

**Parametric Programming
on the
Right-Hand-Side**

Parametric Programming is the analysis of the variation of the solution of an LP when some element (right-hand-side, objective, etc.) varies.

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Consider the optimal value of the LP as a function of the right-hand-side vector, i.e.,

$$Z^*(b) = \left\{ \begin{array}{l} \text{Min } c \cdot x \\ \text{s.t. } Ax \geq b \\ x \geq 0 \end{array} \right\} \stackrel{\text{by duality theory}}{=} \left\{ \begin{array}{l} \text{Max } \pi \cdot b \\ \text{s.t. } \pi A \leq c \\ \pi \geq 0 \end{array} \right\}$$

The function z^* is "evaluated" for some particular right-hand-side b' by solving the LP (either the primal or the dual).

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So we can evaluate $z^*(b')$ by solving the LP

$$z^*(b') = \left\{ \begin{array}{l} \text{Max } \pi \cdot b' \\ \text{s.t. } \pi A \leq c \\ \pi \geq 0 \end{array} \right\}$$

Notice that the feasible region of the dual LP is the same for every argument b' .

We know that a basic solution is optimal for an LP problem, and that there are a finite (but possibly very large!) number of such basic solutions.

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Suppose that we were to number the basic feasible solutions of the dual LP:

$\{ \pi^1, \pi^2, \pi^3, \dots, \pi^K \}$
 where each $\pi^k = c_{B_k} (A_{B_k})^{-1}$ for some dual-feasible basis B_k
 (K is a finite number, no greater than $\binom{n+m}{m}$.)

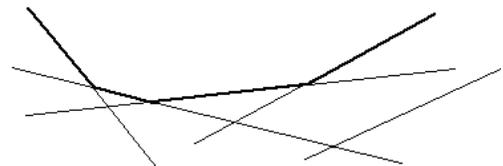
In theory, then, we could evaluate $z^*(b)$ by enumerating the K basic feasible solutions of the dual, evaluating the dual objective $\pi^k b$ at each, and selecting the maximum such value.

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Therefore, $z^*(b) = \text{Maximum}_{k=1,2,\dots,K} \{ \pi^k b \}$

That is, $z^*(b)$ is the maximum of a family of linear functions, $\pi^k b, k=1,2,\dots,K$

which is a piecewise linear convex function!



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Let us restrict our analysis of the function $z^(b)$ to a study of its behavior along a line (rather than everywhere in m -dimensional space!)*

That is, we assume an initial right-hand-side vector (b) is given, and a direction (d), and study the behavior of $z^(b + \lambda d)$, considered as a function of the scalar parameter λ .*

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Consider the solution of the LP

$$P_\lambda : \begin{array}{l} z^*(\lambda) = \text{minimum } c \cdot x \\ \text{s.t. } Ax = b + \lambda d \\ x \geq 0 \end{array}$$

where d is an m -vector and λ is a scalar.

$$z^*(\lambda) = \text{maximum}_{k=1,2,\dots,K} \{ \pi^k (b + \lambda d) \}$$

$$= \text{maximum}_{k=1,2,\dots,K} \left\{ \underbrace{\pi^k b}_{\text{intercept}} + \underbrace{(\pi^k d)}_{\text{slope}} \lambda \right\} \leftarrow \begin{array}{l} \text{Linear functions} \\ \text{of } \lambda \end{array}$$

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Example

$$z^*(\lambda) = \text{minimum } -x_1 - x_2$$

subject to

$$\begin{cases} 2x_1 + x_2 \leq 8 + 2\lambda \\ x_1 + 2x_2 \leq 7 + 7\lambda \\ x_2 \leq 3 + 2\lambda \\ x_1 \geq 0, x_2 \geq 0 \end{cases}$$

$$= \text{maximum } \{ (8\pi_1 + 7\pi_2 + 2\pi_3) + (2\pi_1 + 7\pi_2 + 2\pi_3)\lambda \}$$

s.t.

$$\begin{cases} 2\pi_1 + \pi_2 \leq -1 \\ \pi_1 + 2\pi_2 + \pi_3 \leq -1 \\ \pi_1 \leq 0, \pi_2 \leq 0, \pi_3 \leq 0 \end{cases}$$

nonpositive because of direction of \leq in the primal!

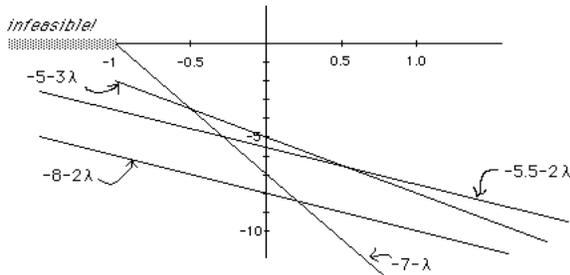
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π_1	π_2	π_3	intercept $\pi^k b$	slope $\pi^k d$	Basis
0	0	0	0	0	3 4 5
-0.5	0	0	-4	-1	1 4 5
-1	0	0	-8	-2	2 4 5
-0.333	-0.333	0	-5	-3	1 2 5
0	-1	0	-7	-7	1 3 5
0	-0.5	0	-3.5	-3.5	2 3 5
0	-1	1	-4	-5	1 2 3
-0.5	0	-0.5	-5.5	-2	1 2 4
0	0	-1	-3	-2	2 3 4

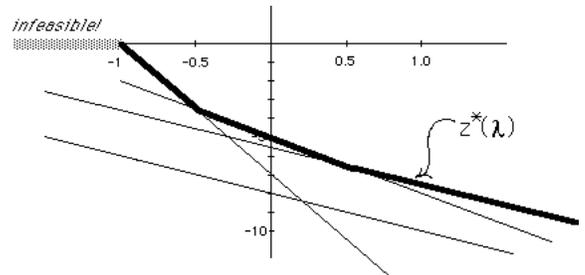
(columns #1,3,4 are dependent & do not form a basis!)

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Of the nine basic solutions, four are dual feasible. Therefore, $z^*(\lambda)$ is the maximum of four linear functions:



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In this example, with only nine basic dual solutions, it was possible to enumerate all of them, test each for feasibility, and then maximize the corresponding linear functions

$$\pi^k b + (\pi^k d)\lambda$$

However, for most problems, the number of basic dual solutions is astronomical and enumerating them is practically impossible.

(Usually, only a few of these basic dual solutions actually determine z^ .)*

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Let's consider again the parametric LP P_λ :

$$z^*(\lambda) = \text{minimum } -x_1 - x_2$$

subject to

$$\begin{cases} 2x_1 + x_2 \leq 8 + 2\lambda \\ x_1 + 2x_2 \leq 7 + 7\lambda \\ x_2 \leq 3 + 2\lambda \\ x_1 \geq 0, x_2 \geq 0 \end{cases}$$

Determine the optimal value function $z^*(\lambda)$ as well as $x_1^*(\lambda)$ and $x_2^*(\lambda)$ [i.e., the optimal solution] for all values of $\lambda \in (-\infty, +\infty)$

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The initial tableau:

-z	1	2	3	4	5	B	Δ
1	-1	-1	0	0	0	0	0
0	2	1	1	0	0	8	2
0	1	2	0	1	0	7	7
0	0	1	0	0	1	3	2

change column (pointing to column 3)

RHS is $B + \lambda \Delta$

slack variables (under columns 3, 4, 5)

Let's start with $\lambda = 0$, and investigate the LP as λ increases.

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The optimal tableau for $\lambda = 0$:

-z	1	2	3	4	5	B	Δ
1	0	0	0.333	0.333	0	5	3
0	1	0	0.667	-0.333	0	3	-1
0	0	1	-0.333	0.667	0	2	4
0	0	0	0.333	-0.667	1	1	-2

(this column updated during each pivot) (pointing to column 3)

The optimal solution of $P(0)$ is $z^*(0) = -5$,
 at $x_1^*(0) = 3, x_2^*(0) = 2, x_3^*(0) = 1,$
 $x_4^*(0) = x_5^*(0) = 0$

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Expressed as functions of λ , the basic solution is:

-z	1	2	5		B	Δ
1	0	0	0		5	3
0	1	0	0		3	-1
0	0	1	0		2	4
0	0	0	1		1	-2

 $\Rightarrow \begin{cases} z(\lambda) = -5 + 3\lambda \\ x_1^*(\lambda) = 3 - \lambda \\ x_2^*(\lambda) = 2 + 4\lambda \\ x_3^*(\lambda) = 1 - 2\lambda \\ x_4^*(\lambda) = x_5^*(\lambda) = 0 \end{cases}$

Note that these are linear functions of λ !

The optimality criterion (reduced cost ≥ 0) is independent of the parameter λ , and so the current basis remains optimal so long as the basic variables

$$\begin{cases} x_1^*(\lambda) = 3 - \lambda \\ x_2^*(\lambda) = 2 + 4\lambda \\ x_3^*(\lambda) = 1 - 2\lambda \end{cases}$$

remain feasible, i.e., nonnegative.

For what values of λ is $x^*(\lambda) \geq 0$?

We solve the inequalities

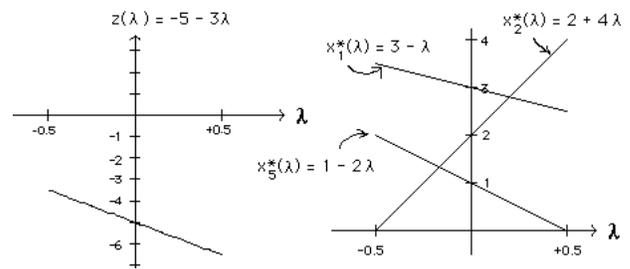
$$\begin{aligned} x_1^*(\lambda) &= 3 - \lambda \geq 0 \\ x_2^*(\lambda) &= 2 + 4\lambda \geq 0 \\ x_3^*(\lambda) &= 1 - 2\lambda \geq 0 \end{aligned}$$

for λ :

$$\begin{aligned} 3 - \lambda &\geq 0 \Rightarrow \lambda \leq 3 \\ 2 + 4\lambda &\geq 0 \Rightarrow \lambda \geq -1/2 \\ 1 - 2\lambda &\geq 0 \Rightarrow \lambda \leq 1/2 \end{aligned}$$

That is, the basic solution is feasible for all $\lambda \in [-1/2, 1/2]$ (and, in particular, for $\lambda = 0$).

Plot of $z(\lambda)$ & $x_i^*(\lambda)$ over the interval $[-0.5, +0.5]$:



Let us now increase λ from its initial value (0) to the upper limit for which the basis is feasible, i.e., $+1/2$. The basic solution then becomes

$$\begin{aligned} x_1^*(+1/2) &= 3 - 1/2 = 2.5 \\ x_2^*(+1/2) &= 2 + 4(1/2) = 4.0 \\ x_3^*(+1/2) &= 1 - 2(1/2) = 0 \end{aligned}$$

Any further increase in λ would result in infeasible (i.e., negative) values for x_3^* !

In order to increase the parameter λ further, x_3 must leave the basis, since it would otherwise become negative. In order to remove x_3 from the basis, we perform a DUAL SIMPLEX pivot:

-z	1	2	3	4	5	B	Δ
1	0	0	0.333	0.333	0	5	3
0	1	0	0.667	-0.333	0	3	-1
0	0	1	-0.333	0.667	0	2	4
0	0	0	0.333	-0.667	1	1	-2

pivot in this row! $\leftarrow \begin{cases} = 0 \text{ when} \\ \lambda = 1/2 \end{cases}$

-z	1	2	3	4	5	B	Δ
1	0	0	0.333	0.333	0	5	3
0	1	0	0.667	-0.333	0	3	-1
0	0	1	-0.333	0.667	0	2	4
0	0	0	0.333	-0.667	1	1	-2

-z	1	2	3	4	5	B	Δ
1	0	0	0.5	0	0.5	5.5	2
0	1	0	0.5	0	-0.5	2.5	0
0	0	1	0	0	1	3	2
0	0	0	-0.5	1	-1.5	-1.5	3

For this basis, the basic solution is

-z	1	2	4		B	Δ
1	0	0	0		5.5	2
0	1	0	0		2.5	0
0	0	1	0		3	2
0	0	0	1		-1.5	3

 $\Rightarrow \begin{cases} z(\lambda) = -5.5 - 2\lambda \\ x_1^*(\lambda) = 2.5 \\ x_2^*(\lambda) = 3 + 2\lambda \\ x_4^*(\lambda) = -1.5 + 3\lambda \end{cases}$

Notice that as λ increases, no basic variable decreases. Since the optimality criterion (reduced costs ≥ 0) does not depend upon λ , this basis is optimal for all $\lambda \geq 0.5$

That is, if we solve the system of inequalities

$$\begin{aligned} x_1^*(\lambda) = 2.5 &\geq 0 \implies (\text{no restriction on } \lambda) \\ x_2^*(\lambda) = 3 + 2\lambda &\geq 0 \implies \lambda \geq -1.5 \\ x_4^*(\lambda) = -1.5 + 3\lambda &\geq 0 \implies \lambda \geq 0.5 \end{aligned}$$

we see that it is satisfied for $\lambda \in [0.5, +\infty)$

Let us now investigate $P(\lambda)$ for $\lambda < -0.5$
Consider the tableau which was optimal for $-0.5 \leq \lambda \leq +0.5$

-z	1	2	3	4	5	B	Δ
1	0	0	0.333	0.333	0	5	3
0	1	0	0.667	-0.333	0	3	-1
0	0	1	-0.333	0.667	0	2	4
0	0	0	0.333	-0.667	1	1	-2

Recall that the lower limit of the parameter, $\lambda \geq -1/2$, derives from $x_2^*(\lambda) \geq 0$, i.e., $x_2^*(-1/2) = 0$
A further decrease in λ requires that x_2 be removed from the basis (by a dual simplex pivot)

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-z	1	2	3	4	5	B	Δ
1	0	0	0.333	0.333	0	5	3
0	1	0	0.667	-0.333	0	3	-1
0	0	1	0.333	0.667	0	2	4
0	0	0	0.333	-0.667	1	1	-2

pivot here!

the dual simplex pivot yields

-z	1	2	3	4	5	B	Δ
1	0	1	0	1	0	7	7
0	1	2	0	1	0	7	7
0	0	-3	1	-2	0	-6	-12
0	0	1	0	0	1	3	2

The new basic solution is:

-z	1	3	5	B	Δ
1	0	0	0	7	7
0	1	0	0	7	7
0	0	1	0	-6	-12
0	0	0	1	3	2

$$\Rightarrow \begin{cases} z(\lambda) = -7 - 7\lambda \\ x_1^*(\lambda) = 7 + 7\lambda \\ x_3^*(\lambda) = -6 - 12\lambda \\ x_5^*(\lambda) = 3 + 2\lambda \end{cases}$$

To find the interval for which this basic solution is feasible (& therefore optimal), solve
 $x_1^*(\lambda) = 7 + 7\lambda \geq 0 \implies \lambda \geq -1$
 $x_3^*(\lambda) = -6 - 12\lambda \geq 0 \implies \lambda \leq -0.5$
 $x_5^*(\lambda) = 3 + 2\lambda \geq 0 \implies \lambda \geq -1.5$
 that is, $\lambda \in [-1.0, -0.5]$

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When λ decreases to -1.0 , $x_1^*(\lambda)$ decreases to 0 and must be removed from the basis to allow any further decrease in the parameter. We therefore attempt another dual simplex pivot:

-z	1	2	3	4	5	B	Δ
1	0	1	0	1	0	7	7
0	1	2	0	1	0	7	7
0	0	-3	1	-2	0	-6	-12
0	0	1	0	0	1	3	2

pivot row

Because there is no negative element in the pivot row, x_1 cannot be removed from the basis, and it is evident that $P(\lambda)$ is infeasible for $\lambda < -1.0$

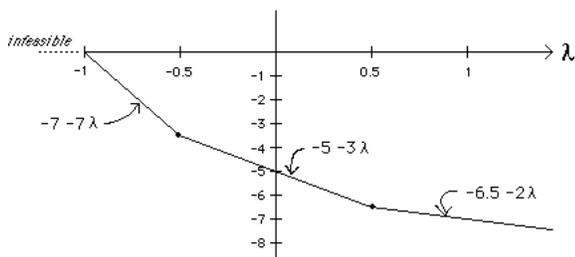
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Summary of Parametric Analysis:

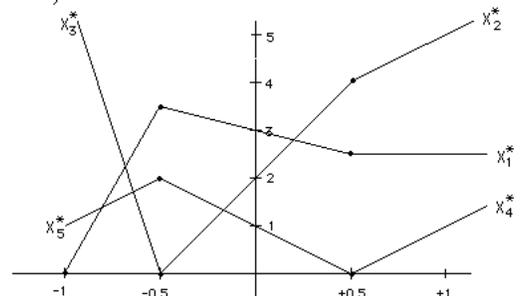
λ	$(-\infty, -1]$	$[-1, -0.5]$	$[-0.5, +0.5]$	$[+0.5, +\infty)$
x_1^*	<i>infeasible</i>	$7 + 7\lambda$	$3 - \lambda$	2.5
x_2^*		0	$2 + 4\lambda$	$3 + 2\lambda$
x_3^*		$-6 - 12\lambda$	0	0
x_4^*		0	0	$-1.5 + 3\lambda$
x_5^*		$3 + 2\lambda$	$1 - 2\lambda$	0

Plot of z vs λ :



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Plot of x_j^* vs. λ



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Example

Minimize $x_1 + x_2 + 7x_3 + 3x_4 + x_5 + 2x_6$
 subject to
 $x_1 + 2x_2 - x_3 - x_4 + x_5 + 2x_6 = 16 - \lambda$
 $x_2 - 3x_4 - x_5 + 3x_6 = 2 + \lambda$
 $-x_1 - 3x_3 + 3x_4 + x_5 = -4 + \lambda$
 $x_j \geq 0, j=1,2, \dots, 6$

Initial tableau:

-Z	1	2	3	4	5	6	B	Δ
1	1	1	7	3	1	2	0	0
0	1	2	-1	-1	1	2	16	-1
0	0	1	0	-3	-1	3	2	-1
0	-1	0	-3	3	1	0	-4	1

Optimal tableau ($\lambda = 0$)

-Z	1	2	3	4	5	6	B	Δ
1	0	0	7	3	0	2	-12	1.5
0	1	0	2	-1	0	-1	6	-1.5
0	0	1	-1	-1	0	2	4	0.5
0	0	0	-1	2	1	-1	2	-0.5

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-Z	1	2	3	4	5	6	B	Δ
1	0	0	7	3	0	2	-12	1.5
0	1	0	2	-1	0	-1	6	-1.5
0	0	1	-1	-1	0	2	4	0.5
0	0	0	-1	2	1	-1	2	-0.5

Parametric Analysis

Least Upper Bound (LUB): 4
 = $\text{Min}\{-6 - 2 \div -1.5, -0.5\} = \text{Min}\{4, 4\}$
 RHS at LUB is $-10 \ 0 \ 6 \ 0$
 Greatest Lower Bound (GLB): -8
 = $\text{Max}\{-4 \div 0.5\} = \text{Max}\{-8\}$
 RHS at GLB is $-16 \ 18 \ 0 \ 6$
 Range of parameters LAMBDA within which basis is feasible:
 $[-8, 4]$

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Dual Simplex pivot

-Z	1	2	3	4	5	6	B	Δ
1	0	0	7	3	0	2	-12	1.5
0	1	0	2	-1	0	-1	6	-1.5
0	0	1	-1	-1	0	2	4	0.5
0	0	0	-1	2	1	-1	2	-0.5

New tableau

-Z	1	2	3	4	5	6	B	Δ
1	2	0	11	1	0	0	0	-1.5
0	-1	0	-2	1	0	1	-6	1.5
0	2	1	3	-3	0	0	16	-2.5
0	-1	0	-3	3	1	0	-4	1

Dual Simplex Pivot

-Z	1	2	3	4	5	6	B	Δ
1	2	0	11	1	0	0	0	-1.5
0	-1	0	-2	1	0	1	-6	1.5
0	2	1	3	-3	0	0	16	-2.5
0	-1	0	-3	3	1	0	-4	1

Parametric Analysis

Least Upper Bound (LUB): 6.4
 = $\text{Min}\{-16 \div -2.5\} = \text{Min}\{6.4\}$
 RHS at LUB is $-16 \ 3.6 \ 0 \ 2.4$
 Greatest Lower Bound (GLB): 4
 = $\text{Max}\{6.4 \div 1.5, 1\} = \text{Max}\{4, 4\}$
 RHS at GLB is $-10 \ 0 \ 6 \ 0$
 Range of parameters LAMBDA within which basis is feasible:
 $[4, 6.4]$

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-Z	1	2	3	4	5	6	B	Δ
1	2.67	0.333	12	0	0	0	5.33	-2.33
0	-0.333	0.333	-1	0	0	1	-0.667	0.667
0	-0.667	-0.333	-1	1	0	0	-5.33	0.833
0	1	1	0	0	1	0	12	-1.5

Parametric Analysis

Least Upper Bound (LUB): 8
 = $\text{Min}\{-12 \div -1.5\} = \text{Min}\{8\}$
 RHS at LUB is $-21.3 \ 4.67 \ 1.33 \ 0$
 Greatest Lower Bound (GLB): 6.4
 = $\text{Max}\{0.667, 5.33 \div 0.667, 0.833\} = \text{Max}\{1, 6.4\}$
 RHS at GLB is $-16 \ 3.6 \ 0 \ 2.4$
 Range of parameters LAMBDA within which basis is feasible:
 $[6.4, 8]$

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A dual simplex pivot in row #4 is not possible:

-Z	1	2	3	4	5	6	B	Δ
1	2.67	0.333	12	0	0	0	5.33	-2.33
0	-0.333	0.333	-1	0	0	1	-0.667	0.667
0	-0.667	-0.333	-1	1	0	0	-5.33	0.833
0	1	1	0	0	1	0	12	-1.5

The LP is infeasible for $\lambda > 8$

Let's return to the optimal tableau for $\lambda = 0$:

-Z	1	2	3	4	5	6	B	Δ
1	0	0	7	3	0	2	-12	1.5
0	1	0	2	-1	0	-1	6	-1.5
0	0	1	-1	-1	0	2	4	0.5
0	0	0	-1	2	1	-1	2	-0.5

Parametric Analysis

Least Upper Bound (LUB): 4
 = $\text{Min}\{-6 - 2\lambda, -1.5 - 0.5\lambda\} = \text{Min}\{4, 4\}$
 RHS at LUB is -10 0 6 0
 Greatest Lower Bound (GLB): -8
 = $\text{Max}\{-4 + 0.5\lambda\} = \text{Max}\{-8\}$
 RHS at GLB is -16 18 0 6
 Range of parameters LAMBDA within which basis is feasible:
 [-8, 4]

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Dual Simplex Pivot

-Z	1	2	3	4	5	6	B	Δ
1	0	0	7	3	0	2	-12	1.5
0	1	0	2	-1	0	-1	6	-1.5
0	0	1	-1	-1	0	2	4	0.5
0	0	0	-1	2	1	-1	2	-0.5

New tableau

-Z	1	2	3	4	5	6	B	Δ
1	0	3	4	0	0	8	0	3
0	1	-1	3	0	0	-3	2	-2
0	0	-1	1	1	0	-2	-4	-0.5
0	0	2	-3	0	1	3	10	0.5

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-Z	1	2	3	4	5	6	B	Δ
1	0	3	4	0	0	8	0	3
0	1	-1	3	0	0	-3	2	-2
0	0	-1	1	1	0	-2	-4	-0.5
0	0	2	-3	0	1	3	10	0.5

Parametric Analysis

Least Upper Bound (LUB): -8
 = $\text{Min}\{-2 - 4\lambda, -2 - 0.5\lambda\} = \text{Min}\{-8\}$
 RHS at LUB is -16 18 0 6
 Greatest Lower Bound (GLB): -20
 = $\text{Max}\{-10 + 0.5\lambda\} = \text{Max}\{-20\}$
 RHS at GLB is -40 42 6 0
 Range of parameters LAMBDA within which basis is feasible:
 [-20, -8]

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Dual Simplex Pivot

-Z	1	2	3	4	5	6	B	Δ
1	0	3	4	0	0	8	0	3
0	1	-1	3	0	0	-3	2	-2
0	0	-1	1	1	0	-2	-4	-0.5
0	0	2	-3	0	1	3	10	0.5

New tableau

-Z	1	2	3	4	5	6	B	Δ
1	0	5.67	0	0	1.33	12	13.3	3.67
0	1	1	0	0	1	0	12	-1.5
0	0	-0.333	0	1	0.333	-1	-0.667	-0.333
0	0	-0.667	1	0	-0.333	-1	-3.33	-0.167

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-Z	1	2	3	4	5	6	B	Δ
1	0	5.67	0	0	1.33	12	13.3	3.67
0	1	1	0	0	1	0	12	-1.5
0	0	-0.333	0	1	0.333	-1	-0.667	-0.333
0	0	-0.667	1	0	-0.333	-1	-3.33	-0.167

Parametric Analysis

Least Upper Bound (LUB): -20
 = $\text{Min}\{-12 - 0.667\lambda, 3.33 - 1.5\lambda, -0.333 - 0.167\lambda\}$
 = $\text{Min}\{8 - 2\lambda\}$
 RHS at LUB is -60 42 6 0
 No Lower Bound
 Range of parameters LAMBDA within which basis is feasible:
 [-1.8E308, -20]

i.e., $-\infty$

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