

# Parametric Programming on the Right-Hand-Side

*Parametric Programming* is the analysis of the variation of the solution of an LP when some element (right-hand-side, objective, etc.) varies.

Consider the optimal value of the LP as a function of the right-hand-side vector, i.e.,

$$Z^*(b) = \left\{ \begin{array}{l} \text{Min } c \cdot x \\ \text{s.t. } Ax \geq b \\ x \geq 0 \end{array} \right\} \stackrel{\text{by duality theory}}{=} \left\{ \begin{array}{l} \text{Max } \pi \cdot b \\ \text{s.t. } \pi \cdot A \leq c \\ \pi \geq 0 \end{array} \right\}$$

The function  $z^*$  is "evaluated" for some particular right-hand-side  $b'$  by solving the LP (either the primal or the dual).

So we can evaluate  $z^*(b')$  by solving the LP

$$z^*(b') = \left\{ \begin{array}{l} \text{Max } \pi \cdot b' \\ \text{s.t. } \pi \cdot A \leq c \\ \pi \geq 0 \end{array} \right\}$$

Notice that the feasible region of the dual LP is the same for every argument  $b'$ .

We know that a basic solution is optimal for an LP problem, and that there are a finite (but possibly very large!) number of such basic solutions.

Suppose that we were to number the basic feasible solutions of the dual LP:

$$\{ \pi^1, \pi^2, \pi^3, \dots, \pi^K \}$$

where each  $\pi^k = c_{B_k} (A^{B_k})^{-1}$  for some dual-feasible basis  $B_k$

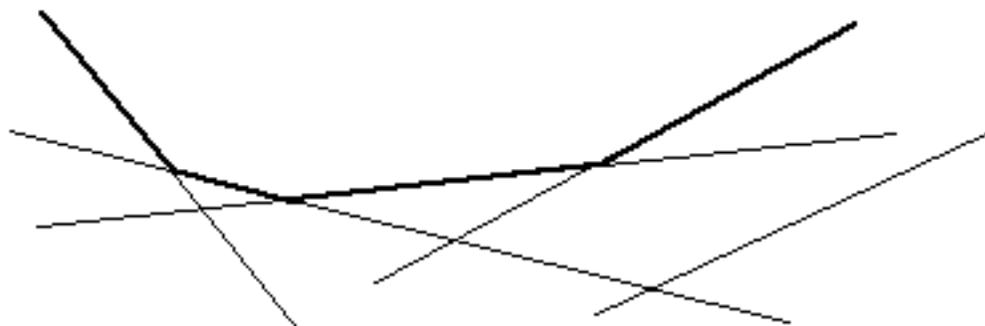
( $K$  is a finite number, no greater than  $\binom{n+m}{m}$ .)

In theory, then, we could evaluate  $z^*(b)$  by enumerating the  $K$  basic feasible solutions of the dual, evaluating the dual objective  $\pi^k b$  at each, and selecting the maximum such value.

Therefore,  $z^*(b) = \text{Maximum}_{k=1,2,\dots,K} \{ \pi^k b \}$

That is,  $z^*(b)$  is the maximum of a family of linear functions,  $\pi^k b$ ,  $k=1,2,\dots,K$

which is a piecewise linear convex function!



*Let us restrict our analysis of the function  $z^*(b)$  to a study of its behavior along a line (rather than everywhere in  $m$ -dimensional space!)*

*That is, we assume an initial right-hand-side vector  $(b)$  is given, and a direction  $(d)$ , and study the behavior of  $z^*(b + \lambda d)$ , considered as a function of the scalar parameter  $\lambda$ .*

Consider the solution of the LP

$$P_\lambda : \quad z^*(\lambda) = \text{minimum } c x \\ \text{s.t. } Ax = b + \lambda d \\ x \geq 0$$

where  $d$  is an  $m$ -vector and  $\lambda$  is a scalar.

$$z^*(\lambda) = \text{maximum}_{k=1,2,\dots,K} \{ \pi^k(b + \lambda d) \} \\ = \text{maximum}_{k=1,2,\dots,K} \left\{ \underbrace{\pi^k b}_{\text{intercept}} + \underbrace{(\pi^k d)}_{\text{slope}} \lambda \right\} \leftarrow \begin{array}{l} \text{linear functions} \\ \text{of } \lambda \end{array}$$

**Example**

$$z^*(\lambda) = \text{minimum } -x_1 - x_2$$

$$\text{subject to } \begin{cases} 2x_1 + x_2 \leq 8 + 2\lambda \\ x_1 + 2x_2 \leq 7 + 7\lambda \\ x_2 \leq 3 + 2\lambda \\ x_1 \geq 0, x_2 \geq 0 \end{cases}$$

$$= \text{maximum } \{ (8\pi_1 + 7\pi_2 + 2\pi_3) + (2\pi_1 + 7\pi_2 + 2\pi_3)\lambda \}$$

$$\text{s.t. } \begin{cases} 2\pi_1 + \pi_2 \leq -1 \\ \pi_1 + 2\pi_2 + \pi_3 \leq -1 \\ \pi_1 \leq 0, \pi_2 \leq 0, \pi_3 \leq 0 \end{cases}$$

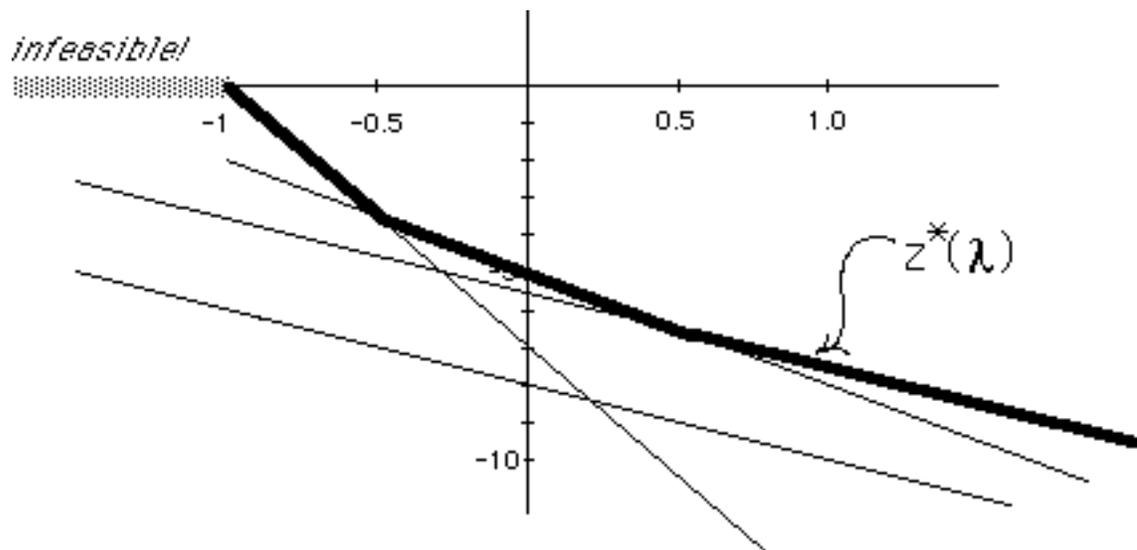
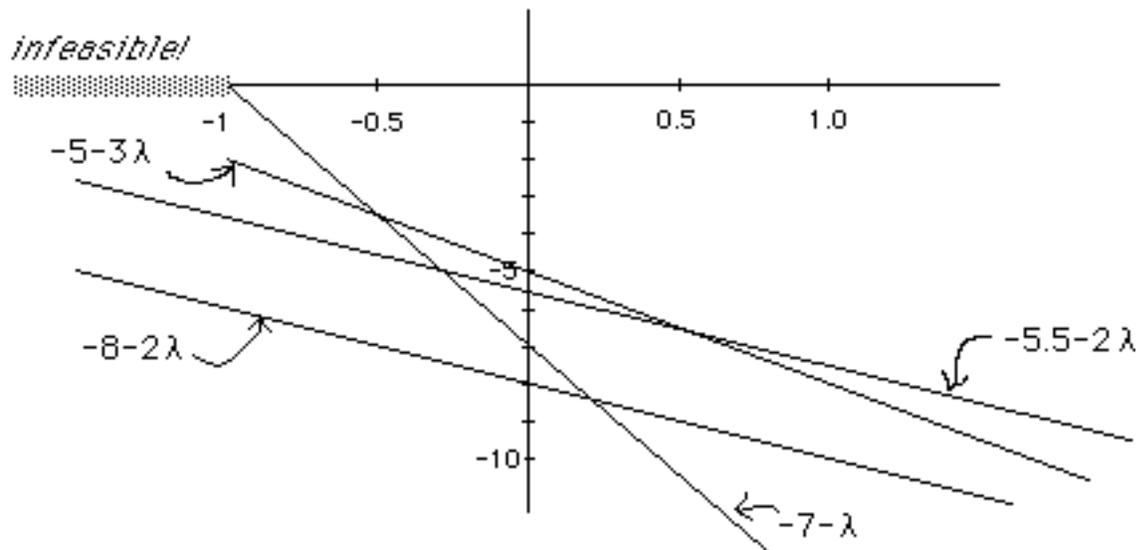
*nonpositive because of direction of  $\leq$  in the primal!*

$\pi_1$	$\pi_2$	$\pi_3$	<i>intercept</i> $\pi^k \mathbf{b}$	<i>slope</i> $\pi^k \mathbf{d}$	<i>Basis</i>
0	0	0	0	0	3 4 5
-0.5	0	0	-4	-1	1 4 5
-1	0	0	-8	-2	2 4 5
-0.333	-0.333	0	-5	-3	1 2 5
0	-1	0	-7	-7	1 3 5
0	-0.5	0	-3.5	-3.5	2 3 5
0	-1	1	-4	-5	1 2 3
-0.5	0	-0.5	-5.5	-2	1 2 4
0	0	-1	-3	-2	2 3 4

**feasible**

(columns #1,3,4 are dependent & do not form a basis!)

Of the nine basic solutions, four are dual feasible. Therefore,  $z^*(\lambda)$  is the maximum of four linear functions:



*In this example, with only nine basic dual solutions, it was possible to enumerate all of them, test each for feasibility, and then maximize the corresponding linear functions*

$$\pi^k b + (\pi^k d) \lambda$$

*However, for most problems, the number of basic dual solutions is astronomical and enumerating them is practically impossible.*

*(Usually, only a few of these basic dual solutions actually determine  $z^*$ .)*

Let's consider again the parametric LP  $P_\lambda$ :

$$z^*(\lambda) = \begin{array}{l} \text{minimum } -x_1 - x_2 \\ \text{subject to } \left\{ \begin{array}{l} 2x_1 + x_2 \leq 8 + 2\lambda \\ x_1 + 2x_2 \leq 7 + 7\lambda \\ x_2 \leq 3 + 2\lambda \\ x_1 \geq 0, x_2 \geq 0 \end{array} \right. \end{array}$$

Determine the optimal value function  $z^*(\lambda)$  as well as  $x_1^*(\lambda)$  and  $x_2^*(\lambda)$  [i.e., the optimal solution] for all values of  $\lambda \in (-\infty, +\infty)$

The initial tableau:

$-z$	<b>1</b>	<b>2</b>	<b>3</b>	<b>4</b>	<b>5</b>	<b>B</b>	$\Delta$
<b>1</b>	<b>-1</b>	<b>-1</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>	<b>0</b>
<b>0</b>	<b>2</b>	<b>1</b>	<b>1</b>	<b>0</b>	<b>0</b>	<b>8</b>	<b>2</b>
<b>0</b>	<b>1</b>	<b>2</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>7</b>	<b>7</b>
<b>0</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>0</b>	<b>1</b>	<b>3</b>	<b>2</b>

*change column*  $\swarrow$

*RHS is*  
 $B + \lambda \Delta$

$\underbrace{\hspace{10em}}$   
*slack variables*

Let's start with  $\lambda = 0$ , and investigate the LP as  $\lambda$  increases.

The optimal tableau for  $\lambda = 0$ :

$-z$	<b>1</b>	<b>2</b>	<b>3</b>	<b>4</b>	<b>5</b>	<b>B</b>	$\Delta$
<i>(min.)</i> <b>1</b>	<b>0</b>	<b>0</b>	<b>0.333</b>	<b>0.333</b>	<b>0</b>	<b>5</b>	<b>3</b>
<b>0</b>	<b>1</b>	<b>0</b>	<b>0.667</b>	<b>-0.333</b>	<b>0</b>	<b>3</b>	<b>-1</b>
<b>0</b>	<b>0</b>	<b>1</b>	<b>-0.333</b>	<b>0.667</b>	<b>0</b>	<b>2</b>	<b>4</b>
<b>0</b>	<b>0</b>	<b>0</b>	<b>0.333</b>	<b>-0.667</b>	<b>1</b>	<b>1</b>	<b>-2</b>

$\uparrow$   
*(this column updated during each pivot)*

The optimal solution of  $P(0)$  is  $z^*(0) = -5$ ,  
 at  $x_1^*(0) = 3$ ,  $x_2^*(0) = 2$ ,  $x_5^*(0) = 1$ ,  
 $x_3^*(0) = x_4^*(0) = 0$

Expressed as functions of  $\lambda$ , the basic solution is:

$$\begin{array}{cccc|c|c}
 -z & 1 & 2 & 5 & \mathbf{B} & \Delta \\
 \hline
 1 & 0 & 0 & 0 & 5 & 3 \\
 0 & 1 & 0 & 0 & 3 & -1 \\
 0 & 0 & 1 & 0 & 2 & 4 \\
 0 & 0 & 0 & 1 & 1 & -2
 \end{array}
 \Rightarrow
 \begin{cases}
 z(\lambda) = -5 + 3\lambda \\
 x_1^*(\lambda) = 3 - \lambda \\
 x_2^*(\lambda) = 2 + 4\lambda \\
 x_5^*(\lambda) = 1 - 2\lambda \\
 x_3^*(\lambda) = x_4^*(\lambda) = 0
 \end{cases}$$

Note that these are linear functions of  $\lambda$ !

The optimality criterion (reduced cost  $\geq 0$ ) is independent of the parameter  $\lambda$ , and so the current basis remains optimal so long as the basic variables

$$\begin{aligned}
 x_1^*(\lambda) &= 3 - \lambda \\
 x_2^*(\lambda) &= 2 + 4\lambda \\
 x_5^*(\lambda) &= 1 - 2\lambda
 \end{aligned}$$

remain feasible, i.e., nonnegative.

*For what values of  $\lambda$  is  $x^*(\lambda) \geq 0$ ?*

We solve the inequalities

$$x_1^*(\lambda) = 3 - \lambda \geq 0$$

$$x_2^*(\lambda) = 2 + 4\lambda \geq 0$$

$$x_5^*(\lambda) = 1 - 2\lambda \geq 0$$

for  $\lambda$ :

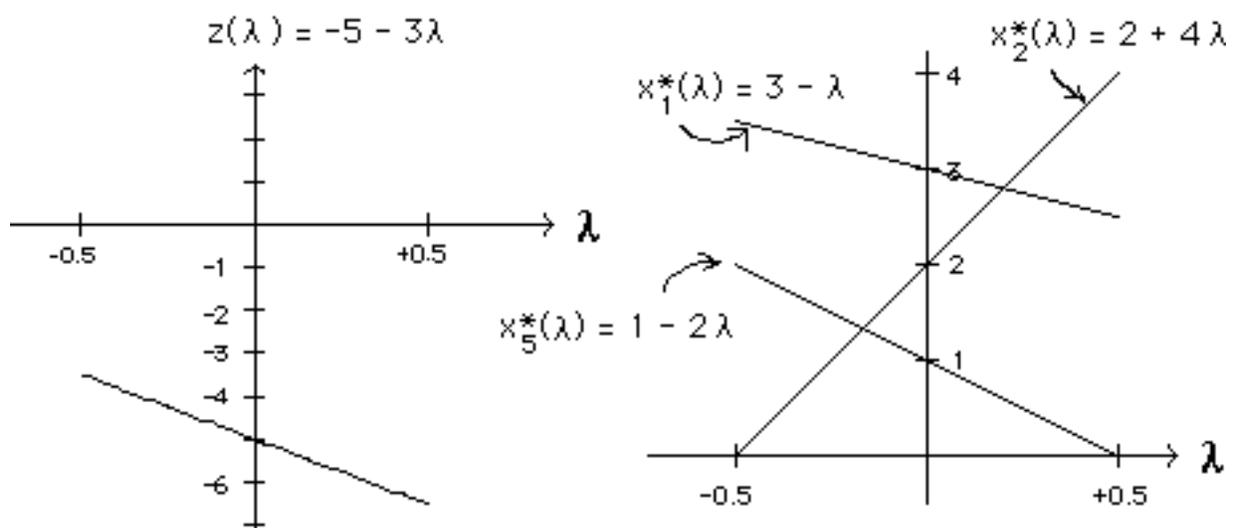
$$3 - \lambda \geq 0 \Rightarrow \lambda \leq 3$$

$$2 + 4\lambda \geq 0 \Rightarrow \lambda \geq -\frac{1}{2}$$

$$1 - 2\lambda \geq 0 \Rightarrow \lambda \leq +\frac{1}{2}$$

That is, the basic solution is feasible for all  $\lambda \in [-\frac{1}{2}, +\frac{1}{2}]$  (and, in particular, for  $\lambda = 0$ ).

Plot of  $z(\lambda)$  &  $x_j^*(\lambda)$  over the interval  $[-0.5, +0.5]$ :



Let us now increase  $\lambda$  from its initial value (0) to the upper limit for which the basis is feasible, i.e.,  $+1/2$ . The basic solution then becomes

$$x_1^*(+1/2) = 3 - 1/2 = 2.5$$

$$x_2^*(+1/2) = 2 + 4(1/2) = 4.0$$

$$x_5^*(+1/2) = 1 - 2(1/2) = 0$$

Any further increase in  $\lambda$  would result in infeasible (i.e., negative) values for  $x_5^*$  !

In order to increase the parameter  $\lambda$  further,  $x_5$  must leave the basis, since it would otherwise become negative. In order to remove  $x_5$  from the basis, we perform a DUAL SIMPLEX pivot:

$-z$	1	2	3	4	5	B	$\Delta$
1	0	0	0.333	0.333	0	5	3
0	1	0	0.667	-0.333	0	3	-1
0	0	1	-0.333	0.667	0	2	4
0	0	0	0.333	-0.667	1	1	-2

*pivot in this row!*

$\leftarrow \begin{cases} = 0 \text{ when} \\ \lambda = 1/2 \end{cases}$

-Z	1	2	3	4	5	B	Δ
1	0	0	0.333	0.333	0	5	3
0	1	0	0.667	-0.333	0	3	-1
0	0	1	-0.333	0.667	0	2	4
0	0	0	0.333	-0.667	1	1	-2

*pivot* ↗

-Z	1	2	3	4	5	B	Δ
1	0	0	0.5	0	0.5	5.5	2
0	1	0	0.5	0	-0.5	2.5	0
0	0	1	0	0	1	3	2
0	0	0	-0.5	1	-1.5	-1.5	3

For this basis, the basic solution is

-Z	1	2	...	4	...	B	Δ
1	0	0		0		5.5	2
0	1	0		0		2.5	0
0	0	1		0		3	2
0	0	0		1		-1.5	3

$$\Rightarrow \begin{cases} z(\lambda) = -5.5 - 2\lambda \\ x_1^*(\lambda) = 2.5 \\ x_2^*(\lambda) = 3 + 2\lambda \\ x_4^*(\lambda) = -1.5 + 3\lambda \end{cases}$$

Notice that as  $\lambda$  increases, no basic variable decreases. Since the optimality criterion (reduced costs  $\geq 0$ ) does not depend upon  $\lambda$ , this basis is optimal for all  $\lambda \geq 0.5$

That is, if we solve the system of inequalities

$$x_1^*(\lambda) = 2.5 \geq 0 \implies (\text{no restriction on } \lambda)$$

$$x_2^*(\lambda) = 3 + 2\lambda \geq 0 \implies \lambda \geq -1.5$$

$$x_4^*(\lambda) = -1.5 + 3\lambda \geq 0 \implies \lambda \geq 0.5$$

we see that it is satisfied for  $\lambda \in [0.5, +\infty)$

Let us now investigate  $P(\lambda)$  for  $\lambda < -0.5$   
 Consider the tableau which was optimal for  
 $-0.5 \leq \lambda \leq +0.5$

-z	1	2	3	4	5	B	$\Delta$
1	0	0	0.333	0.333	0	5	3
0	1	0	0.667	-0.333	0	3	-1
0	0	1	-0.333	0.667	0	2	4
0	0	0	0.333	-0.667	1	1	-2

Recall that the lower limit of the parameter,  
 $\lambda \geq -1/2$ , derives from  $x_2^*(\lambda) \geq 0$ , i.e.,  $x_2^*(-1/2) = 0$

A further decrease in  $\lambda$  requires that  $x_2$  be removed  
 from the basis (by a dual simplex pivot)

-Z	1	2	3	4	5	B	Δ
1	0	0	0.333	0.333	0	5	3
0	1	0	0.667	-0.333	0	3	-1
0	0	1	-0.333	0.667	0	2	4 ← pivot row
0	0	0	0.333	-0.667	1	1	-2

*pivot here!*

*the dual simplex pivot yields*

⇒

-Z	1	2	3	4	5	B	Δ
1	0	1	0	1	0	7	7
0	1	2	0	1	0	7	7
0	0	-3	1	-2	0	-6	-12
0	0	1	0	0	1	3	2

The new basic solution is:

-Z	1	3	5	B	Δ
1	0	0	0	7	7
0	1	0	0	7	7
0	0	1	0	-6	-12
0	0	0	1	3	2

$$\Rightarrow \begin{cases} z(\lambda) = -7 - 7\lambda \\ x_1^*(\lambda) = 7 + 7\lambda \\ x_3^*(\lambda) = -6 - 12\lambda \\ x_5^*(\lambda) = 3 + 2\lambda \end{cases}$$

To find the interval for which this basic solution is feasible (& therefore optimal), solve

$$\begin{cases} x_1^*(\lambda) = 7 + 7\lambda \geq 0 \\ x_3^*(\lambda) = -6 - 12\lambda \geq 0 \\ x_5^*(\lambda) = 3 + 2\lambda \geq 0 \end{cases} \Rightarrow \begin{cases} \lambda \geq -1 \\ \lambda \leq -0.5 \\ \lambda \geq -1.5 \end{cases} \quad \text{that is, } \lambda \in [-1.0, -0.5]$$

When  $\lambda$  decreases to  $-1.0$ ,  $x_1^*(\lambda)$  decreases to  $0$  and must be removed from the basis to allow any further decrease in the parameter. We therefore attempt another dual simplex pivot:

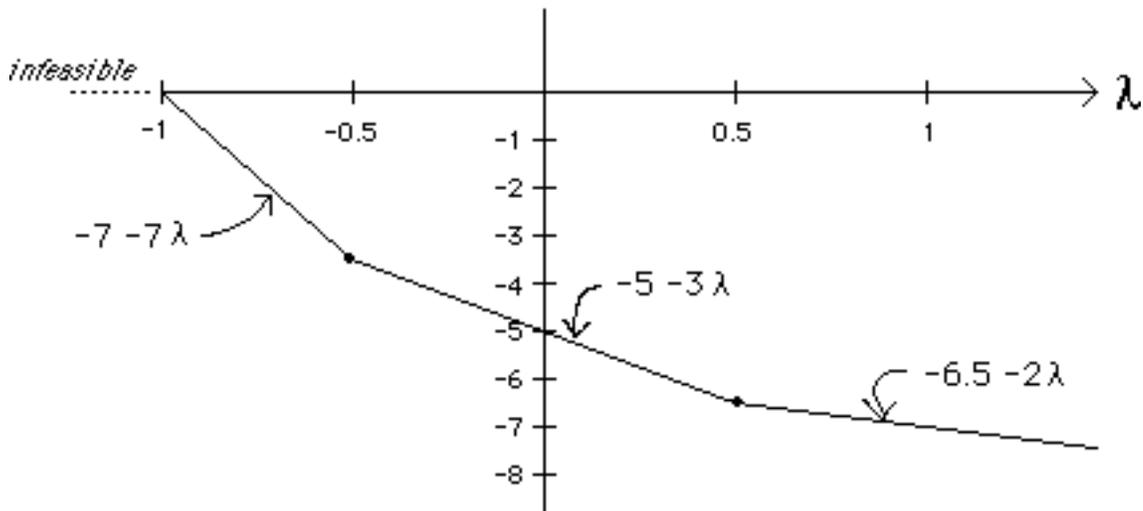
$-Z$	<b>1</b>	<b>2</b>	<b>3</b>	<b>4</b>	<b>5</b>	<b>B</b>	$\Delta$
<b>1</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>7</b>	<b>7</b>
<b>0</b>	<b>1</b>	<b>2</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>7</b>	<b>7</b> ← <i>pivot row</i>
<b>0</b>	<b>0</b>	<b>-3</b>	<b>1</b>	<b>-2</b>	<b>0</b>	<b>-6</b>	<b>-12</b>
<b>0</b>	<b>0</b>	<b>1</b>	<b>0</b>	<b>0</b>	<b>1</b>	<b>3</b>	<b>2</b>

Because there is no negative element in the pivot row,  $x_1$  cannot be removed from the basis, and it is evident that  $P(\lambda)$  is infeasible for  $\lambda < -1.0$

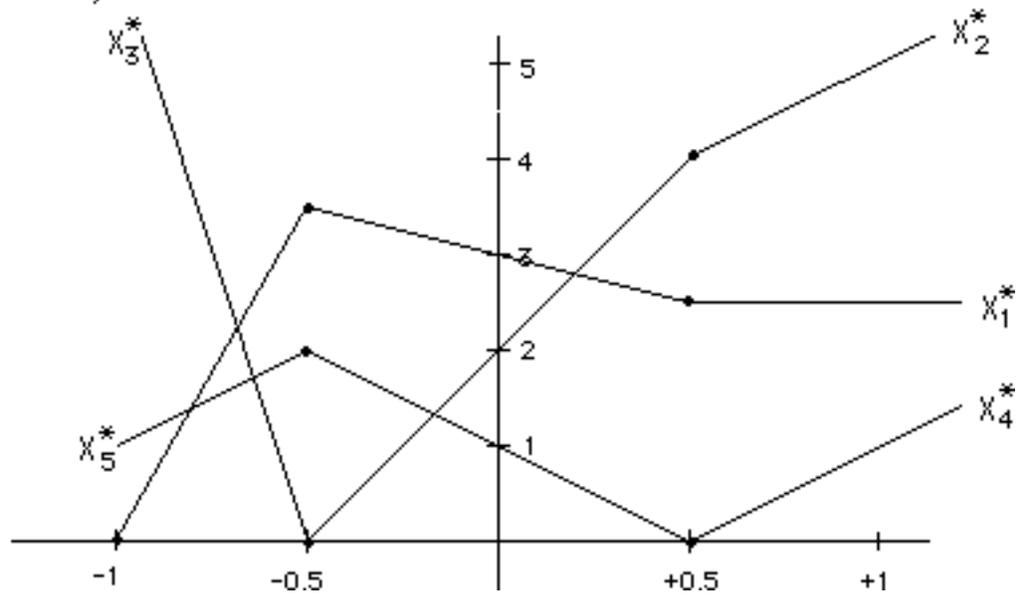
### Summary of Parametric Analysis:

$\lambda$	$(-\infty, -1]$	$[-1, -0.5]$	$[-0.5, +0.5]$	$[+0.5, +\infty)$
$x_1^*$	<i>infeasible</i>	$7 + 7\lambda$	$3 - \lambda$	$2.5$
$x_2^*$		$0$	$2 + 4\lambda$	$3 + 2\lambda$
$x_3^*$		$-6 - 12\lambda$	$0$	$0$
$x_4^*$		$0$	$0$	$-1.5 + 3\lambda$
$x_5^*$		$3 + 2\lambda$	$1 - 2\lambda$	$0$

*Plot of  $z$  vs  $\lambda$  :*



*Plot of  $x_j^*$  vs.  $\lambda$*



## Example

Minimize  $x_1 + x_2 + 7x_3 + 3x_4 + x_5 + 2x_6$   
 subject to

$$\begin{aligned} x_1 + 2x_2 - x_3 - x_4 + x_5 + 2x_6 &= 16 - \lambda \\ x_2 - 3x_4 - x_5 + 3x_6 &= 2 + \lambda \\ -x_1 - 3x_3 + 3x_4 + x_5 &= -4 + \lambda \\ x_j &\geq 0, j=1,2, \dots, 6 \end{aligned}$$

Initial tableau:

-Z	1	2	3	4	5	6	B	Δ
1	1	1	7	3	1	2	0	0
0	1	2	-1	-1	1	2	16	-1
0	0	1	0	-3	-1	3	2	1
0	-1	0	-3	3	1	0	-4	1

Optimal tableau (  $\lambda = 0$  )

-Z	1	2	3	4	5	6	B	Δ
1	0	0	7	3	0	2	-12	1.5
0	1	0	2	-1	0	-1	6	-1.5
0	0	1	-1	-1	0	2	4	0.5
0	0	0	-1	2	1	-1	2	-0.5

-Z	1	2	3	4	5	6	B	$\Delta$
1	0	0	7	3	0	2	-12	1.5
0	1	0	2	-1	0	-1	6	-1.5
0	0	1	-1	-1	0	2	4	0.5
0	0	0	-1	2	1	-1	2	-0.5

**Parametric Analysis**

Least Upper Bound (LUB): 4  
 =  $\text{Min}\{-6 \cdot -2 \div -1.5 \cdot -0.5\} = \text{Min}\{4 \cdot 4\}$   
 RHS at LUB is -10 0 6 0  
 Greatest Lower Bound (GLB): -8  
 =  $\text{Max}\{-4 \div 0.5\} = \text{Max}\{-8\}$   
 RHS at GLB is -16 18 0 6  
 Range of parameters LAMBDA within which basis is feasible:  
 [ -8 , 4 ]

**Dual Simplex pivot**

-Z	1	2	3	4	5	6	B	$\Delta$
1	0	0	7	3	0	2	-12	1.5
0	1	0	2	-1	0	-1	6	-1.5
0	0	1	-1	-1	0	2	4	0.5
0	0	0	-1	2	1	-1	2	-0.5

**New tableau**

-Z	1	2	3	4	5	6	B	$\Delta$
1	2	0	11	1	0	0	0	-1.5
0	-1	0	-2	1	0	1	-6	1.5
0	2	1	3	-3	0	0	16	-2.5
0	-1	0	-3	3	1	0	-4	1

-Z	1	2	3	4	5	6	B	Δ
1	2	0	11	1	0	0	0	-1.5
0	-1	0	-2	1	0	1	-6	1.5
0	2	1	3	-3	0	0	16	-2.5
0	-1	0	-3	3	1	0	-4	1

Parametric Analysis

Least Upper Bound (LUB): 6.4  
 =  $\text{Min}\{-16 \div -2.5\} = \text{Min}\{6.4\}$   
 RHS at LUB is  $-16 \ 3.6 \ 0 \ 2.4$   
 Greatest Lower Bound (GLB): 4  
 =  $\text{Max}\{6.4 \div 1.5 \ 1\} = \text{Max}\{4 \ 4\}$   
 RHS at GLB is  $-10 \ 0 \ 6 \ 0$   
 Range of parameters LAMBDA within which basis is feasible:  
 [ 4 , 6.4 ]

Dual Simplex Pivot

-Z	1	2	3	4	5	6	B	Δ
1	2	0	11	1	0	0	0	-1.5
0	-1	0	-2	1	0	1	-6	1.5
0	2	1	3	-3	0	0	16	-2.5
0	-1	0	-3	3	1	0	-4	1

←

-Z	1	2	3	4	5	6	B	Δ
1	2.67	0.333	12	0	0	0	5.33	-2.33
0	-0.333	0.333	-1	0	0	1	-0.667	0.667
0	-0.667	-0.333	-1	1	0	0	-5.33	0.833
0	1	1	0	0	1	0	12	-1.5

-Z	1	2	3	4	5	6	B	Δ
1	2.67	0.333	12	0	0	0	5.33	-2.33
0	-0.333	0.333	-1	0	0	1	-0.667	0.667
0	-0.667	-0.333	-1	1	0	0	-5.33	0.833
0	1	1	0	0	1	0	12	-1.5

Parametric Analysis

Least Upper Bound (LUB): 8  
 =  $\text{Min}\{-12 \div -1.5\} = \text{Min}\{8\}$   
 RHS at LUB is  $-21.3 \ 4.67 \ 1.33 \ 0$   
 Greatest Lower Bound (GLB): 6.4  
 =  $\text{Max}\{0.667 \ 5.33 \div 0.667 \ 0.833\} = \text{Max}\{1 \ 6.4\}$   
 RHS at GLB is  $-16 \ 3.6 \ 0 \ 2.4$   
 Range of parameters LAMBDA within which basis is feasible:  
 [ 6.4 , 8 ]

A dual simplex pivot in row #4 is not possible:

-Z	1	2	3	4	5	6	B	Δ
1	2.67	0.333	12	0	0	0	5.33	-2.33
0	-0.333	0.333	-1	0	0	1	-0.667	0.667
0	-0.667	-0.333	-1	1	0	0	-5.33	0.833
0	1	1	0	0	1	0	12	-1.5

The LP is infeasible for  $\lambda > 8$

Let's return to the optimal tableau for  $\lambda = 0$ :

-Z	1	2	3	4	5	6	B	$\Delta$
1	0	0	7	3	0	2	-12	1.5
0	1	0	2	-1	0	-1	6	-1.5
0	0	1	-1	-1	0	2	4	0.5
0	0	0	-1	2	1	-1	2	-0.5

**Parametric Analysis**

Least Upper Bound (LUB): 4  
 =  $\text{Min}\{-6 - 2 \div -1.5 \quad -0.5\} = \text{Min}\{4 \quad 4\}$   
 RHS at LUB is -10 0 6 0  
 Greatest Lower Bound (GLB): -8  
 =  $\text{Max}\{-4 \div 0.5\} = \text{Max}\{-8\}$   
 RHS at GLB is -16 18 0 6  
 Range of parameters LAMBDA within which basis is feasible:  
 [ -8 , 4 ]

**Dual Simplex Pivot**

-Z	1	2	3	4	5	6	B	$\Delta$
1	0	0	7	3	0	2	-12	1.5
0	1	0	2	-1	0	-1	6	-1.5
0	0	1	-1	-1	0	2	4	0.5
0	0	0	-1	2	1	-1	2	-0.5

←

**New tableau**

-Z	1	2	3	4	5	6	B	$\Delta$
1	0	3	4	0	0	8	0	3
0	1	-1	3	0	0	-3	2	-2
0	0	-1	1	1	0	-2	-4	-0.5
0	0	2	-3	0	1	3	10	0.5

-Z	1	2	3	4	5	6	B	Δ
1	0	3	4	0	0	8	0	3
0	1	-1	3	0	0	-3	2	-2
0	0	-1	1	1	0	-2	-4	-0.5
0	0	2	-3	0	1	3	10	0.5

**Parametric Analysis**

Least Upper Bound (LUB): -8  
 =  $\text{Min}\{-2 \cdot 4 \div -2 \cdot -0.5\} = \text{Min}\{1 \cdot -8\}$   
 RHS at LUB is -16 18 0 6  
 Greatest Lower Bound (GLB): -20  
 =  $\text{Max}\{-10 \div 0.5\} = \text{Max}\{-20\}$   
 RHS at GLB is -40 42 6 0  
 Range of parameters LAMBDA within which basis is feasible:  
 [ -20 , -8 ]

**Dual Simplex Pivot**

-Z	1	2	3	4	5	6	B	Δ
1	0	3	4	0	0	8	0	3
0	1	-1	3	0	0	-3	2	-2
0	0	-1	1	1	0	-2	-4	-0.5
0	0	2	-3	0	1	3	10	0.5

←

**New tableau**

-Z	1	2	3	4	5	6	B	Δ
1	0	5.67	0	0	1.33	12	13.3	3.67
0	1	1	0	0	1	0	12	-1.5
0	0	-0.333	0	1	0.333	-1	-0.667	-0.333
0	0	-0.667	1	0	-0.333	-1	-3.33	-0.167

-Z	1	2	3	4	5	6	B	$\Delta$
1	0	5.67	0	0	1.33	12	13.3	3.67
0	1	1	0	0	1	0	12	-1.5
0	0	-0.333	0	1	0.333	-1	-0.667	-0.333
0	0	-0.667	1	0	-0.333	-1	-3.33	-0.167

### Parametric Analysis

Least Upper Bound (LUB): -20  
 =  $\text{Min}\{-12 \ 0.667 \ 3.33 \div \ -1.5 \ -0.333 \ -0.167\}$   
 =  $\text{Min}\{ \ 8 \ -2 \ -20\}$   
 RHS at LUB is -60 42 6 0

No Lower Bound

Range of parameters LAMBDA within which basis is feasible:  
 [ -1.8E308 , -20 ]

i.e.,  $-\infty$  

