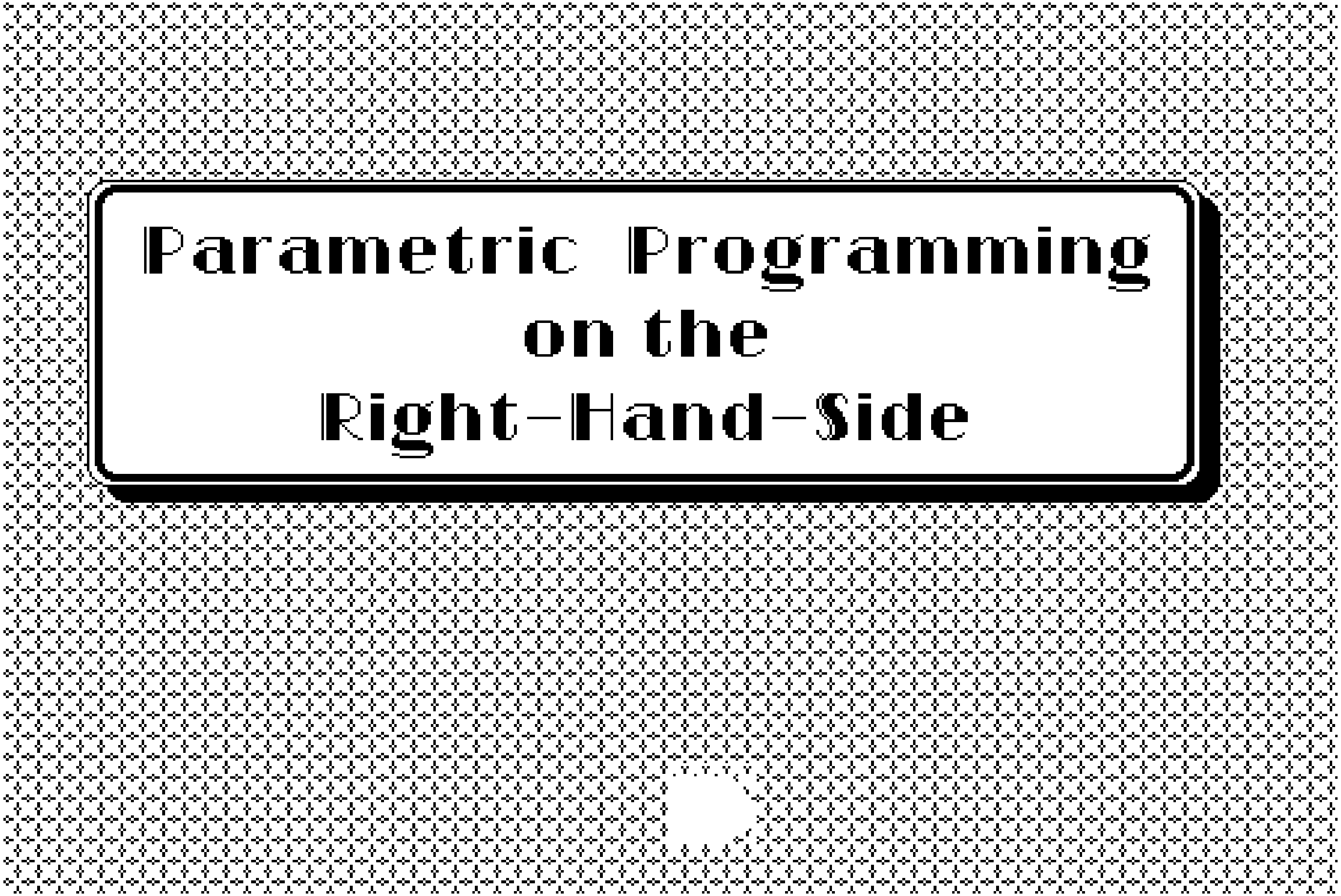


Parametric Programming on the Right-Hand-Side

Parametric Programming is the analysis of the variation of the solution of an LP when some element (right-hand-side, objective, etc.) varies.

Consider the optimal value of the $\mathbb{L}\,\mathbb{P}$ as a function of the right-hand-side vector, i.e.,

$$Z^*(b) = \left\{ \begin{array}{ll} \text{Min} & c\ x \\ \text{s.t.} & Ax \geq b \\ & x \geq 0 \end{array} \right\} \underset{\text{by duality theory}}{=\!=} \left\{ \begin{array}{ll} \text{Max} & \pi\, b \\ \text{s.t.} & \pi A \leq c \\ & \pi \geq 0 \end{array} \right\}$$

The function z^* is "evaluated" for some particular right-hand-side b' by solving the LP (either the primal or the dual).

So we can evaluate $z^*(b')$ by solving the LP

$$z^*(b') = \begin{cases} \text{Max} & \pi b' \\ \text{s.t.} & \pi A \le c \\ & \pi \ge 0 \end{cases}$$

Notice that the feasible region of the dual LP is the same for every argument b'.

We know that a basic solution is optimal for an LP problem, and that there are a finite (but possibly very large!) number of such basic solutions.

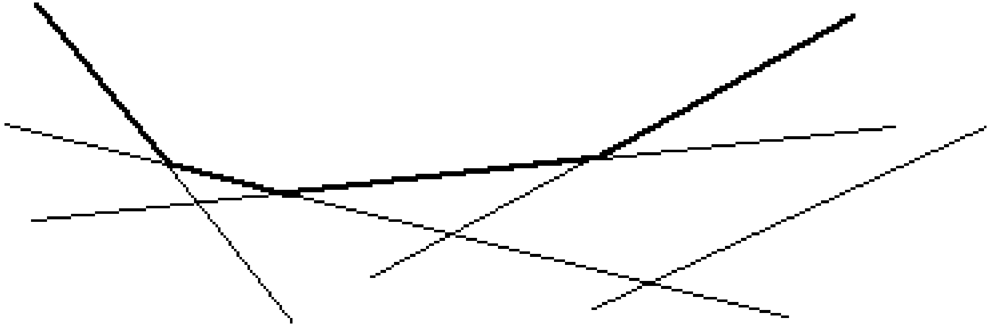
Suppose that we were to number the basic feasible solutions of the dual LP:
$\{ \pi^1, \pi^2, \pi^3, \ldots \pi^K \}$
where each $\pi^k = c_{B_k}(A^{B_k})^{-1}$ for some dual-feasible basis B_k
(K is a finite number, no greater than $\binom{n+m}{m}$.)

In theory, then, we could evaluate $z^*(b)$ by enumerating the K basic feasible solutions of the dual, evaluating the dual objective $\pi^k b$ at each, and selecting the maximum such value.

Therefore, $z^*(b) = \underset{k=1,2,\ldots K}{\text{Maximum}} \left\{ \pi^k b \right\}$

That is, $z^*(b)$ is the maximum of a family of linear functions, $\pi^k b$, $k=1,2,\ldots K$

which is a piecewise linear convex function!



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Let us restrict our analysis of the function $z^*(b)$ to a study of its behavior along a line (rather than everywhere in m-dimensional space!)

That is, we assume an initial right-hand-side vector (b) is given, and a direction (d), and study the behavior of $z^*(b+\lambda d)$, considered as a function of the scalar parameter λ.

Consider the solution of the LP

$$P_\lambda : \quad z^*(\lambda) = \text{minimum } cx \quad \text{s.t. } Ax = b + \lambda d, \quad x \geq 0$$

where d is an m-vector and λ is a scalar.

$$z^*(\lambda) = \underset{k=1,2,\ldots K}{\text{maximum}} \left\{ \pi^k (b + \lambda d) \right\}$$

$$= \underset{k=1,2,\ldots K}{\text{maximum}} \left\{ \underbrace{\pi^k b}_{\text{intercept}} + \underbrace{(\pi^k d)}_{\text{slope}} \lambda \right\} \leftarrow \text{linear functions of } \lambda$$

Example

$$z^*(\lambda) = \text{minimum } -x_1 - x_2$$

$$\text{subject to } \begin{cases} 2x_1 + x_2 \le 8 + 2\lambda \\ x_1 + 2x_2 \le 7 + 7\lambda \\ x_2 \le 3 + 2\lambda \\ x_1 \ge 0,\ x_2 \ge 0 \end{cases}$$

$$= \text{maximum } \{ (8\pi_1+7\pi_2+2\pi_3) + (2\pi_1+7\pi_2+2\pi_3)\lambda \}$$

$$\text{s.t. } \begin{array}{l} 2\pi_1 + \pi_2 \le -1 \\ \pi_1 + 2\pi_2 + \pi_3 \le -1 \\ \pi_1 \le 0,\ \pi_2 \le 0,\ \pi_3 \le 0 \end{array}$$

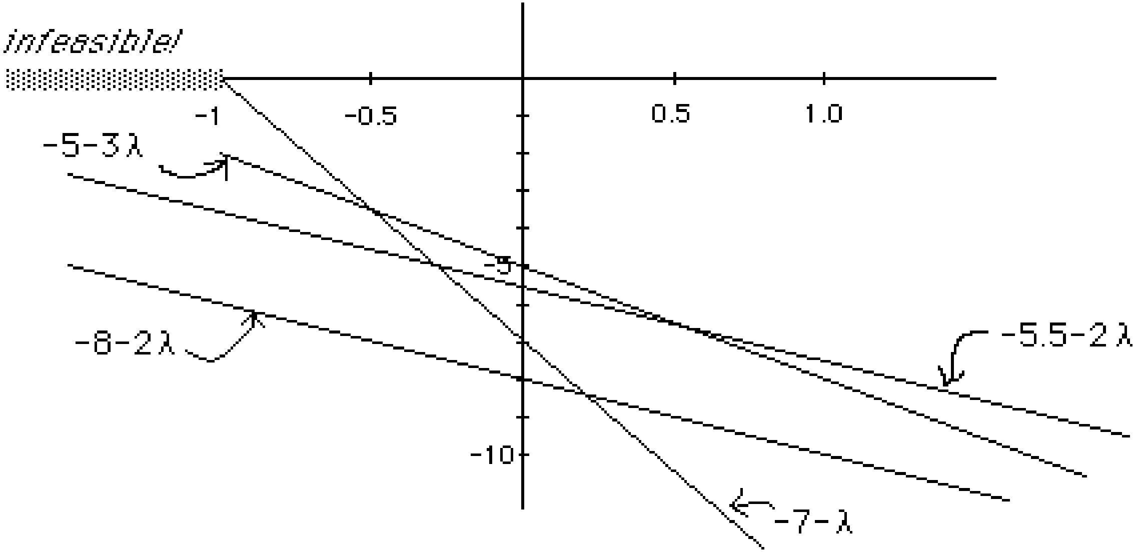
nonpositive because of direction of ≤ in the primal!

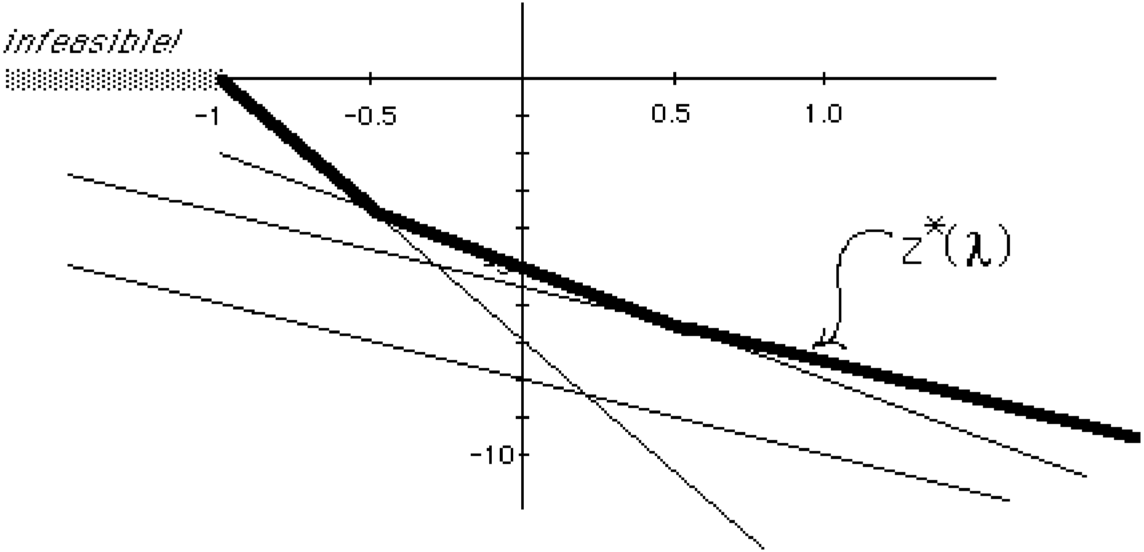
			intercept	*slope*		
π_1	π_2	π_3	$\pi^k b$	$\pi^k d$	*Basis*	
0	0	0	0	0	3 4 5	
-0.5	0	0	-4	-1	1 4 5	
-1	0	0	-8	-2	2 4 5	⇐
-0.333	-0.333	0	-5	-3	1 2 5	⇐
0	-1	0	-7	-7	1 3 5	⇐
0	-0.5	0	-3.5	-3.5	2 3 5	
0	-1	1	-4	-5	1 2 3	
-0.5	0	-0.5	-5.5	-2	1 2 4	⇐
0	0	-1	-3	-2	2 3 4	

feasible

(columns #1,3,4 are dependent & do not form a basis!)

Of the nine basic solutions, four are dual feasible. Therefore, $z^*(\lambda)$ is the maximum of four linear functions:




infeasible!
-1
-0.5
0.5
1.0
-10
$z^*(\lambda)$

In this example, with only nine basic dual solutions, it was possible to enumerate all of them, test each for feasibility, and then maximize the corresponding linear functions

$$\pi^k b + (\pi^k d)\lambda$$

However, for most problems, the number of basic dual solutions is astronomical and enumerating them is practically impossible.

(Usually, only a few of these basic dual solutions actually determine z.)*

Let's consider again the parametric LP $\mathbf{P}_\lambda$:

$$z^*(\lambda) = \text{minimum } -x_1 - x_2$$
$$\text{subject to } \begin{cases} 2x_1 + x_2 \leq 8 + 2\lambda \\ x_1 + 2x_2 \leq 7 + 7\lambda \\ x_2 \leq 3 + 2\lambda \\ x_1 \geq 0,\ x_2 \geq 0 \end{cases}$$

Determine the optimal value function $z^*(\lambda)$ as well as $x_1^*(\lambda)$ and $x_2^*(\lambda)$ [i.e., the optimal solution] for all values of $\lambda \varepsilon (-\infty, +\infty)$

The initial tableau:

change column

−z	1	2	3	4	5	B	Δ
1	¯1	¯1	0	0	0	0	0
0	2	1	1	0	0	8	2
0	1	2	0	1	0	7	7
0	0	1	0	0	1	3	2

slack variables

RHS is

B + λ Δ

Let's start with λ = 0, and investigate the LP as λ increases.

The optimal tableau for $\lambda = 0$:

	−z	1	2	3	4	5	B	Δ
(min)	1	0	0	0.333	0.333	0	5	3
	0	1	0	0.667	¯0.333	0	3	¯1
	0	0	1	¯0.333	0.667	0	2	4
	0	0	0	0.333	¯0.667	1	1	¯2

↑

(this column updated during each pivot)

The optimal solution of $P(0)$ is $z^*(0) = -5$, at $x_1^*(0) = 3$, $x_2^*(0) = 2$, $x_5^*(0) = 1$,

$x_3^*(0)=x_4^*(0)=0$

Expressed as functions of λ, the basic solution is:

−z	1	2	5	B	Δ
1	0	0	0	5	3
0	1	0	0	3	⁻1
0	0	1	0	2	4
0	0	0	1	1	⁻2

$$\Longrightarrow \begin{cases} z(\lambda) = -5 + 3\lambda \\ x_1^*(\lambda) = 3 - \lambda \\ x_2^*(\lambda) = 2 + 4\lambda \\ x_5^*(\lambda) = 1 - 2\lambda \\ x_3^*(\lambda) = x_4^*(\lambda) = 0 \end{cases}$$

Note that these are linear functions of λ!

The optimality criterion (reduced cost ≥ 0) is independent of the parameter λ, and so the current basis remains optimal so long as the basic variables

$$x_1^*(\lambda) = 3 - \lambda$$
$$x_2^*(\lambda) = 2 + 4\lambda$$
$$x_5^*(\lambda) = 1 - 2\lambda$$

remain feasible, i.e., nonnegative.

For what values of λ is $x^*(\lambda) \geq 0$?

We solve the inequalities

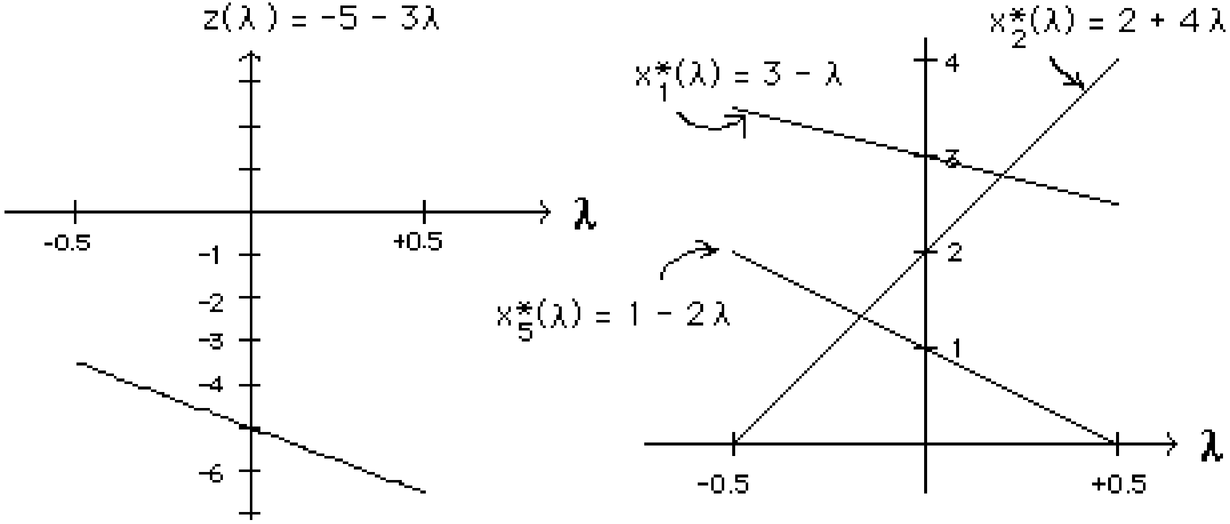
$$x_1^*(\lambda) = 3 - \lambda \ge 0$$
$$x_2^*(\lambda) = 2 + 4\lambda \ge 0$$
$$x_5^*(\lambda) = 1 - 2\lambda \ge 0$$

for λ:

$$3 - \lambda \ge 0 \Rightarrow \lambda \le 3$$
$$2 + 4\lambda \ge 0 \Rightarrow \lambda \ge -1/2$$
$$1 - 2\lambda \ge 0 \Rightarrow \lambda \le +1/2$$

That is, the basic solution is feasible for all $\lambda \in [-1/2, +1/2]$ (and, in particular, for $\lambda = 0$).

Plot of $z(\lambda)$ & $x_j^*(\lambda)$ over the interval [-0.5, +0.5]:



Let us now increase λ from its initial value (0) to the upper limit for which the basis is feasible, i.e., $+\frac{1}{2}$. The basic solution then becomes

$$x_1^*(+\tfrac{1}{2}) = 3 - \tfrac{1}{2} = 2.5$$

$$x_2^*(+\tfrac{1}{2}) = 2 + 4(\tfrac{1}{2}) = 4.0$$

$$x_5^*(+\tfrac{1}{2}) = 1 - 2(\tfrac{1}{2}) = 0$$

Any further increase in λ would result in infeasible (i.e., negative) values for x_5^* !

In order to increase the parameter λ further, x_5 must leave the basis, since it would otherwise become negative. In order to remove x_5 from the basis, we perform a DUAL SIMPLEX pivot:

−z	1	2	3	4	5	B	Δ
1	0	0	0.333	0.333	0	5	3
0	1	0	0.667	¯0.333	0	3	¯1
0	0	1	¯0.333	0.667	0	2	4
0	0	0	0.333	¯0.667	1	1	¯2

← $= 0$ when $\lambda = 1/2$

pivot in this row!

−z	1	2	3	4	5	B	Δ
1	0	0	0.333	0.333	0	5	3
0	1	0	0.667	¯0.333	0	3	¯1
0	0	1	¯0.333	0.667	0	2	4
0	0	0	0.333	$\boxed{¯0.667}$	1	1	¯2

pivot

−z	1	2	3	4	5	B	Δ
1	0	0	0.5	0	0.5	5.5	2
0	1	0	0.5	0	¯0.5	2.5	0
0	0	1	0	0	1	3	2
0	0	0	¯0.5	1	¯1.5	¯1.5	3

For this basis, the basic solution is

$$\begin{array}{cccccc|c|c} -z & 1 & 2 & \cdots & 4 & \cdots & B & \Delta \\ \hline 1 & 0 & 0 & & 0 & & 5.5 & 2 \\ 0 & 1 & 0 & & 0 & & 2.5 & 0 \\ 0 & 0 & 1 & & 0 & & 3 & 2 \\ 0 & 0 & 0 & & 1 & & -1.5 & 3 \end{array} \Rightarrow \begin{cases} z(\lambda) = -5.5 - 2\lambda \\ x_1^*(\lambda) = 2.5 \\ x_2^*(\lambda) = 3 + 2\lambda \\ x_4^*(\lambda) = -1.5+3\lambda \end{cases}$$

Notice that as λ increases, no basic variable decreases. Since the optimality criterion (reduced costs ≥ 0) does not depend upon λ, this basis is optimal for all $\lambda \geq 0.5$

That is, if we solve the system of inequalities

$$x_1^*(\lambda) = 2.5 \geq 0 \Longrightarrow \text{(no restriction on } \lambda \text{.)}$$

$$x_2^*(\lambda) = 3 + 2\lambda \geq 0 \Longrightarrow \lambda \geq {}^-1.5$$

$$x_4^*(\lambda) = -1.5+3\lambda \geq 0 \Longrightarrow \lambda \geq 0.5$$

we see that it is satisfied for $\lambda \in [\, 0.5, +\infty \,)$

Let us now investigate $P(\lambda)$ for $\lambda < -0.5$
Consider the tableau which was optimal for
$-0.5 \leq \lambda \leq +0.5$

−z	1	2	3	4	5	B	Δ
1	0	0	0.333	0.333	0	5	3
0	1	0	0.667	¯0.333	0	3	¯1
0	0	1	¯0.333	0.667	0	2	4
0	0	0	0.333	¯0.667	1	1	¯2

Recall that the lower limit of the parameter,
$\lambda \geq -\frac{1}{2}$, derives from $x_2^*(\lambda) \geq 0$, i.e., $x_2^*(-\frac{1}{2}) = 0$

A further decrease in λ requires that x_2 be removed from the basis (by a dual simplex pivot)

−Z	1	2	3	4	5	B	Δ	
1	0	0	0.333	0.333	0	5	3	
0	1	0	0.667	¯0.333	0	3	¯1	
0	0	1	¯0.333	0.667	0	2	4	← pivot row
0	0	0	0.333	¯0.667	1	1	¯2	

pivot here!

the dual simplex pivot yields

⟹

−Z	1	2	3	4	5	B	Δ
1	0	1	0	1	0	7	7
0	1	2	0	1	0	7	7
0	0	¯3	1	¯2	0	¯6	¯12
0	0	1	0	0	1	3	2

The new basic solution is:

$-z$	1	3	5	B	Δ
1	0	0	0	7	7
0	1	0	0	7	7
0	0	1	0	-6	-12
0	0	0	1	3	2

$$\Rightarrow \begin{cases} z(\lambda) = -7 - 7\lambda \\ x_1^*(\lambda) = 7 + 7\lambda \\ x_3^*(\lambda) = -6 - 12\lambda \\ x_5^*(\lambda) = 3 + 2\lambda \end{cases}$$

To find the interval for which this basic solution is feasible (& therefore optimal), solve

$$\begin{aligned} x_1^*(\lambda) &= 7 + 7\lambda \geq 0 \\ x_3^*(\lambda) &= -6 - 12\lambda \geq 0 \\ x_5^*(\lambda) &= 3 + 2\lambda \geq 0 \end{aligned} \Longrightarrow \begin{cases} \lambda \geq -1 \\ \lambda \leq -0.5 \\ \lambda \geq -1.5 \end{cases}$$

that is,

$\lambda \varepsilon [-1.0,-0.5]$

When λ decreases to −1.0, $x_1^*(\lambda)$ decreases to 0 and must be removed from the basis to allow any further decrease in the parameter. We therefore attempt another dual simplex pivot:

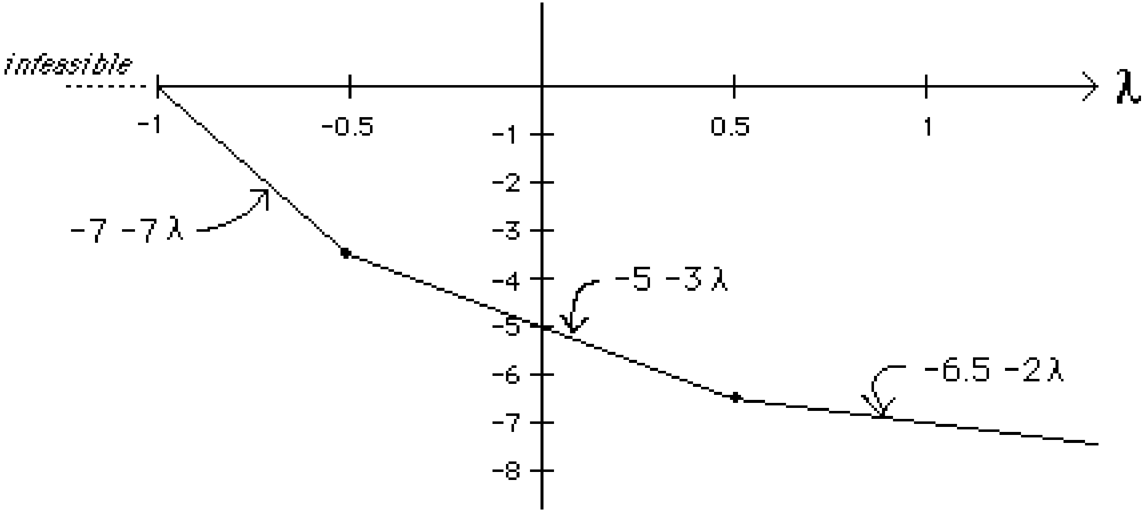
¯z	1	2	3	4	5	B	Δ	
1	0	1	0	1	0	7	7	
0	1	2	0	1	0	7	7	⟵ *pivot row*
0	0	¯3	1	¯2	0	¯6	¯12	
0	0	1	0	0	1	3	2	

Because there is no negative element in the pivot row, x_1 cannot be removed from the basis, and it is evident that $P(\lambda)$ is infeasible for $\lambda < -1.0$

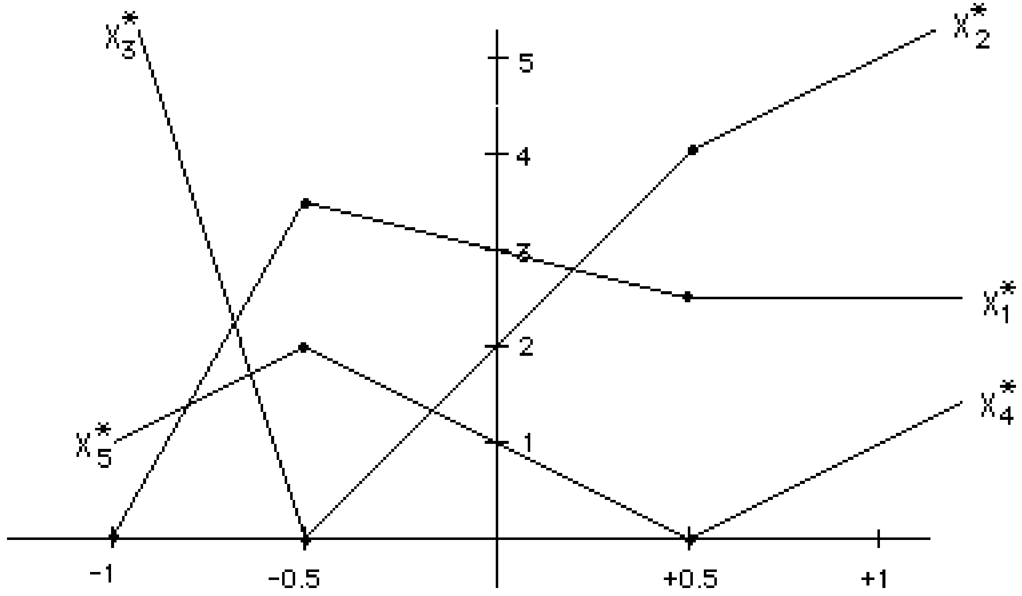
Summary of Parametric Analysis:

λ	$(-\infty, -1]$	$[-1, -0.5]$	$[-0.5, +0.5]$	$[+0.5, +\infty)$
x_1^*	*infeasible*	$7 + 7\lambda$	$3 - \lambda$	2.5
x_2^*		0	$2 + 4\lambda$	$3 + 2\lambda$
x_3^*		$-6 - 12\lambda$	0	0
x_4^*		0	0	$-1.5 + 3\lambda$
x_5^*		$3 + 2\lambda$	$1 - 2\lambda$	0

Plot of z vs λ :



Plot of x_j^* vs. λ



Example

Minimize $x_1 + x_2 + 7x_3 + 3x_4 + x_5 + 2x_6$

subject to

$$\begin{aligned} x_1 + 2x_2 - x_3 - x_4 + x_5 + 2x_6 &= 16 - \lambda \\ x_2 - 3x_4 - x_5 + 3x_6 &= 2 + \lambda \\ -x_1 - 3x_3 + 3x_4 + x_5 &= -4 + \lambda \end{aligned}$$

$x_j \geq 0,\ j=1,2, \ldots 6$

Initial tableau:

−Z	1	2	3	4	5	6	B	Δ
1	1	1	7	3	1	2	0	0
0	1	2	−1	−1	1	2	16	−1
0	0	1	0	−3	−1	3	2	1
0	−1	0	−3	3	1	0	−4	1

Optimal tableau ($\lambda = 0$)

−Z	1	2	3	4	5	6	B	Δ
1	0	0	7	3	0	2	−12	1.5
0	1	0	2	−1	0	−1	6	−1.5
0	0	1	−1	−1	0	2	4	0.5
0	0	0	−1	2	1	−1	2	−0.5

¯Z	1	2	3	4	5	6	B	Δ
1	0	0	7	3	0	2	¯12	1.5
0	1	0	2	¯1	0	¯1	6	¯1.5
0	0	1	¯1	¯1	0	2	4	0.5
0	0	0	¯1	2	1	¯1	2	¯0.5

Parametric Analysis

```
Least Upper Bound (LUB): 4
         = Min{¯6 ¯2 ÷ ¯1.5 ¯0.5} = Min {4 4}
        RHS at LUB is ¯10 0 6 0
Greatest Lower Bound (GLB): ¯8
         = Max{¯4 ÷ 0.5} = Max {¯8}
        RHS at GLB is ¯16 18 0 6
Range of parameters LAMBDA within which basis is feasible:
        [ ¯8 , 4 ]
```

Dual Simplex pivot

¯z	1	2	3	4	5	6	B	Δ	
1	0	0	7	3	0	2	¯12	1.5	
0	1	0	2	¯1	0	¯1	6	¯1.5	←
0	0	1	¯1	¯1	0	2	4	0.5	
0	0	0	¯1	2	1	¯1	2	¯0.5	

New tableau

¯z	1	2	3	4	5	6	B	Δ
1	2	0	11	1	0	0	0	¯1.5
0	¯1	0	¯2	1	0	1	¯6	1.5
0	2	1	3	¯3	0	0	16	¯2.5
0	¯1	0	¯3	3	1	0	¯4	1

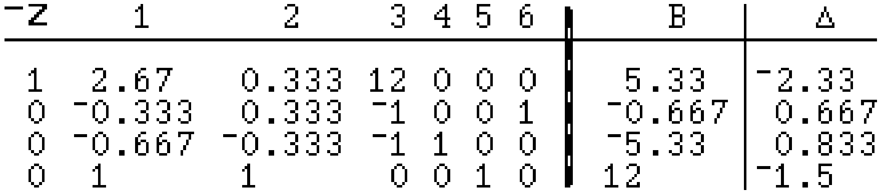
```
¯Z  1  2  3   4  5  6 ┃   B │   Δ
─────────────────────────────────────
 1   2  0 11  1  0  0 ┃   0 │ ¯1.5
 0  ¯1  0 ¯2  1  0  1 ┃  ¯6 │  1.5
 0   2  1  3 ¯3  0  0 ┃  16 │ ¯2.5
 0  ¯1  0 ¯3  3  1  0 ┃  ¯4 │  1
```

```
                 Parametric Analysis
Least Upper Bound (LUB): 6.4
          = Min{¯16 ÷ ¯2.5} = Min {6.4}
         RHS at LUB is ¯16 3.6 0 2.4
Greatest Lower Bound (GLB): 4
          = Max{6 4 ÷ 1.5 1} = Max {4 4}
         RHS at GLB is ¯10 0 6 0
Range of parameters LAMBDA within which basis is feasible:
         [ 4 , 6.4 ]
```

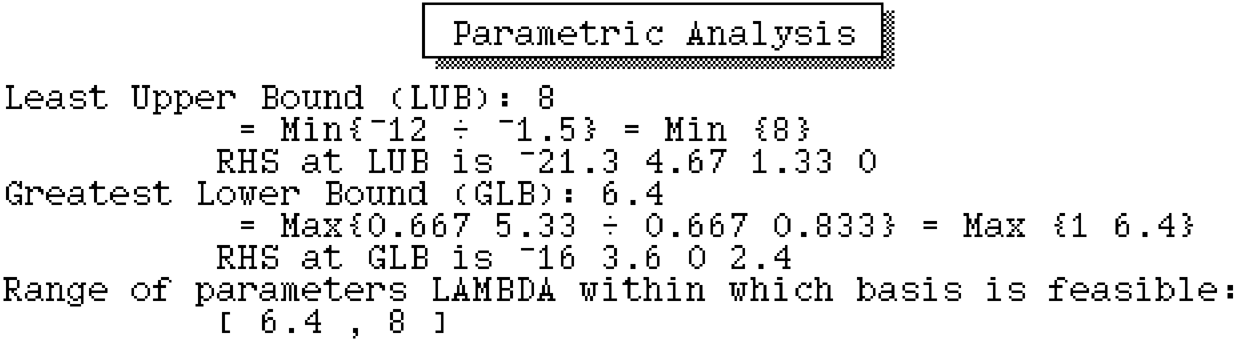
Dual Simplex Pivot

¯Z	1	2	3	4	5	6	B	Δ	
1	2	0	11	1	0	0	0	¯1.5	
0	¯1	0	¯2	1	0	1	¯6	1.5	
0	2	1	3	¯3	0	0	16	¯2.5	←
0	¯1	0	¯3	3	1	0	¯4	1	

¯z	1	2	3	4	5	6	B	Δ
1	2.67	0.333	12	0	0	0	5.33	¯2.33
0	¯0.333	0.333	¯1	0	0	1	¯0.667	0.667
0	¯0.667	¯0.333	¯1	1	0	0	¯5.33	0.833
0	1	1	0	0	1	0	12	¯1.5



```
¯Z   1       2      3 4 5 6 ┃    B    │  Δ
─────────────────────────────┃─────────┼──────
 1   2.67    0.333 12 0 0 0 ┃  5.33   │¯2.33
 0 ¯0.333    0.333 ¯1 0 0 1 ┃ ¯0.667  │ 0.667
 0 ¯0.667  ¯0.333 ¯1 1 0 0 ┃ ¯5.33   │ 0.833
 0   1       1        0 0 1 0 ┃ 12      │¯1.5
```



```
                    Parametric Analysis
Least Upper Bound (LUB): 8
          = Min{¯12 ÷ ¯1.5} = Min {8}
         RHS at LUB is ¯21.3 4.67 1.33 0
Greatest Lower Bound (GLB): 6.4
          = Max{0.667 5.33 ÷ 0.667 0.833} = Max {1 6.4}
         RHS at GLB is ¯16 3.6 0 2.4
Range of parameters LAMBDA within which basis is feasible:
         [ 6.4 , 8 ]
```

A dual simplex pivot in row #4 is not possible:

¯z	1	2	3	4	5	6	B	Δ
1	2.67	0.333	12	0	0	0	5.33	¯2.33
0	¯0.333	0.333	¯1	0	0	1	¯0.667	0.667
0	¯0.667	¯0.333	¯1	1	0	0	¯5.33	0.833
0	1	1	0	0	1	0	12	¯1.5

The LP is infeasible for $\lambda > 8$

Let's return to the optimal tableau for $\lambda = 0$:

¯Z	1	2	3	4	5	6	B	Δ
1	0	0	7	3	0	2	¯12	1.5
0	1	0	2	¯1	0	¯1	6	¯1.5
0	0	1	¯1	¯1	0	2	4	0.5
0	0	0	¯1	2	1	¯1	2	¯0.5

```
                    Parametric Analysis
Least Upper Bound (LUB): 4
          = Min{¯6 ¯2 ÷ ¯1.5 ¯0.5} = Min {4 4}
         RHS at LUB is ¯10 0 6 0
Greatest Lower Bound (GLB): ¯8
          = Max{¯4 ÷ 0.5} = Max {¯8}
         RHS at GLB is ¯16 18 0 6
Range of parameters LAMBDA within which basis is feasible:
         [ ¯8 , 4 ]
```

Dual Simplex Pivot

−Z	1	2	3	4	5	6	B	Δ	
1	0	0	7	3	0	2	−12	1.5	
0	1	0	2	−1	0	−1	6	−1.5	
0	0	1	−1	$\boxed{-1}$	0	2	4	0.5	←
0	0	0	−1	2	1	−1	2	−0.5	

New tableau

−Z	1	2	3	4	5	6	B	Δ
1	0	3	4	0	0	8	0	3
0	1	−1	3	0	0	−3	2	−2
0	0	−1	1	1	0	−2	−4	−0.5
0	0	2	−3	0	1	3	10	0.5

```
¯z 1  2  3 4 5  6 ┃  B │  Δ
─────────────────────────────
 1 0  3  4 0 0  8 ┃  0 │ 3
 0 1 ¯1  3 0 0 ¯3 ┃  2 │¯2
 0 0 ¯1  1 1 0 ¯2 ┃ ¯4 │¯0.5
 0 0  2 ¯3 0 1  3 ┃ 10 │ 0.5
```

Parametric Analysis

```
Least Upper Bound (LUB): ¯8
          = Min{¯2 4 ÷ ¯2 ¯0.5} = Min {1 ¯8}
         RHS at LUB is ¯16 18 0 6
Greatest Lower Bound (GLB): ¯20
          = Max{¯10 ÷ 0.5} = Max {¯20}
         RHS at GLB is ¯40 42 6 0
Range of parameters LAMBDA within which basis is feasible:
         [ ¯20 , ¯8 ]
```

Dual Simplex Pivot

¯Z	1	2	3	4	5	6	B	Δ	
1	0	3	4	0	0	8	0	3	
0	1	¯1	3	0	0	¯3	2	¯2	
0	0	¯1	1	1	0	¯2	¯4	¯0.5	
0	0	2	$\boxed{¯3}$	0	1	3	10	0.5	←

New tableau

¯Z	1	2	3	4	5	6	B	Δ
1	0	5.67	0	0	1.33	12	13.3	3.67
0	1	1	0	0	1	0	12	¯1.5
0	0	¯0.333	0	1	0.333	¯1	¯0.667	¯0.333
0	0	¯0.667	1	0	¯0.333	¯1	¯3.33	¯0.167

```
¯Z 1     2     3 4     5      6 ‖    B     |   Δ
 ─────────────────────────────────────────────────
 1 0   5.67   0 0   1.33   12  ‖ 13.3      |  3.67
 0 1   1      0 0   1       0  ‖ 12        | ¯1.5
 0 0 ¯0.333   0 1   0.333  ¯1  ‖  ¯0.667   | ¯0.333
 0 0 ¯0.667   1 0  ¯0.333  ¯1  ‖  ¯3.33    | ¯0.167
```

```
                    Parametric Analysis

Least Upper Bound (LUB): ¯20
           = Min{¯12 0.667 3.33 ÷ ¯1.5 ¯0.333 ¯0.167}
           = Min {  8 ¯2 ¯20}
          RHS at LUB is ¯60 42 6 0

          No Lower Bound
Range of parameters LAMBDA within which basis is feasible:
          [ ¯1.8E308 , ¯20 ]
```

i.e., - ∞