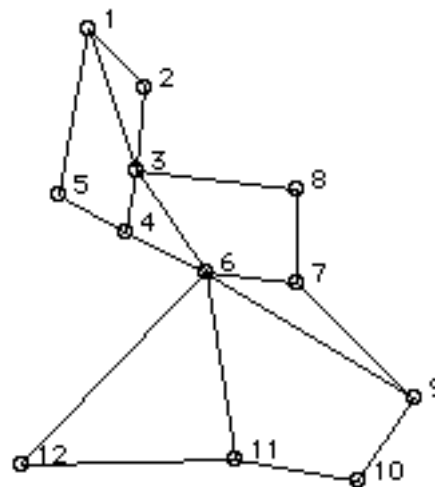
| edge | lower bound | min $\sigma(X)$ |
|-------|-------------|-----------------|
| [A,B] | 2.5 | 3.5 |
| [B,C] | 3 | 3 |
| [C,D] | 2.5 | 2.5 |
| [A,D] | 3.5 | 4 |
| [D,E] | 4 | 4 |

Using the lower bound would have eliminated 3 edges from consideration!

The edge centers needed to be found only for [A,B] & [C,D]

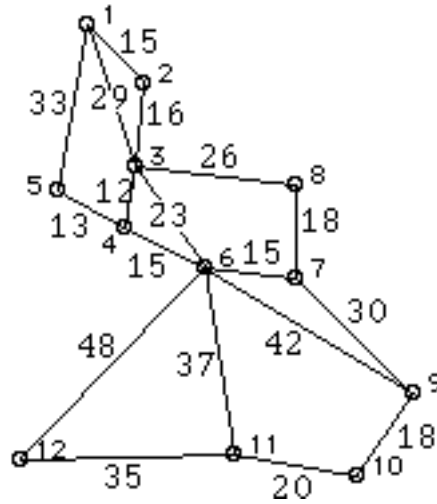
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Example



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Example



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Shortest Path Lengths

| from \ to | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | $\sigma(X)$ |
|-----------|-----|----|----|----|----|----|----|----|----|-----|----|-----|-------------|
| 1 | 0 | 15 | 29 | 41 | 33 | 52 | 67 | 55 | 94 | 109 | 89 | 100 | 109 |
| 2 | 15 | 0 | 16 | 28 | 41 | 39 | 54 | 42 | 81 | 96 | 76 | 87 | 96 |
| 3 | 29 | 16 | 0 | 12 | 25 | 23 | 38 | 26 | 65 | 80 | 60 | 71 | 80 |
| 4 | 41 | 28 | 12 | 0 | 13 | 15 | 30 | 38 | 57 | 72 | 52 | 63 | 72 |
| 5 | 33 | 41 | 25 | 13 | 0 | 28 | 43 | 51 | 70 | 85 | 65 | 76 | 85 |
| 6 | 52 | 39 | 23 | 15 | 28 | 0 | 15 | 33 | 42 | 57 | 37 | 48 | 57 |
| 7 | 67 | 54 | 38 | 30 | 43 | 15 | 0 | 18 | 30 | 48 | 52 | 63 | 67 |
| 8 | 55 | 42 | 26 | 38 | 51 | 33 | 18 | 0 | 48 | 66 | 70 | 81 | 81 |
| 9 | 94 | 81 | 65 | 57 | 70 | 42 | 30 | 48 | 0 | 18 | 38 | 73 | 94 |
| 10 | 109 | 96 | 80 | 72 | 85 | 57 | 48 | 66 | 18 | 0 | 20 | 55 | 109 |
| 11 | 89 | 76 | 60 | 52 | 65 | 37 | 52 | 70 | 38 | 20 | 0 | 35 | 89 |
| 12 | 100 | 87 | 71 | 63 | 76 | 48 | 63 | 81 | 73 | 55 | 35 | 0 | 100 |

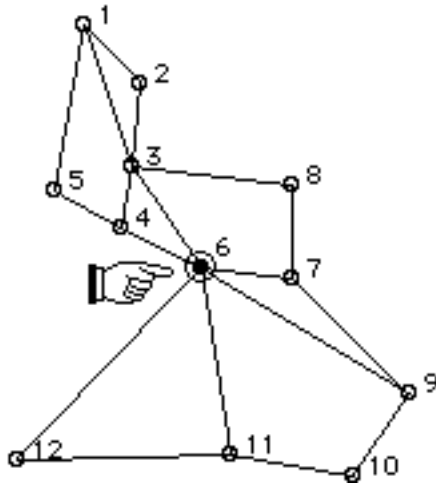
vertex center

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Vertex Center of Network

(Which minimizes the maximum distance to farthest nodes)

Vertex center of the network is at node 6
 where maximum distance (to node 10) is 57



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Lower Bound of σ
on the edges

| i | j | LB |
|----|----|------|
| 6 | 7 | 54.5 |
| 6 | 9 | 54.5 |
| 6 | 11 | 54.5 |
| 6 | 12 | 54.5 |
| 3 | 6 | 57 |
| 4 | 6 | 57 |
| 7 | 8 | 65 |
| 7 | 9 | 65.5 |
| 3 | 8 | 67.5 |
| 3 | 4 | 70 |
| 4 | 5 | 72 |
| 11 | 12 | 77 |
| 1 | 3 | 80 |
| 2 | 3 | 80 |
| 1 | 5 | 80.5 |
| 10 | 11 | 89 |
| 9 | 10 | 92.5 |
| 1 | 2 | 95 |

eliminated by L.B.

Only 4 edges could have edge centers better than the vertex center where $\sigma = 57$

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The function σ on edge [6,7]

Monotonically increasing distance functions: $d(x,k)$ where
 $k = \quad 1 \quad 2 \quad 3 \quad 4 \quad 5 \quad 6 \quad 11 \quad 12$
 $d(i,k) = 52 \quad 39 \quad 23 \quad 15 \quad 28 \quad 0 \quad 37 \quad 48$
 $d(j,k) = 67 \quad 54 \quad 38 \quad 30 \quad 43 \quad 15 \quad 52 \quad 63$

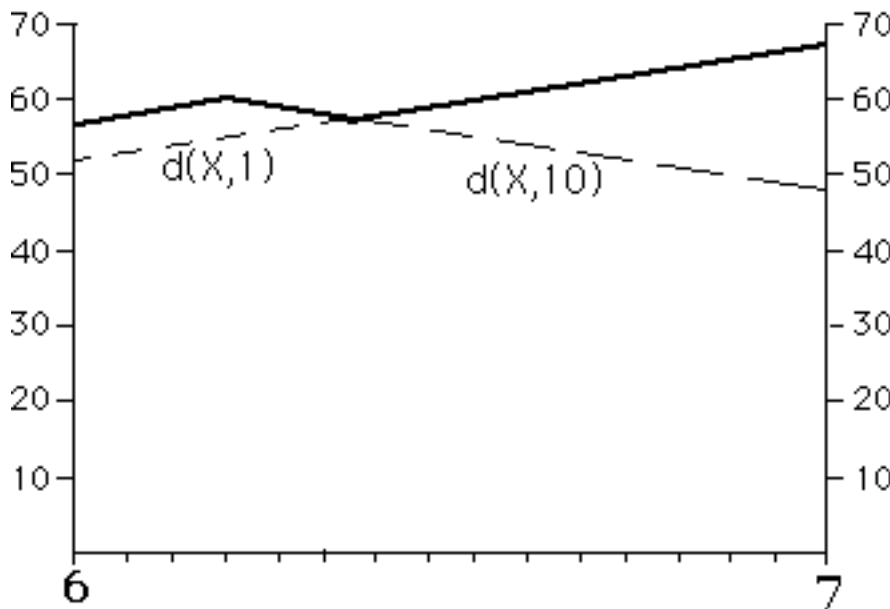
Monotonically decreasing distance functions: $d(x,k)$ where
 $k = \quad 7 \quad 8$
 $d(i,k) = 15 \quad 33$
 $d(j,k) = 0 \quad 18$

Distance functions which increase to a peak at a point Δ units from i , then decrease: $d(x,k)$ where

$k = \quad 9 \quad 10$
 $d(i,k) = 42 \quad 57$
 $d(j,k) = 30 \quad 48$
 $\Delta = \quad 1.5 \quad 3$

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The function σ on edge [6,7]



The edge center is at vertex #6

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The function σ on edge [6,9]

Monotonically increasing distance functions: $d(x,k)$ where

| | | | | | | |
|-----------|----|----|----|----|----|----|
| k= | 1 | 2 | 3 | 4 | 5 | 6 |
| $d(i,k)=$ | 52 | 39 | 23 | 15 | 28 | 0 |
| $d(j,k)=$ | 94 | 81 | 65 | 57 | 70 | 42 |

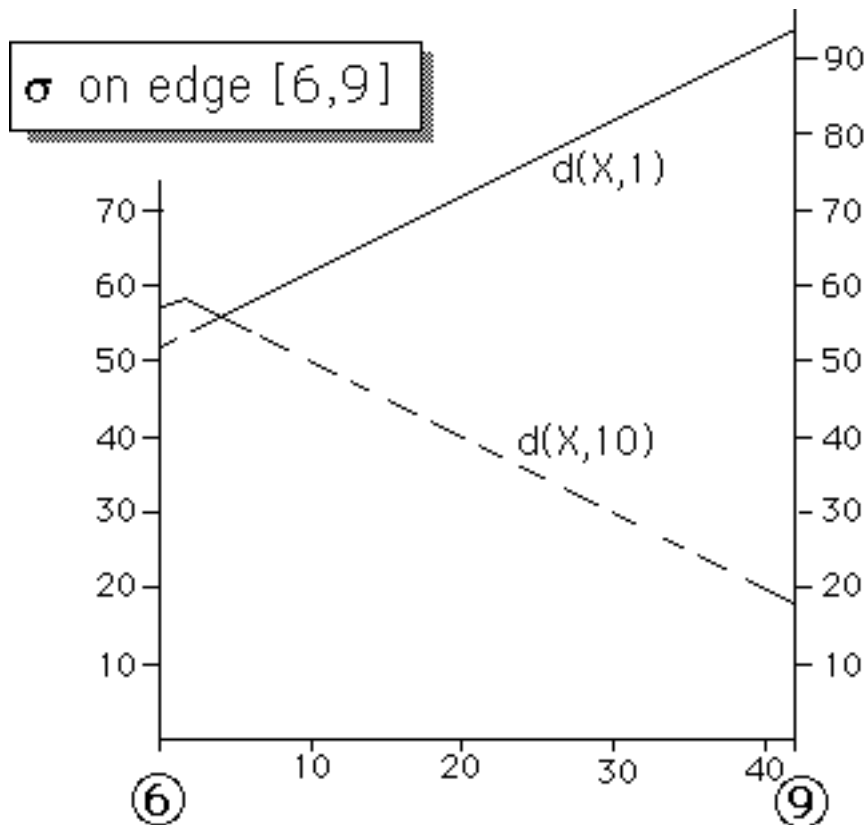
Monotonically decreasing distance functions: $d(x,k)$ where

| | |
|-----------|----|
| k= | 9 |
| $d(i,k)=$ | 42 |
| $d(j,k)=$ | 0 |

Distance functions which increase to a peak at a point Δ units from i , then decrease: $d(x,k)$ where

| | | | | | |
|-----------|------|------|-----|------|------|
| k= | 7 | 8 | 10 | 11 | 12 |
| $d(i,k)=$ | 15 | 33 | 57 | 37 | 48 |
| $d(j,k)=$ | 30 | 48 | 18 | 38 | 73 |
| $\Delta=$ | 28.5 | 28.5 | 1.5 | 21.5 | 33.5 |

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The edge center is 4 units from vertex 6, with $\sigma(X) = 56$

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The function σ on edge [6,11]

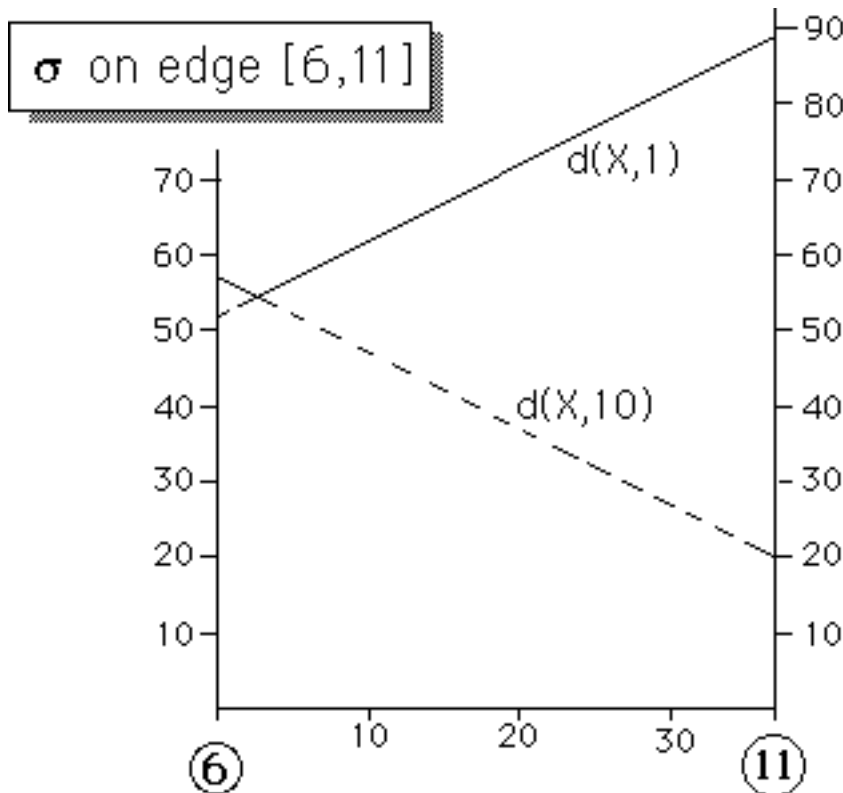
Monotonically increasing distance functions: $d(x,k)$ where
 $k=$ 1 2 3 4 5 6 7 8
 $d(i,k)=$ 52 39 23 15 28 0 15 33
 $d(j,k)=$ 89 76 60 52 65 37 52 70

Monotonically decreasing distance functions: $d(x,k)$ where
 $k=$ 10 11
 $d(i,k)=$ 57 37
 $d(j,k)=$ 20 0

Distance functions which increase to a peak at a point Δ units from i , then decrease: $d(x,k)$ where

$k=$ 9 12
 $d(i,k)=$ 42 48
 $d(j,k)=$ 38 35
 $\Delta=$ 16.5 12

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The edge center is 2.5 units from vertex #6, with $\sigma(X) = 54.5$

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The function σ on edge [6,12]

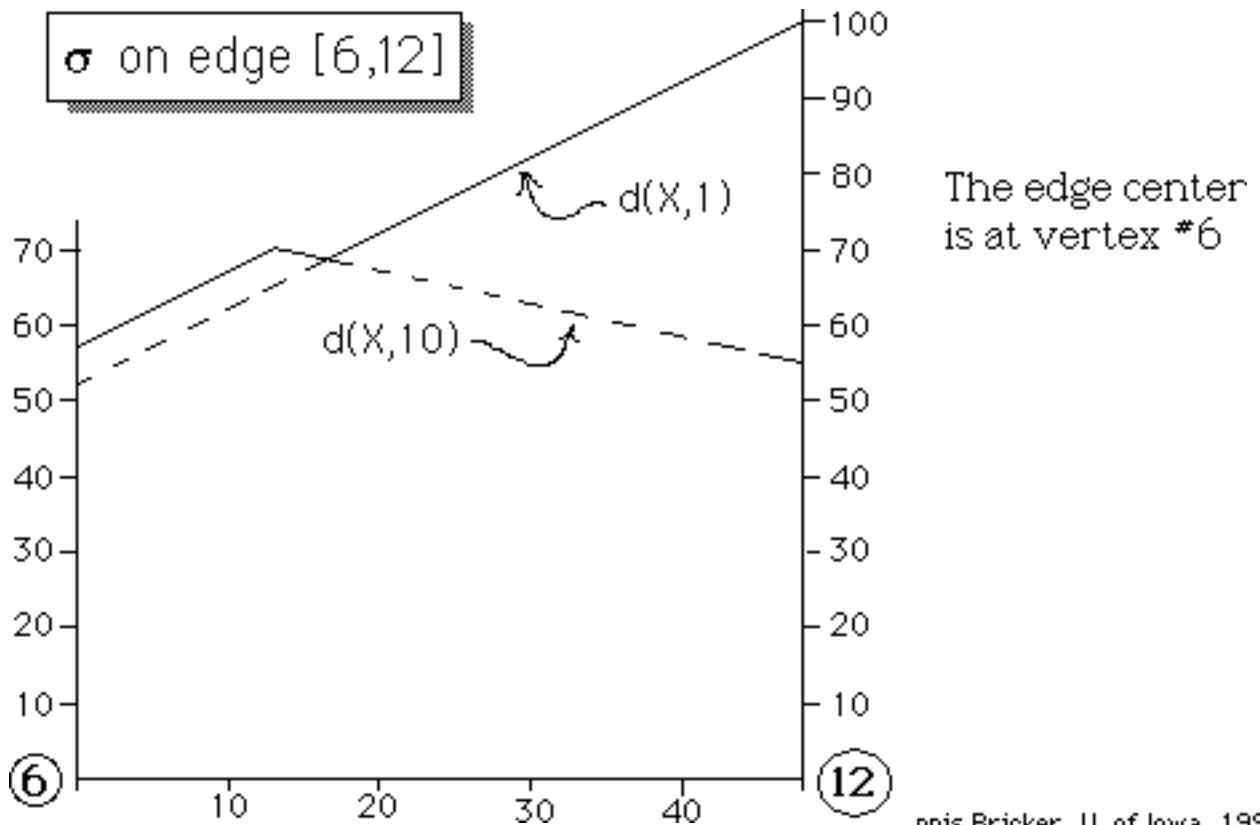
Monotonically increasing distance functions: $d(x,k)$ where
 $k=$ 1 2 3 4 5 6 7 8
 $d(i,k)=$ 52 39 23 15 28 0 15 33
 $d(j,k)=$ 100 87 71 63 76 48 63 81

Monotonically decreasing distance functions: $d(x,k)$ where
 $k=$ 12
 $d(i,k)=$ 48
 $d(j,k)=$ 0

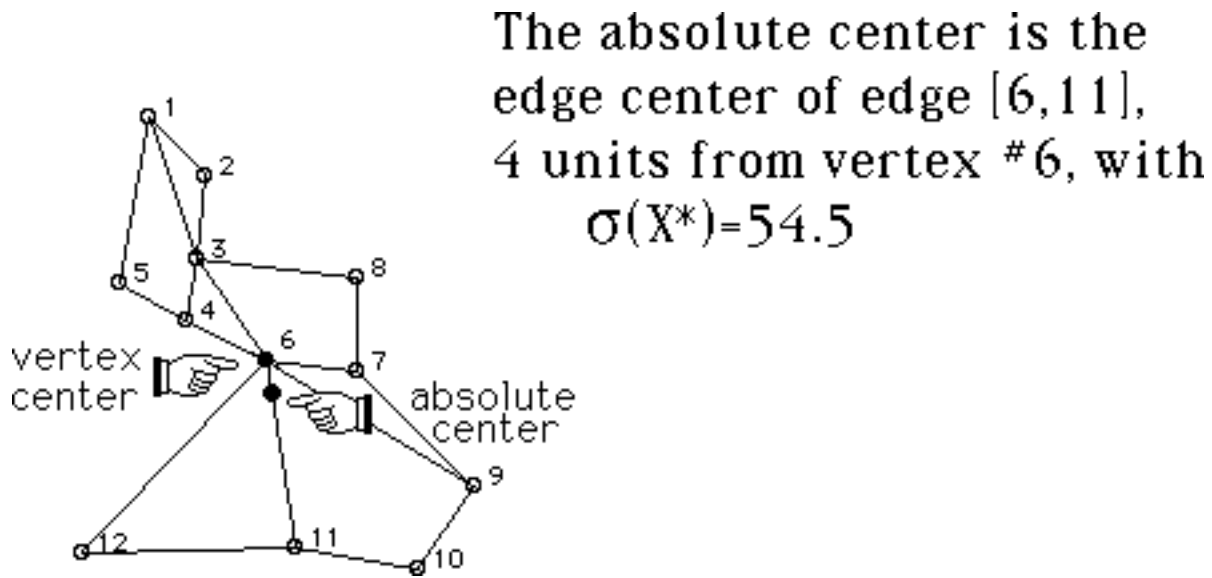
Distance functions which increase to a peak at a point Δ units from i , then decrease: $d(x,k)$ where

$k=$ 9 10 11
 $d(i,k)=$ 42 57 37
 $d(j,k)=$ 73 55 35
 $\Delta=$ 39.5 23 23

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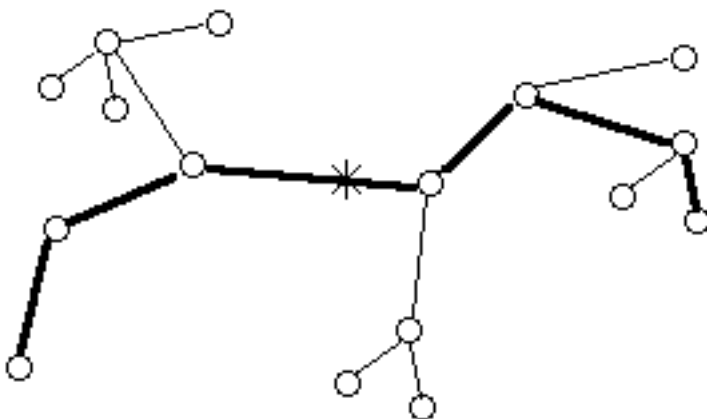
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Center of a Tree

A center of a tree lies at the midpoint of the longest elementary chain in the tree.



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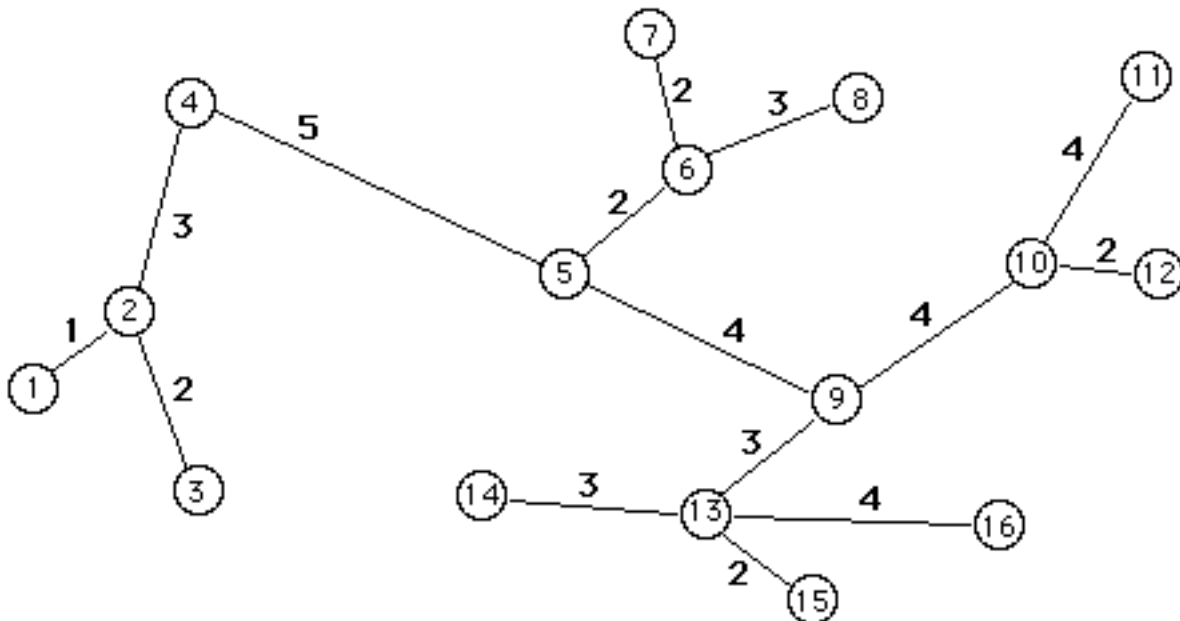
Finding Center of a Tree

0. Choose arbitrarily a point X of the tree.
1. Find the vertex *farthest* from X . Call this vertex V_1 . (This will have degree 1.)
2. Find the vertex *farthest* from V_1 . Call this vertex V_2 . (This will also have degree 1.)
3. Find the midpoint X^* of the unique elementary path from V_1 to V_2 . X^* will be the *absolute* center of the tree, and the vertex nearest to X^* will be the *vertex* center.

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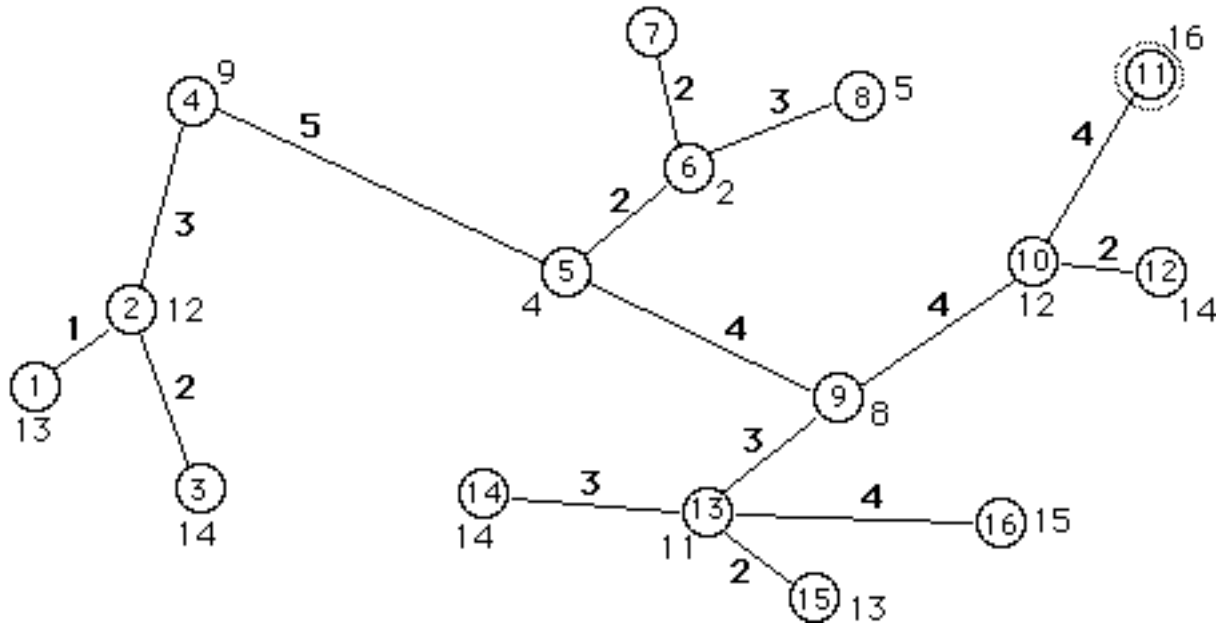
Example

Find the absolute & vertex centers of the tree:



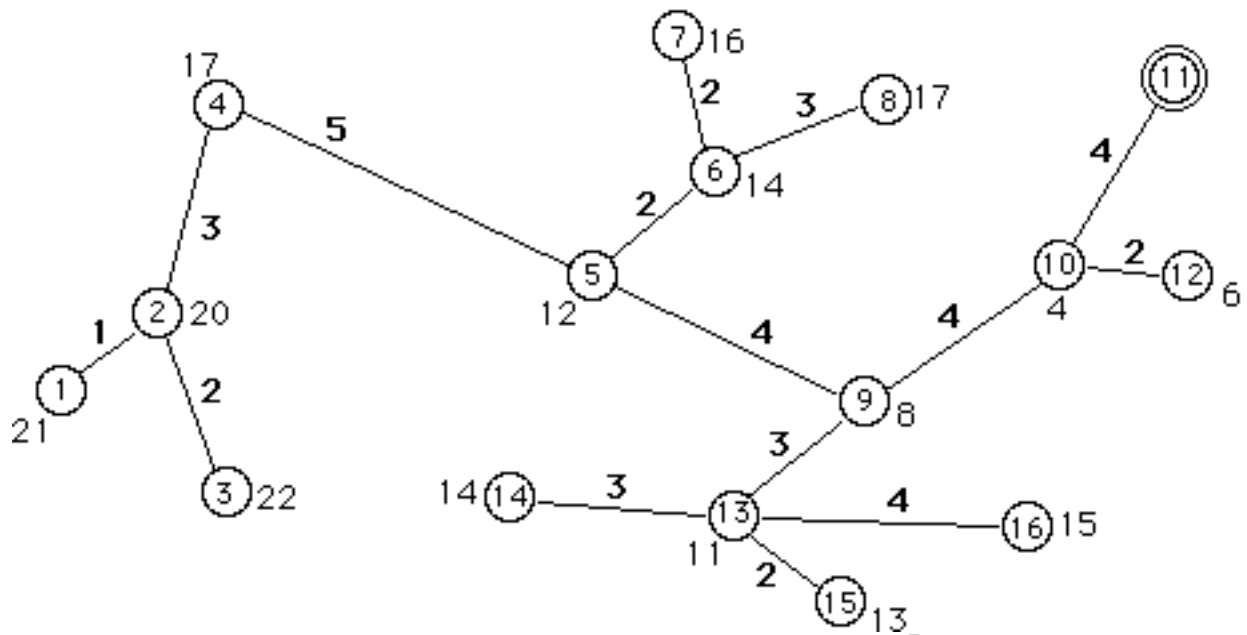
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Arbitrarily choose vertex 7. Label each vertex with its distance from vertex 7, to find that farthest from #7: (vertex #11)



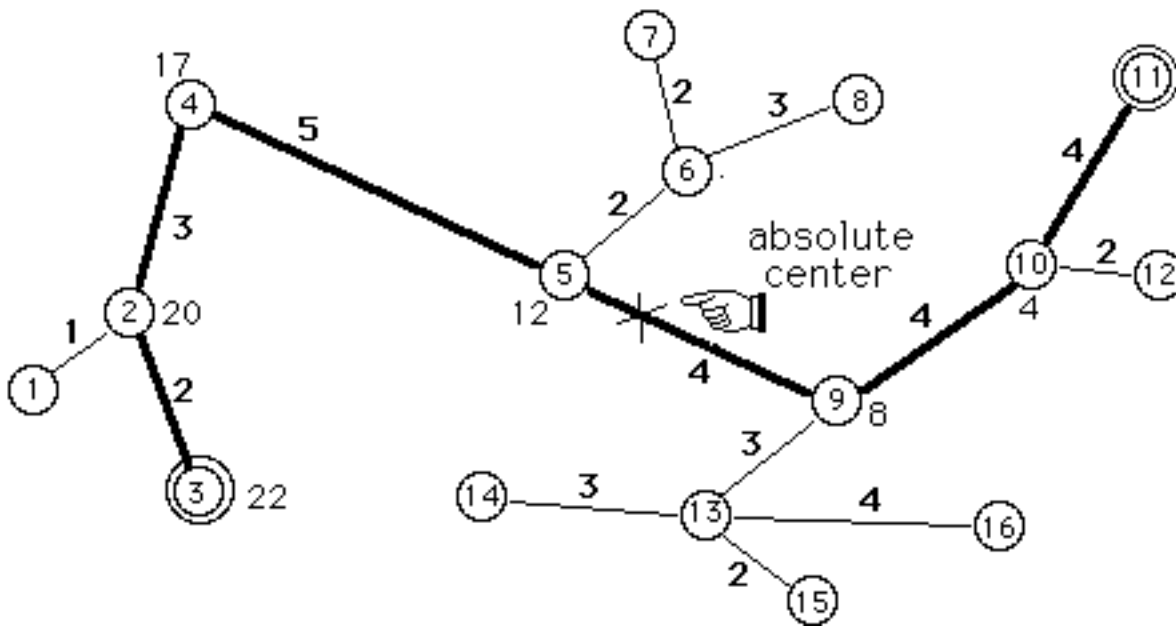
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Now label the vertices with their distances from vertex #11, to find that farthest from #11: vertex 3.



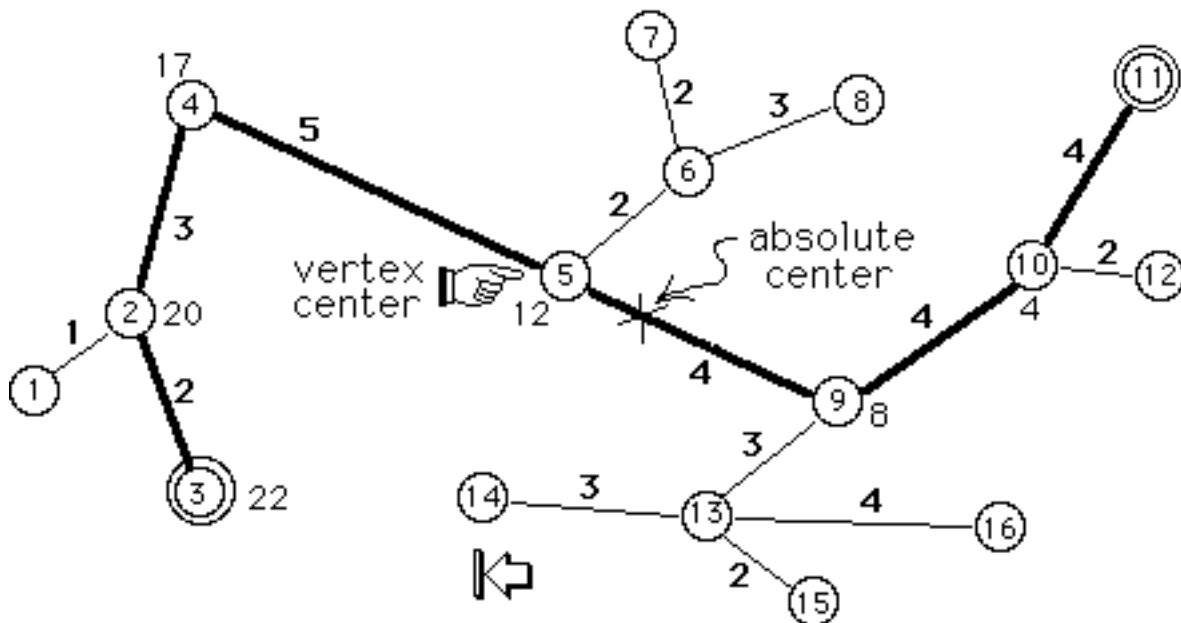
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The midpoint of the chain from vertex 11 to vertex 3 is a distance 11 from vertex 11, on the edge [5,9]



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The vertex center of the tree is at vertex #5, the vertex nearest to the absolute center.



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