

**56:272 Integer Programming & Network Flows**  
**Final Examination -- December 14, 1998**

- Part A: Answer any **four** of the **five** problems. (15 points each)
1. Transportation problem \_\_\_\_\_
  2. Integer LP Model Formulation \_\_\_\_\_
  3. Balas' Additive Algorithm (Implicit Enumeration) \_\_\_\_\_
  4. Generalized Assignment Problem \_\_\_\_\_
  5. Benders' Decomposition of Capacitated Plant Location Problem \_\_\_\_\_
- Part B: Answer any **seven** of the **eight** problems. (10 points each)
6. Branch-&-bound solution of knapsack problem \_\_\_\_\_
  7. Asymmetric Traveling salesman problem \_\_\_\_\_
  8. Primal Simplex Algorithm for Networks \_\_\_\_\_
  9. Gomory's Cutting Plane Algorithm \_\_\_\_\_
  10. Chinese Postman Problem \_\_\_\_\_
  11. Vertex Penalty Method for TSP \_\_\_\_\_
  12. Shortest Path Problem \_\_\_\_\_
  13. Location in a network \_\_\_\_\_
- Total possible: 130 points

\*\*\*\*\* PART A \*\*\*\*\*

**1. Transportation Problem:** Consider the transportation problem with the tableau below:

|          |          |          |          |          |          |        |
|----------|----------|----------|----------|----------|----------|--------|
|          | <b>A</b> | <b>B</b> | <b>C</b> | <b>D</b> | <b>E</b> | supply |
| <b>F</b> | 6        | 3        | 4        | 1        | 2        | 4      |
| <b>G</b> | 4        | 2        | 5        | 6        | 3        | 7      |
| <b>H</b> | 5        | 1        | 7        | 2        | 4        | 9      |
| demand   | 5        | 4        | 4        | 5        | 2        |        |

- a. If the ordinary simplex tableau were to be written for this problem, it would have \_\_\_ non-redundant constraint rows, not including the objective row, and \_\_\_ columns (not including -z and the right-hand-side).
- b. This problem will have \_\_\_ basic variables (not including -z).
- c. Find an initial basic feasible solution using the "Northwest Corner Method". (Write the values of the variables in the tableau above.)
- d. What are the values of the dual variables for the solution in (c)? (Note that  $V_A$  has been assigned the value zero)  $U_F = \underline{\quad}$ ,  $U_G = \underline{\quad}$ ,  $U_H = \underline{\quad}$ ,  
 $V_A = \underline{0}$ ,  $V_B = \underline{\quad}$ ,  $V_C = \underline{\quad}$ ,  $V_D = \underline{\quad}$ ,  $V_E = \underline{\quad}$ .
- e. What is the reduced cost of the variable  $X_{FD}$ ? \_\_\_\_\_
- f. Will increasing  $X_{FD}$  improve the objective function? \_\_\_\_\_
- g. Regardless of whether the answer to (f) is "yes" or "no", what variable must leave the basis if  $X_{FD}$  enters? \_\_\_\_\_
- h. What will be the value of  $X_{FD}$  if it is entered into the solution as in (g)? \_\_\_\_\_
- i. Which variables (if any), if it were entered into the solution, would result in a degenerate solution?  
 \_\_\_\_\_

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**2. Integer LP Model Formulation** The board of directors of a large manufacturing firm is considering the set of investments shown below: Let  $R_i$  be the total revenue from investment  $i$  and  $C_i$  the cost (in

\$millions) to make investment  $i$ . The board wishes to maximize total revenue and invest no more than a total of 50 million dollars.

| Investment<br>$i$ | Revenue<br>$R_i$ | Cost<br>$C_i$ | Condition             |
|-------------------|------------------|---------------|-----------------------|
| 1                 | 1                | 5             | None                  |
| 2                 | 2                | 8             | Only if #1            |
| 3                 | 3                | 12            | Only if #2            |
| 4                 | 4                | 18            | Must if #1 and #2     |
| 5                 | 5                | 24            | Not if #3 or #4       |
| 6                 | 6                | 27            | Not if both #1 and #3 |
| 7                 | 7                | 30            | Only if #4 and not #1 |

Define  $X_i = 1$  if investment  $i$  is selected, else 0.

- Formulate this problem without the "side conditions" as an integer LP.
- This is an example of a special class of integer programming problems called "\_\_\_\_\_ " problems.
- Add a constraint to enforce the condition "Investment #2 only if #1 is selected".
- Add a constraint to enforce the condition "Investment #3 only if #2 is selected".
- Add a constraint to enforce the condition "Investment #4 must be selected if both #1 & #2 are selected".
- Add a constraint to enforce the condition "Investment #5 cannot be selected if both #3 & #4 are selected".
- Add a constraint to enforce the condition "Investment #6 cannot be selected if both #1 & #3 are selected".
- Add a constraint to enforce the condition "Investment #7 only if #4 is selected but #1 is not selected".
- Suppose that investments 3 and 4 are in products which compete against one another to some extent, so that 1 must be subtracted from the total revenue if both of these two investments are selected. Reformulate the problem (without the "side constraints" above.) *Define any new variable(s) which you use.*

✱ ✱ ✱ ✱ ✱ ✱ ✱ ✱ ✱ ✱ ✱ ✱ ✱ ✱ ✱ ✱

**3. Balas' Additive Algorithm (Implicit Enumeration)** Consider the problem

$$\begin{array}{ll}
 \text{Maximize} & X_1 + 8X_2 + 3X_3 - 5X_4 - 3X_5 \\
 \text{subject to} & -2X_1 - X_2 + 2X_3 + 3X_4 - 3X_5 = 1 \\
 & X_2 - 2X_3 - 4X_4 - 2X_5 = -1 \\
 & 2X_1 + 4X_3 + 3X_4 = 1 \\
 & X_i = 0 \text{ or } 1, i=1,2,\dots, 5
 \end{array}$$



b. Suppose that the integer restrictions are relaxed and the problem were solved by the simplex LP algorithm. Will the optimal values of the variables of this LP relaxation necessarily be integer? \_\_\_\_\_

c. Suppose that the machine capacity constraints are relaxed, using the Lagrangian relaxation method with multipliers  $U_1, U_2,$  and  $U_3$ . The first 2 iterations of the subgradient optimization method to maximize the lower bound appears below, where the optimal value was estimated to be 120, and a stepsize parameter was assigned the value 0.75.

```

Lambda = 0.75
Upper bound Z* = 120
Iteration # 1
Multiplier vector U = 0 0 0
Objective function of relaxation:
machine
job
1 10 18 11 17 24 13
2 11 20 25 24 18 13
3 23 11 17 18 18 19

Dual value is A
Variables selected from GUB sets are:
1 3 1 1 2 1
Resources used are: 79 24 21, (Available: 20 27 38)
Subgradient of Dual Objective is B C D
Stepsize is 0.00861821

Multiplier vector U = 0.508475 0 0
Objective function of relaxation:
machine
job
1 E 29.18 23.20 27.67 36.20 21.13
2 11 20 25 24 18 13
3 23 11 17 18 18 19

Dual value is 77.8305
Variables selected from GUB sets are:
F G H I 2 J
Resources used are: 0 53 50, (Available: 20 27 38)
Subgradient of Dual Objective is 20 26 12
Stepsize is K

```

d. Several values have been omitted from the output. ("Variables selected from GUB sets" refers to the machine selected for each of the jobs.) Compute their values:

A \_\_\_\_\_ B \_\_\_\_\_ C \_\_\_\_\_ D \_\_\_\_\_ E \_\_\_\_\_  
 F \_\_\_\_\_ G \_\_\_\_\_ H \_\_\_\_\_ I \_\_\_\_\_ J \_\_\_\_\_  
 K \_\_\_\_\_

e. What is the "integrality property" of a Lagrangian relaxation?

f. Does this particular Lagrangian relaxation have the integrality property? Circle: Yes No

g. What does your answer in (f) imply about the strength of the lower bound which can be obtained from this relaxation, compared to that of the LP relaxation in (b)? (Is it  $\geq$ ,  $\leq$ , or  $=$  ?)

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**5. Benders' Decomposition of Capacitated Plant Location Problem**

Consider the following problem in which demand in 8 cities is to be satisfied by plants to be built in one or more of cities 1,2,3, & 4:



Number of sources = M = 4  
 Number of destinations = N = 8  
 Total demand: 40

K = capacity,  
 F = fixed cost

Costs, Supplies, Demands

| i \ j  | 1  | 2  | 3  | 4  | 5  | 6  | 7  | 8  | K  | F   |
|--------|----|----|----|----|----|----|----|----|----|-----|
| 1      | 0  | 42 | 75 | 33 | 17 | 38 | 33 | 30 | 19 | 498 |
| 2      | 42 | 0  | 43 | 32 | 56 | 47 | 19 | 52 | 15 | 23  |
| 3      | 75 | 43 | 0  | 46 | 82 | 90 | 62 | 67 | 17 | 89  |
| 4      | 33 | 32 | 46 | 0  | 36 | 61 | 40 | 22 | 17 | 129 |
| Demand | 6  | 1  | 2  | 5  | 5  | 8  | 3  | 10 | 68 |     |

a. State the mixed-integer programming formulation of the problem. How many continuous variables (X) and how many binary (zero-one) variables (Y) are required?

A trial solution was evaluated, in which all four plants are to be open. The result was:

|                              |   |
|------------------------------|---|
| Subproblem Solution          | Generated support is $\alpha Y + b$ , where |
| Plants opened: # 1 2 3 4     | $\alpha = 498 \ 23 \ 89 \ 129$              |
| Minimum transport cost = 666 | & $b = 666$                                 |
| Fixed cost of plants = 739   | That is, $v(Y) \geq \alpha Y + b$           |
| Total = 1405                 | This is support # 1                         |

The Master Problem was next optimized. (A constraint  $\sum_{i=1}^4 K_i Y_i \geq \sum_{j=1}^8 D_j$  was included in the master problem in order to guarantee that only solutions with sufficient capacity to meet the demand were produced.) The result was:

|  |
|--|
| Optimum of Master Problem  |
| Optimal set of plants: 2 3 4   |
| with estimated cost <span style="border: 1px solid black; display: inline-block; width: 50px; height: 15px; vertical-align: middle;"></span> |

b. What is the value of the estimated cost of  $Y=(0,1,1,1)$  found by the Master problem ("blanked") above?

Next, the subproblem was solved, using the trial set of plants {2,3,4}, with the following results:

|   |   |
|---|---|
| Subproblem Solution                         | Generated support is $\alpha Y + b$ , where |
| Plants opened: # 2 3 4                      | $\alpha = 498 \ 653 \ 1364 \ 639$           |
| Minimum transport cost = 1319               | & $b = -1096$                               |
| Fixed cost of plants = <input type="text"/> | That is, $v(Y) \geq \alpha Y + b$           |
| Total = <input type="text"/>                | This is support # 2                         |

c. What are the two values blanked above?

When the Master problem is solved once more, the result is:

|  |
|--|
| Optimum of Master Problem                |
| Optimal set of plants: 1 2 4             |
| with estimated cost <input type="text"/> |

d. Compute the estimated cost found by the master problem.

The subproblem is again solved, this time with trial set {1,2,4}, yielding:

|                              |   |
|------------------------------|---|
| Subproblem Solution          | Generated support is $\alpha Y + b$ , where     |
| Plants opened: # 1 2 4       | $\alpha = 1144 \ 668 \ 89$ <input type="text"/> |
| Minimum transport cost = 752 | & $b = -1270$                                   |
| Fixed cost of plants = 650   | That is, $v(Y) \geq \alpha Y + b$               |
| Total = 1402                 | This is support # 3                             |

e. Using the information below from the subproblem solution, compute the blanked value of  $\alpha_4$  above.

|                   |                                   |
|-------------------|-----------------------------------|
| Optimal Shipments | Dual Variables                    |
| f                 | Supply constraints                |
| r                 | i= 1 2 3 4                        |
| o                 | U[i]= 34 43 0 43                  |
| m                 | Demand constraints                |
| 1                 | j= 1 2 3 4 5 6 7 8                |
| 2                 | V[j]= -34 -43 0 -43 -17 4 -24 -21 |
| 4                 |                                   |

(Demand pt #9 is dummy demand for excess capacity.)

|                                     |    |    |    |    |    |    |    |
|-------------------------------------|----|----|----|----|----|----|----|
| Reduced costs: COST - $U \cdot + V$ |    |    |    |    |    |    |    |
| 0                                   | 51 | 41 | 42 | 0  | 0  | 23 | 17 |
| 33                                  | 0  | 0  | 32 | 30 | 0  | 0  | 30 |
| 109                                 | 86 | 0  | 89 | 99 | 86 | 86 | 88 |
| 24                                  | 32 | 3  | 0  | 10 | 14 | 21 | 0  |

f. The solution of the transportation problem (the shipments) is degenerate. Based upon the information given, which variable(s) is/are basic but zero in the solution?

\*\*\*\*\* PART B \*\*\*\*\*

**6. Knapsack Problem:** A "knapsack" is to be filled to maximize the value of the contents, subject to a weight restriction. A maximum of 1 unit of each item may be included.

|                          |                    |                   |                  |                |                  |                     |
|--------------------------|--------------------|-------------------|------------------|----------------|------------------|---------------------|
| Number of items: 4       |                    |                   |                  |                |                  |                     |
| Capacity of knapsack: 10 |                    |                   | Sorted by \$/kg: |                |                  |                     |
| <u>Item</u>              | <u>Weight (kg)</u> | <u>Value (\$)</u> | <u>Item</u>      | <u>Wt (kg)</u> | <u>Value(\$)</u> | <u>Value/Weight</u> |
| 1                        | 3                  | 5                 | 2                | 5              | 9                | 1.8                 |
| 2                        | 5                  | 9                 | 4                | 4              | 7                | 1.75                |
| 3                        | 2                  | 3                 | 1                | 3              | 5                | 1.66667             |
| 4                        | 4                  | 7                 | 3                | 2              | 3                | 1.5                 |

a. Formulate this problem as a 0-1 ILP problem:

*Output of the branch-&-bound algorithm for this problem appears below:*

```

->>>Subproblem # 1
Forced in:
Forced out:
Free:      1 2 3 4
Fractional solution: selected items = 2 4
                    plus  of item # 
                    value = 17.6667
Rounding down yields value 16
*** NEW INCUMBENT! ***
->>>Subproblem # 2
Forced in:      1
Forced out:
Free:           2 3 4
Fractional solution: selected items = 1 2
                    plus 0.5 of item # 4
                    value = 17.5
Rounding down yields value 14
->>>Subproblem # 3
Forced in:      1 4
Forced out:
Free:           2 3
Fractional solution: selected items = 1 4
                    plus 0.6 of item # 2
                    value = 17.4
Rounding down yields value 12
->>>Subproblem # 4
Forced in:      1 2 4
Forced out:
Free:           3
Infeasible!
<<<Subproblem # 4 fathomed.
->>>Subproblem # 5
Forced in:      1 4
Forced out:      2
Free:           3
Integer solution: selected items = 1 3 4
                    Value= 15
<<<Subproblem # 5 fathomed.

```

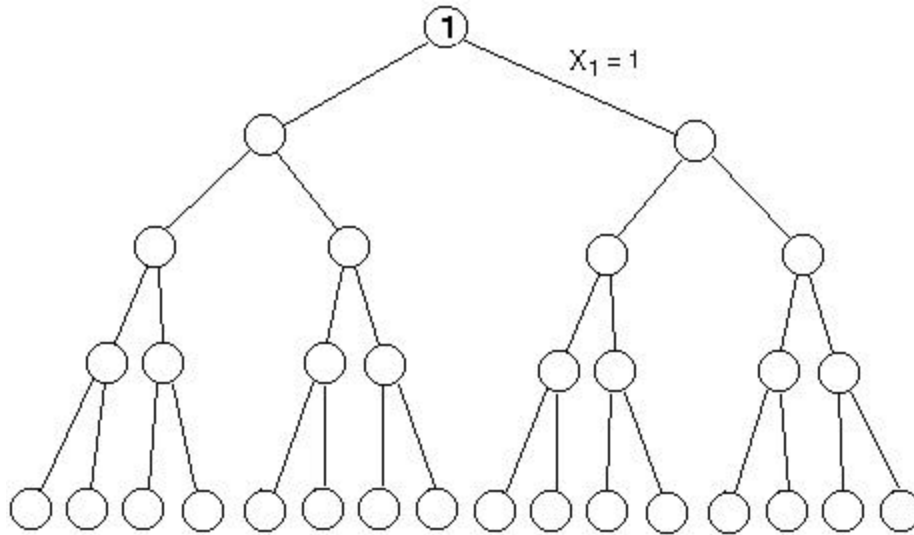
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<<<Subproblem # 3 fathomed.
->>>Subproblem # 6
Forced in: 1
Forced out: 4
Free: 2 3
Integer solution: selected items = 1 2 3
Value= 17
*** NEW INCUMBENT! ***
<<<Subproblem # 6 fathomed.
<<<Subproblem # 2 fathomed.
->>>Subproblem # 7
Forced in:
Forced out: 1
Free: 2 3 4
Fractional solution: selected items = 
plus  of item # 
value = 
Rounding down yields value 16
<<<Subproblem # 7 fathomed.
<<<Subproblem # 1 fathomed.

```

b. Complete the blanks in the output above.

c. Below is a "complete" search tree. Using the information in the above output, complete the numbering of the nodes which were explicitly searched (leaving other nodes which were implicitly searched blank). Also write on each branch the variable which was fixed (& its value):



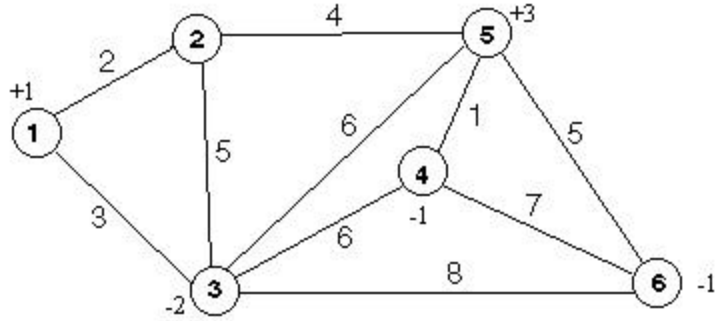
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**7. Asymmetric Traveling Salesman Problem.** Five products are to be manufactured weekly on the same machine. The table below gives the cost of switching the machine from one product to another product. (Assume that this is also the cost of switching to the last product of the week to the first product to be scheduled the following week!)

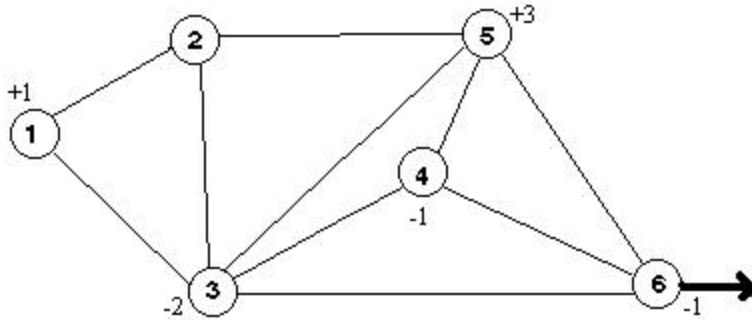
|       |   |   |   |   |   |   |
|-------|---|---|---|---|---|---|
| to:   | A | B | C | D | E |   |
| from: | A | - | 3 | 5 | 3 | 6 |
|       | B | 2 | - | 3 | 2 | 8 |
|       | C | 6 | 3 | - | 6 | 7 |
|       | D | 3 | 1 | 4 | - | 8 |





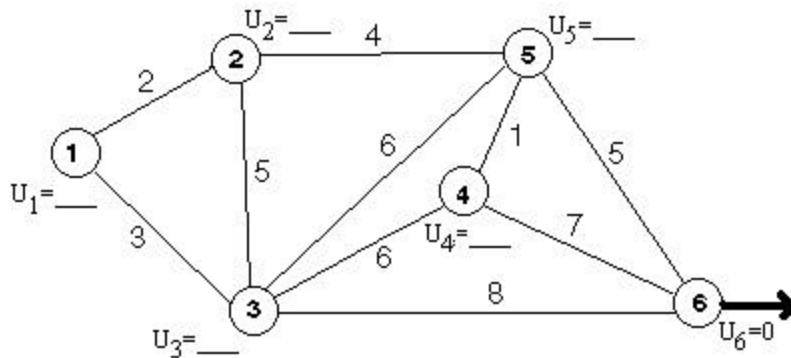


a. Find the minimum spanning tree of this network, and indicate it below.



b. Using the minimum spanning tree (plus artificial "root" arc at node 6) as an initial basis, compute the corresponding basic solution, i.e., flows. Indicate these flows above (indicating direction of the flow).

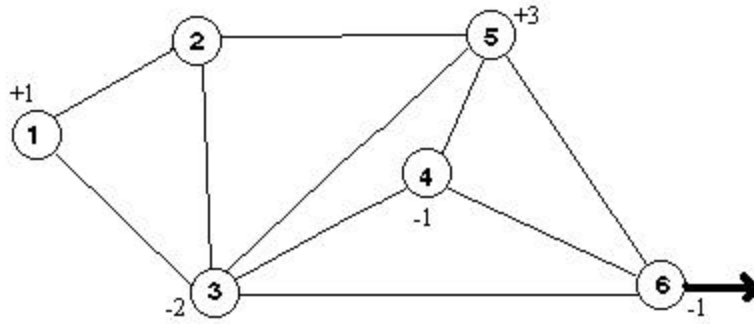
c. Using the same basis, compute the dual variables (simplex multipliers) for each node (letting the dual variable for node 6 be zero), and indicate below:



d. Choose one arc not in the rooted spanning tree, and "price" it, i.e., compute its reduced cost. (*Choose both an edge and a direction!*)

Would entering this arc into the basis result in an improvement? \_\_\_\_

e. Regardless of whether the arc you selected in (d) should enter the basis, explain how to enter the arc into the basis and how to choose the arc leaving the basis. Indicate the new basis on the network below:



\* \* \* \* \*

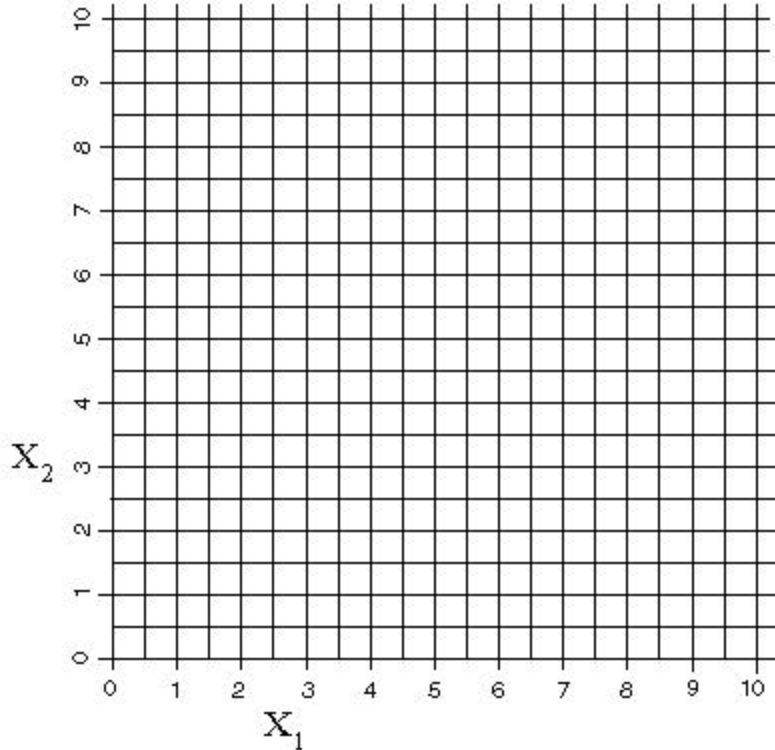
**9. Gomory's Cutting Plane Algorithm.** Consider the problem

$$\begin{aligned}
 &\text{Maximize} && 3X_1 + 5X_2 \\
 &\text{subject to} && -X_1 + X_2 \leq 5 \\
 &&& 9X_1 + 5X_2 \leq 45 \\
 &&& X_1, X_2 \geq 0, \text{ \& integer}
 \end{aligned}$$

After adding slack & surplus variables  $X_3$  &  $X_4$ , respectively, and solving the LP relaxation, we get the optimal tableau:

| $X_1$ | $X_2$ | $X_3$ | $X_4$  | rhs    |       |
|-------|-------|-------|--------|--------|-------|
| 0     | 0     | 0.571 | 2.143  | 36.429 | (max) |
| 0     | 1     | 0.071 | 0.643  | 6.429  |       |
| 1     | 0     | 0.071 | -0.357 | 1.429  |       |

- a. Using the bottom row of the tableau, state a constraint that may be added to the problem to exclude this extreme point of the feasible region of the LP relaxation without excluding any integer feasible solutions.
- b. Express the constraint which you found in (a) in terms of the original variables  $X_1$  and  $X_2$ :  
 $\underline{\hspace{1cm}} X_1 + \underline{\hspace{1cm}} X_2 \geq \underline{\hspace{1cm}}$
- c. Graph, in  $(X_1, X_2)$ -space, the original constraints, the optimum of the LP relaxation, and the new constraint which you chose in (b). Shade the feasible region after adding this constraint.

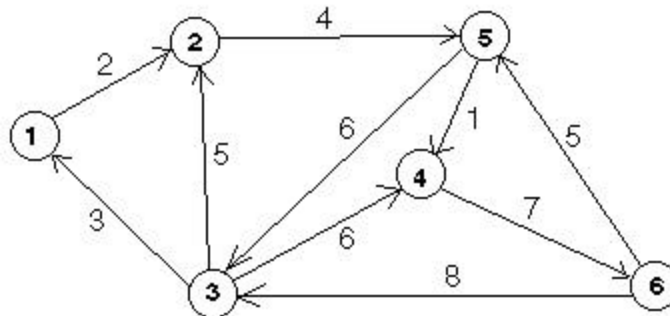


c. Add the new constraint in (b) to the tableau, and indicate by circling the next pivot:

|            | $x_1$ | $x_2$ | $x_3$ | $x_4$  | _____  | rhs   |
|------------|-------|-------|-------|--------|--------|-------|
|            | 0     | 0     | 0.571 | 2.143  | 36.429 | (max) |
|            | 0     | 1     | 0.071 | 0.643  | 6.429  |       |
| added row: | 1     | 0     | 0.071 | -0.357 | 1.429  |       |

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10. Chinese Postman Problem: Consider the directed graph below:



a. Compute the polarity of each node:

node: 1    2    3    4    5    6

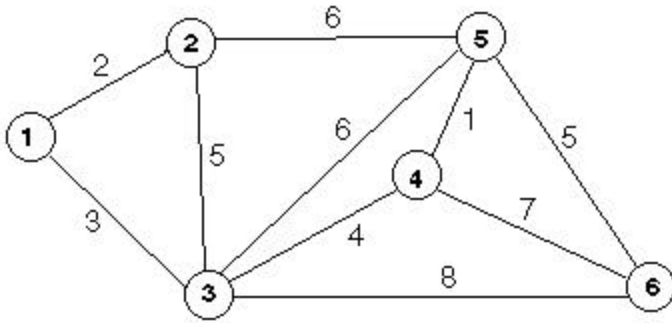
polarity: \_\_\_\_\_

b. Does this directed network possess an Euler tour? \_\_\_\_\_ ... an Euler path? \_\_\_\_\_

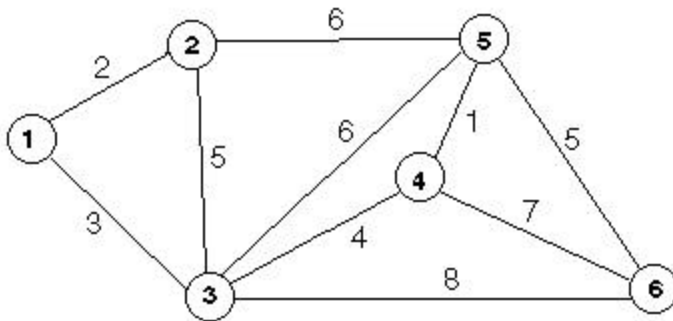
- c. In order for a delivery vehicle to travel each arc at least once with the minimum total distance traveled, which arcs should be traveled more than once?

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**11. Vertex Penalty Method.** Consider the network:



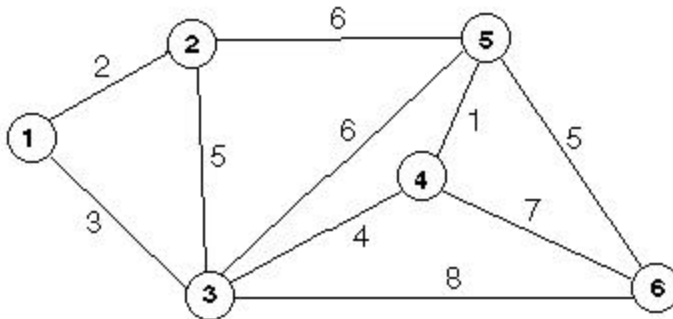
- Indicate the minimum spanning tree of this network.
- Indicate alongside each node its degree in the spanning tree.
- Suppose that, using the vertex penalty method, we wish to find the shortest path (not tour) starting at node 1, which passes through each node of this network. Perform one iteration of this method, using a unit penalty of 1, and indicate the new spanning tree below:



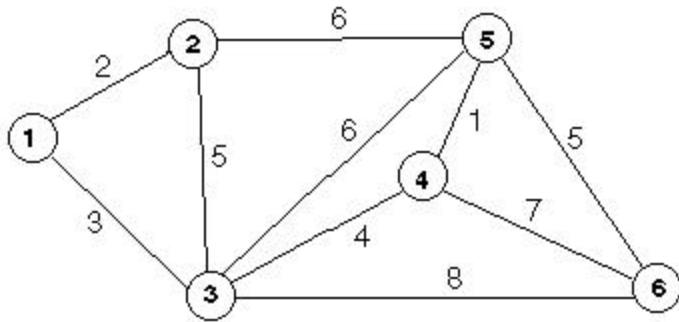
Is this spanning tree a path?

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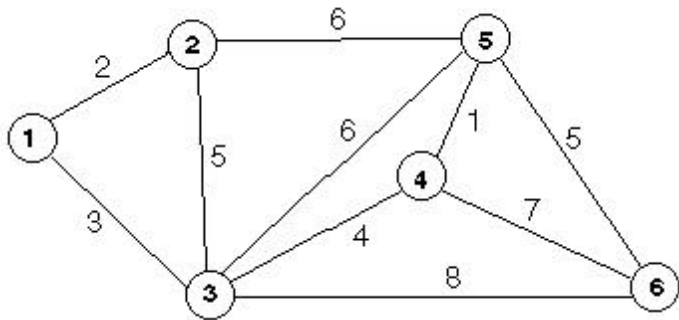
**12. Shortest Path Problem.** Consider again the network:



- It is desired to find the shortest path from node 1 to node 6, using the labeling algorithm of Dijkstra. Indicate the labels after the **first** iteration below:

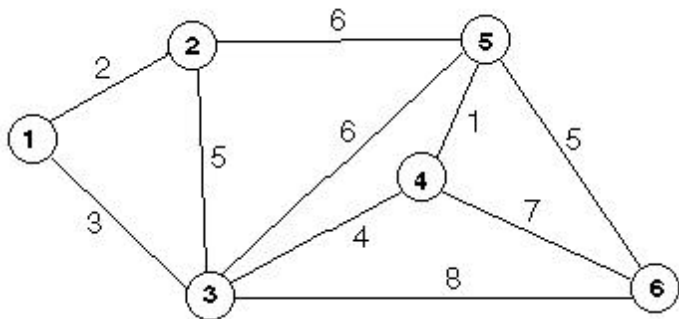


b. Indicate the labels after the second iteration below:



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13. **Location in a network:** Consider again the network,



where the numbers on the edges are distances. Consider the demand at each node to be 1. The table of shortest path lengths found by Floyd's algorithm is:

|    |    |   |   |   |    |
|----|----|---|---|---|----|
| 4  | 2  | 3 | 7 | 8 | 11 |
| 2  | 4  | 5 | 7 | 6 | 11 |
| 3  | 5  | 6 | 4 | 5 | 8  |
| 7  | 7  | 4 | 2 | 1 | 6  |
| 8  | 6  | 5 | 1 | 2 | 5  |
| 11 | 11 | 8 | 6 | 5 | 10 |

- At which node is the median (1-median) of the network? \_\_\_\_
- What is the objective function value of the median problem at this node? \_\_\_\_
- Consider the 2-median problem. What is the objective function value at the solution with nodes 3 and 4 selected? \_\_\_\_
- Which node is the vertex center (node center) of the network? \_\_\_\_
- What is the objective function of the center problem at this node? \_\_\_\_