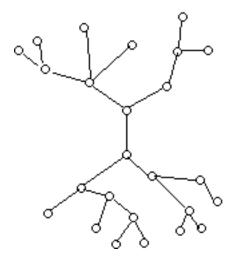
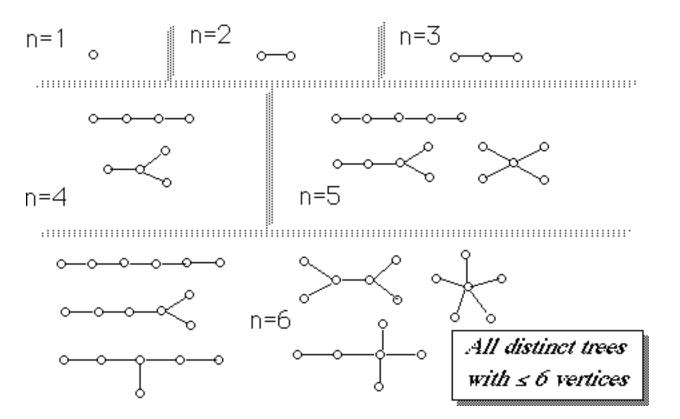


### T ∩ E ∈ : a connected graph without cycles



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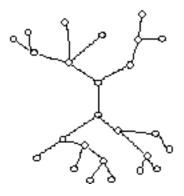
page 3



8/20/00

The following statements about a graph G are equivalent:

- G is a tree
- G is connected with n vertices and n-1 edges
- G has n vertices, n-1 edges, and no cycles
- G is such that each pair of vertices is connected by a *unique* elementary chain

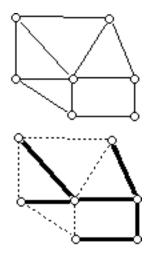


# Spanning tree

A spanning tree of a connected graph G=(V,A) is a tree with vertex set V and an edge set which is a subset of A

Minimum spanning tree

A minimum spanning tree of a *network* is a spanning tree the sum of whose edge lengths are minimal.



### Two algorithms for MST problem:

## r Prim's Algorithm

Beginning with a single node, at each iteration a tree is obtained by adding an edge & node, until ALL nodes have been included.



### <sup>°</sup> Kruskal's Algorithm

Beginning with N trees, each consisting of a single node, at each iteration two trees are combined by adding an edge, until a single tree is obtained.

8/20/00

**Finding a Minimum Spanning Tree (MST) of a Network** (Prim's algorithm)

```
Step 1 (Setup)
```

Select any node to begin the tree

```
Step 2 (Addition)
```

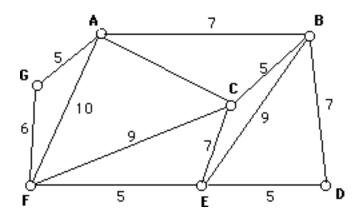
Find a node NOT currently in the tree which is nearest to the set of nodes IN the tree. Add that node and the connecting edge to the tree

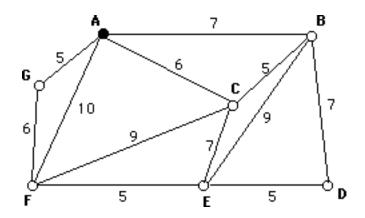
Step 3 (Stopping criterion)

If all nodes are in the tree, STOP; otherwise return to step 2

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Example: Prim's algorithm for MST

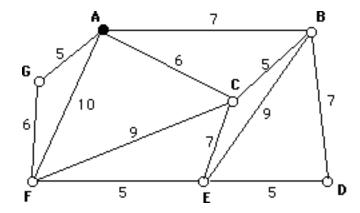




Initially, the tree is empty.

Select (arbitrarily) node A to add to the tree.

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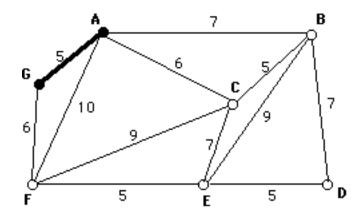


Find the node which is nearest to the nodes of the tree (i.e., node A)

This is node G.

Add it (and edge [A,G]) to the tree

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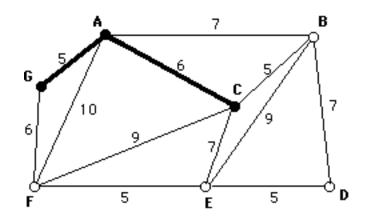


Find the node in the set {B,C,D,E,F} (not in the tree) which is nearest to the nodes {A,G} which are in the tree.

In this case there is a tie!

Break the tie arbitrarily, by selecting node C.

Add node C (and edge [A,C]) to the tree

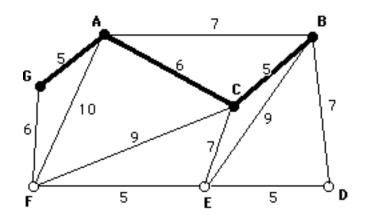


Find the node from the set {B,D,E,F} (not in the tree) which is nearest to the nodes {A,C,G} (in the tree)

This is node B, a distance 5 from the tree.

Add the node B (and the edge [B,C]) to the tree

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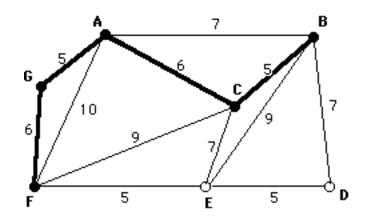


Find the node from the set {D,E,F} which is nearest to the set of nodes in the tree, {A,B,C,G}.

This is node F, a distance of 6 from the tree.

Add node F (and edge [F,G]) to the tree

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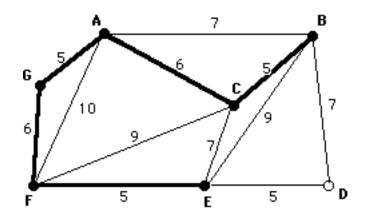


Find the node from the set {D,E} which is nearest to the nodes in the tree, {A,B,C,F,G}.

This is node E, a distance of 5 from the tree.

Add node E (and edge [E,F]) to the tree.

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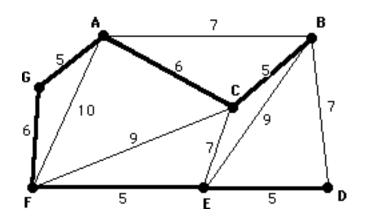


Find the node from the set {D} which is nearest to the nodes {A,B,C,E,F,G} in the tree.

This is node D, a distance of 5 from the tree.

Add node D (and edge [D,E]) to the tree.

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All nodes are now in the tree, so we stop!

@D.L.Bricker, U. of Iowa, 1998

#### Example:

to construct a pipeline to supply gas from Alaska's north slope ("NS") to eight U.S. gas companies, denoted by A through H.

Each mile of "rightof-way which is purchased costs an average of \$1000.

How should the pipeline be routed to minimize the total cost of the right-of-way?

f "rightch is costs an \$1000. the pipeted to he total \$1000. \$1000

ς.		ING	> A	D	C	υ	Ľ	Г	0	П
ניא הממונות במי	NS	0	32	43	41	44	45	53	56	61
3	Α	32	0	12	15	16	17	31	25	32
Ş	В	43	12	0	18	12	11	32	26	28
5	С	41	15	18	0	10	14	23	15	18
È.	D	44	16	12	10	0	5	22	13	16
Ğ.	Е	45	17	11	14	5	0	23	15	12
מוסומווניטס	F	53	31	32	23	22	23	0	7	14
j D	G	56	25	26	15	13	15	7	0	8
3	Н	61	32	28	18	16	12	14	8	0

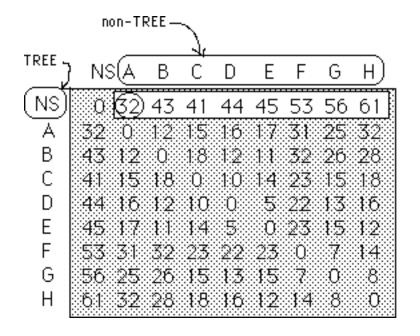
Ц

	NS	λ	В	С	D	Е	F	G	Н
NS	0	32	43	41	44	45	53	56	61
А	32	0	12	15	16	17	31	25	32
В	43	12	0	18	12	11	32	26	28
С				0					
D	44	16	12	10	0	5	22	13	16
E	45 53	17	11	14	5	0	23	15	12
F	53	31	32	23	22	23	0	7	14
G	56	25	26	15	13	15	7	0	8
Н	61	32	28	18	16	12	14	8	0

Arbitrarily select a node to begin the tree.

Let's choose node NS.

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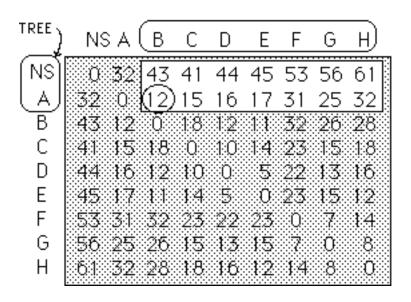
Find the minimum distance from a node NOT in the tree to the node IN the tree.

This is node A.

Add node A (and edge [NS,A]) to the tree.

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#### non-TREE



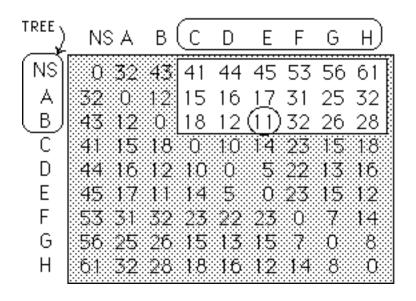
Find the node NOT in the tree which is nearest to the nodes IN the tree.

This is node B, a distance of 12 from node A.

Add node B (& edge [A,B]) to the tree.

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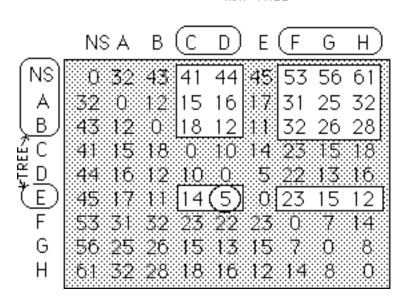


Find the node NOT in the tree which is nearest to the nodes IN the tree.

This is node E, a distance of 11 from node B.

Add node E (& edge [B,E]) to the tree.

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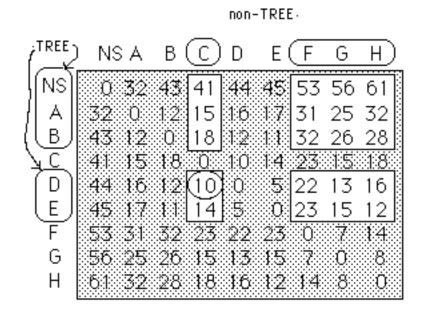
non-TREE.

Find the node NOT in the tree which is nearest to the nodes IN the tree.

This is node D, which is a distance 5 from node E.

Add node D (& edge [D,E]) to the tree.

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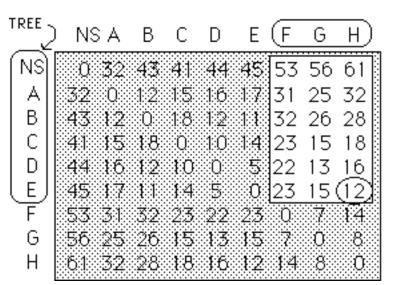


Find the node NOT in the tree which is nearest to the nodes IN the tree.

This is node C, a distance of 10 from node D.

Add node C (& edge [C,D]) to the tree.

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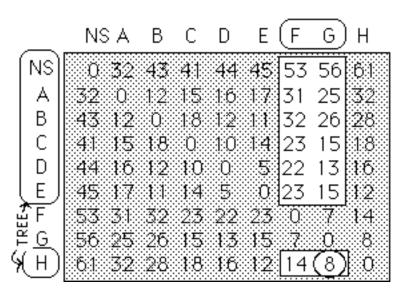
non-TREE.

Find the node NOT in the tree which is nearest to the nodes IN the tree.

This is node H, a distance of 12 from node E.

Add node H (& edge [E,H]) to the tree.

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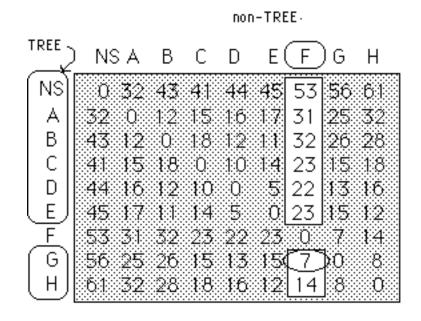
non-TREE.

Find the node NOT in the tree which is nearest to the nodes IN the tree.

This is node G, a distance of 8 from node H.

Add node G (& edge [G,H]) to the tree.

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Find the node NOT in the tree which is nearest to the nodes IN the tree.

This is node F, a distance of 7 from node G.

Add node F (& edge [F,G]) to the tree.

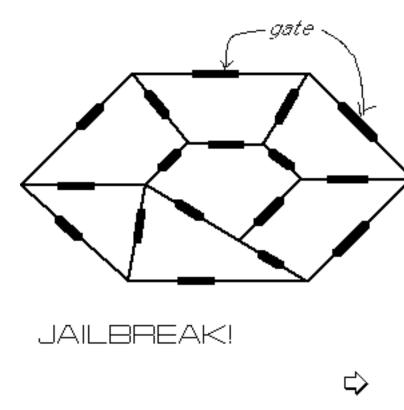
The tree now spans all nine nodes, and is the Minimum Spanning Tree. 8/20/00

#### APL code for Prim's MST algorithm

```
VTREE←MST C;IN;OUT;K;L;ROWMIN;MIN;J
```

```
[1]
      Μ.
[2]
      м
         Compute Minimum Spanning Tree of a graph
[3]
      Θ.
[4]
      IN←,1
                                     A List of nodes in tree
[5]
      OUT←1+1<sup>-</sup>1+1↑ρC
                                     A List of nodes not yet in
[6]
      TREE←(pC)pO
      LENGTH←0
[7]
[8]
                 Find shortest arc joining IN & OUT nodes
[9]
      NEXT:ROWMIN←L/CLIN:OUT]
[10]
      MIN←L/ROWMIN
[11]
      J←ROWMIN1MIN
[12]
                 Add arc from IN node (K) to OUT node (L)
      Θ.
[13] K←IN[J]
[14] L←OUT[C[K;OUT]1MIN]
[15] TREE[K;L]←1
[16]
      OUT+(L≠OUT)/OUT
      IN←IN,L
[17]
[18] LENGTĤ←LENGTH+MIN
      →NEXT IF (pIN)<1↑pC</p>
[19]
     \nabla
```

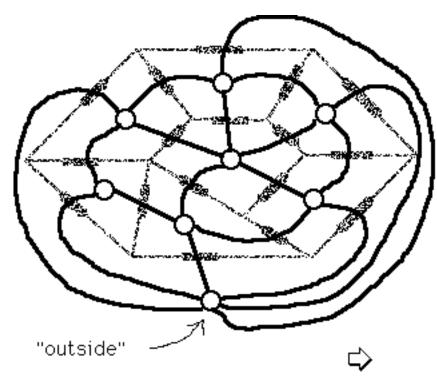
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 Prisoners have been divided into seven groups by walls

• An outside accomplice plans to help them to escape by blowing up some of the gates, using explosives

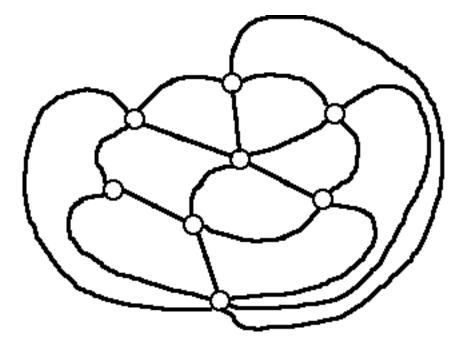
HOW CAN HE DO THIS, DESTROYING AS FEW GATES AS POSSIBLE?



Represent each room, together with the "outside world", by a node, and each gate by an edge.

The problem is to find a spanning tree with the fewest edges!

<sup>@</sup>D.L.Bricker, U. of Iowa, 1998



The number of nodes is 8.

All spanning trees will have seven edges!

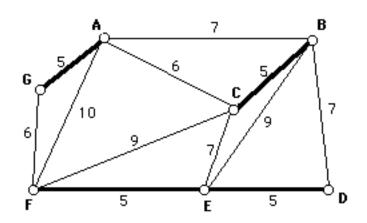
@D.L.Bricker, U. of Iowa, 1998

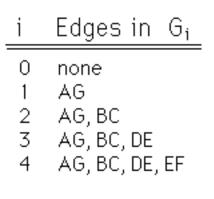
### Kruskal's Algorithm for MST

```
Step 1: Setup
 Let G_0 = (V, \emptyset) and i = 0
Step 2: Addition of Edge
 Find (x,y) which minimizes w(x,y), and set w(x,y) = +\infty
Step 3: Test for cycle
 If the addition of edge (x,y) to the graph G<sub>1</sub> would form a cycle,
    then go to step 2;
 Otherwise, add edge (x,y) to graph G<sub>1</sub> and increment i.
Step 4: Test for termination
 If i < n -1, then return to step 2.
 Otherwise, stop with G_{n-1} = MST
```

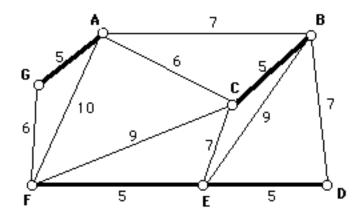
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Example (Kruskal's MST Algorithm)





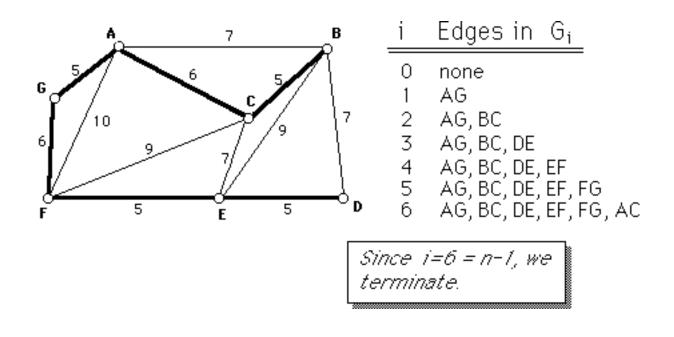
*In each of the first 4 iterations, there is a tie for the minimumlength edge to be added* 

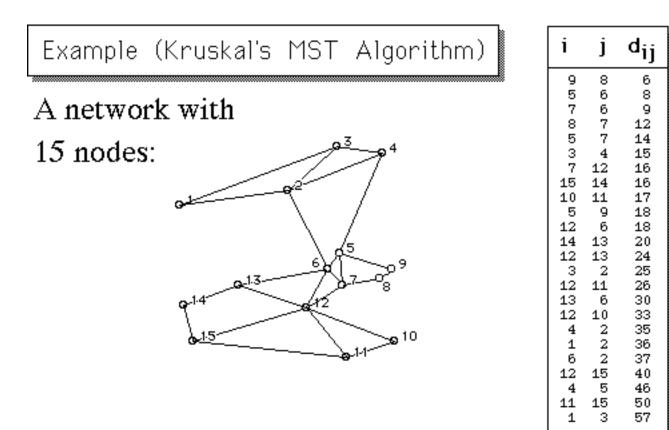


*Next, there is a tie between edges FG and AC* 

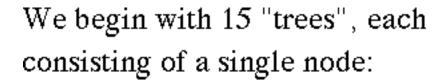
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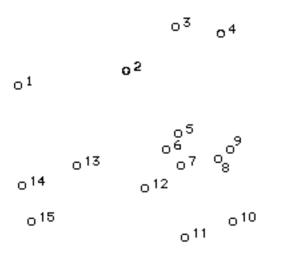
page 34

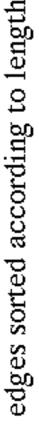




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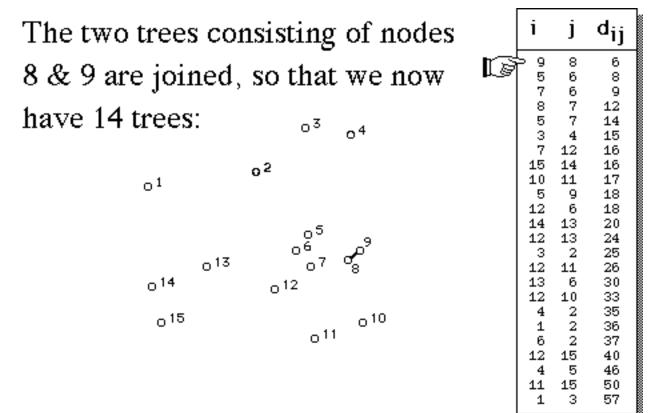


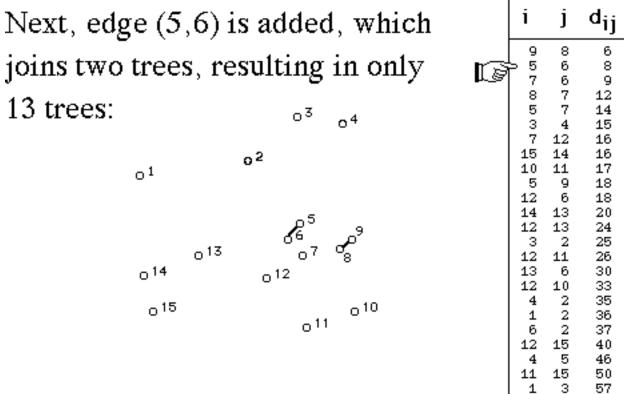




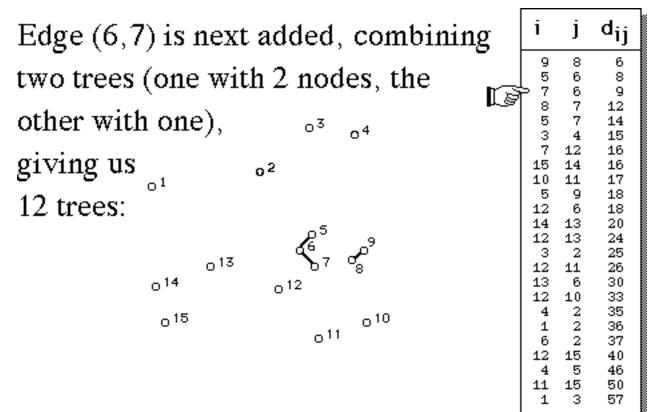
i	j	$d_{ij}$
9	8	6
5	6	8
7	6	9
5 7 8 5	6 7 7	12
		14
3	4	15
7	12	16
15	14	16
10	11	17
5	9	17 18
12	6	18
14	13	20 24
12	13	24
з	2	25
12	11	26
13	6 10	30
12	10	33
4	2 2 2 15	35
1	2	36
6	2	37
12	15	40
4	5	46
11	15	50
1	3	57

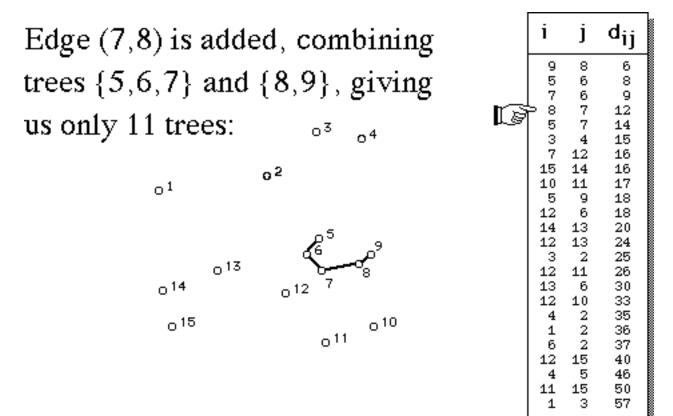
⊚D





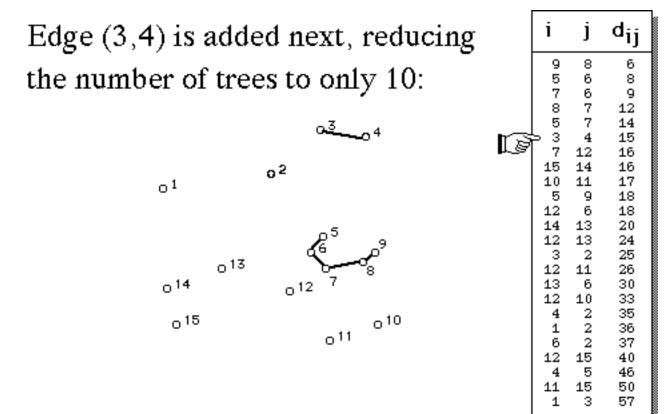
joins two trees, resulting in only 13 trees: 01 o13 o 14 <sub>0</sub>15

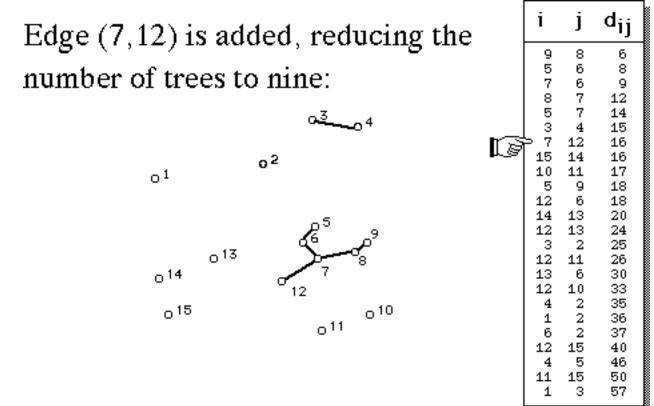




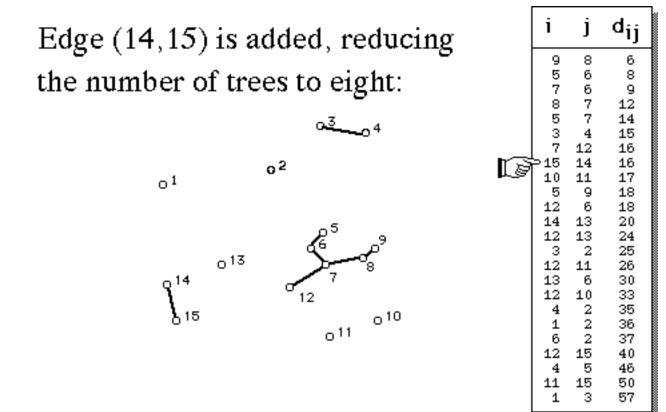
d<sub>ij</sub> If edge (5,7) were added, a cycle 9578537 8 6 7 4 12 14 5-7-6-5 would be formed, and so we "skip" this edge. 53 o4 Íð 15  $o^2$ 10  $\begin{array}{r}
 11 \\
 9 \\
 13 \\
 11 \\
 12 \\
 2 \\
 2 \\
 15 \\
 15 \\
 \end{array}$ 01 o13 o 14 012 o15 o10 o 11 4 11 50 1 з 57

Trees

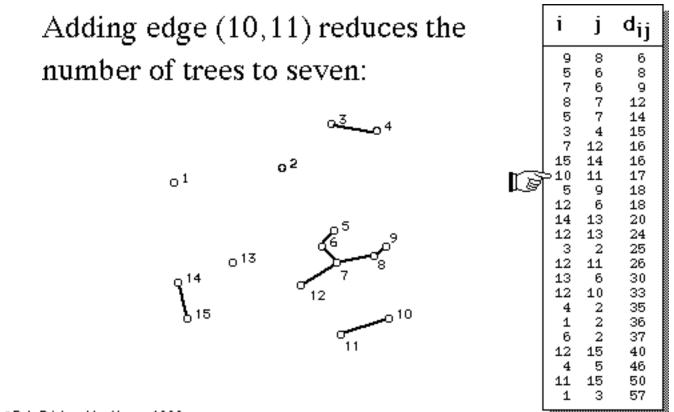




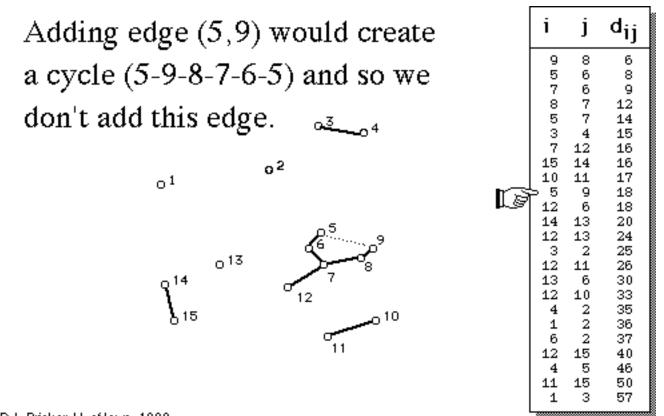
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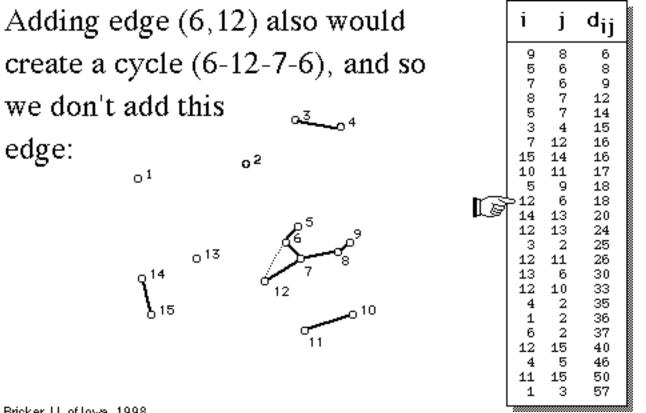


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dii

6

 $\begin{array}{c} 16 \\ 17 \\ 18 \\ 20 \\ 24 \\ 25 \\ 26 \\ 30 \\ 33 \\ 35 \\ 36 \\ 37 \\ 40 \\ 46 \end{array}$ 

50

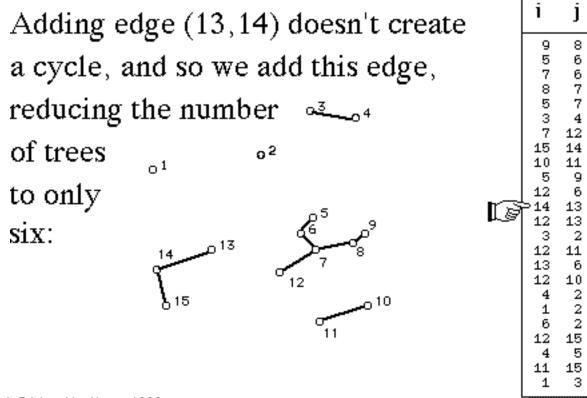
57

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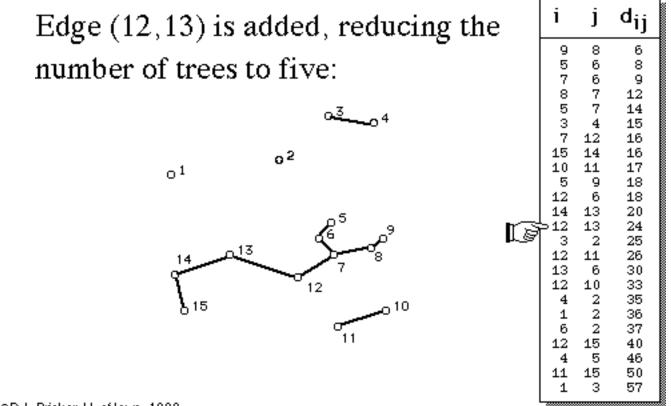
9 6

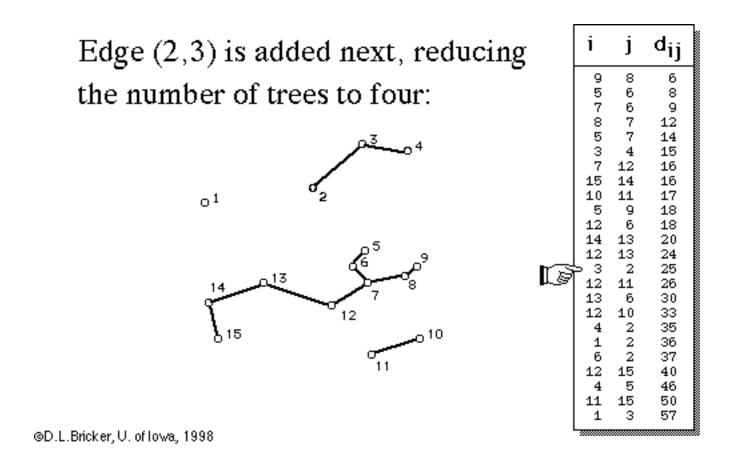
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3



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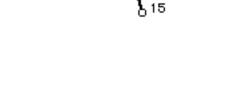
11

.10

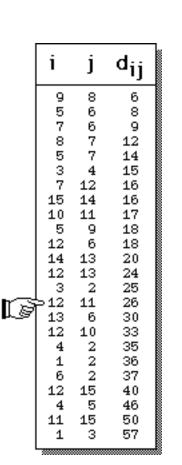
12

Next we add edge (11, 12), to obtain three trees:

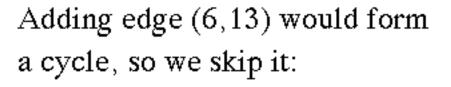
n 13

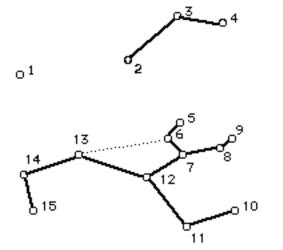


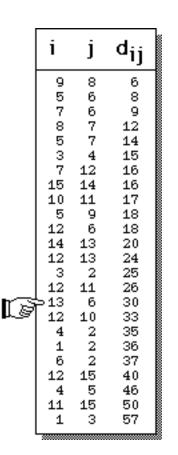
01



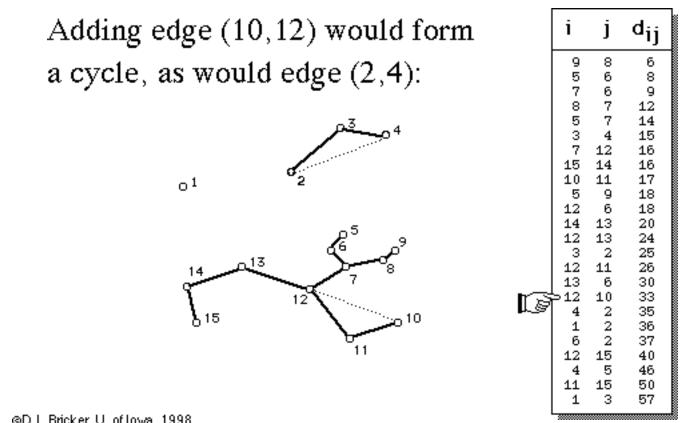
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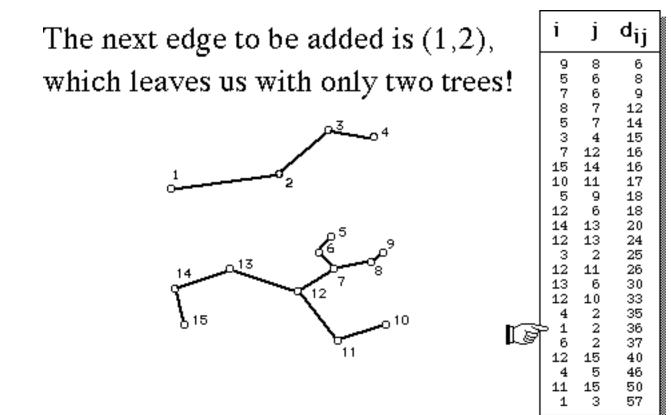




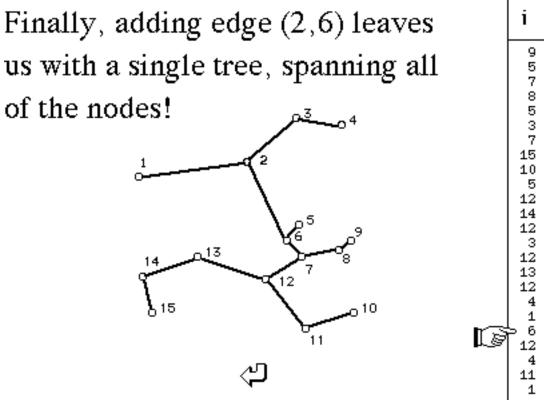
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Trees



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Finally, adding edge (2,6) leaves us with a single tree, spanning all

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dii

6

 $\begin{array}{c} 16\\ 17\\ 18\\ 20\\ 24\\ 25\\ 26\\ 30\\ 33\\ 35\\ 36\\ 37\\ 40\\ 46 \end{array}$ 

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57

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12

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